The **OpenCL C Specification**

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Khronos OpenCL Working Group

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6. The OpenCL C Programming Language

This section describes the OpenCL C programming language used to create kernels that are executed on OpenCL device(s). The OpenCL C programming language (also referred to as OpenCL C) is based on the ISO/IEC 9899:1999 C language specification (a.k.a. C99 specification) with specific extensions and restrictions. Please refer to the ISO/IEC 9899:1999 specification for a detailed description of the language grammar. This section describes modifications and restrictions to ISO/IEC 9899:1999 supported in OpenCL C.

6.1 Supported Data Types

The following data types are supported.

6.1.1 Built-in Scalar Data Types

Table 6.1 describes the list of built-in scalar data types.

<table>
<thead>
<tr>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bool(^1)</td>
<td>A conditional data type which is either true or false. The value true expands to the integer constant 1 and the value false expands to the integer constant 0.</td>
</tr>
<tr>
<td>char</td>
<td>A signed two’s complement 8-bit integer.</td>
</tr>
<tr>
<td>unsigned char, uchar</td>
<td>An unsigned 8-bit integer.</td>
</tr>
<tr>
<td>short</td>
<td>A signed two’s complement 16-bit integer.</td>
</tr>
<tr>
<td>unsigned short, ushort</td>
<td>An unsigned 16-bit integer.</td>
</tr>
<tr>
<td>int</td>
<td>A signed two’s complement 32-bit integer.</td>
</tr>
<tr>
<td>unsigned int, uint</td>
<td>An unsigned 32-bit integer.</td>
</tr>
<tr>
<td>long</td>
<td>A signed two’s complement 64-bit integer.</td>
</tr>
<tr>
<td>unsigned long, ulong</td>
<td>An unsigned 64-bit integer.</td>
</tr>
<tr>
<td>float</td>
<td>A 32-bit floating-point. The float data type must conform to the IEEE 754 single precision storage format.</td>
</tr>
<tr>
<td>double(^2)</td>
<td>A 64-bit floating-point. The double data type must conform to the IEEE 754 double precision storage format.</td>
</tr>
<tr>
<td>half</td>
<td>A 16-bit floating-point. The half data type must conform to the</td>
</tr>
</tbody>
</table>

\(^1\) When any scalar value is converted to bool, the result is 0 if the value compares equal to 0; otherwise, the result is 1.
\(^2\) The double scalar type is an optional type that is supported if CL_DEVICE_DOUBLE_FP_CONFIG in table 4.3 for a device is not zero.
IEEE 754-2008 half precision storage format.

**size_t**  
The unsigned integer type of the result of the `sizeof` operator. This is a 32-bit unsigned integer if CL\_DEVICE\_ADDRESS\_BITS defined in table 4.3 is 32-bits and is a 64-bit unsigned integer if CL\_DEVICE\_ADDRESS\_BITS is 64-bits.

**ptrdiff_t**  
A signed integer type that is the result of subtracting two pointers. This is a 32-bit signed integer if CL\_DEVICE\_ADDRESS\_BITS defined in table 4.3 is 32-bits and is a 64-bit signed integer if CL\_DEVICE\_ADDRESS\_BITS is 64-bits.

**intptr_t**  
A signed integer type with the property that any valid pointer to `void` can be converted to this type, then converted back to pointer to `void`, and the result will compare equal to the original pointer. This is a 32-bit signed integer if CL\_DEVICE\_ADDRESS\_BITS defined in table 4.3 is 32-bits and is a 64-bit signed integer if CL\_DEVICE\_ADDRESS\_BITS is 64-bits.

**uintptr_t**  
An unsigned integer type with the property that any valid pointer to `void` can be converted to this type, then converted back to pointer to `void`, and the result will compare equal to the original pointer. This is a 32-bit signed integer if CL\_DEVICE\_ADDRESS\_BITS defined in table 4.3 is 32-bits and is a 64-bit signed integer if CL\_DEVICE\_ADDRESS\_BITS is 64-bits.

**void**  
The `void` type comprises an empty set of values; it is an incomplete type that cannot be completed.

### Table 6.1  Built-in Scalar Data Types

Most built-in scalar data types are also declared as appropriate types in the OpenCL API (and header files) that can be used by an application. The following table describes the built-in scalar data type in the OpenCL C programming language and the corresponding data type available to the application:

<table>
<thead>
<tr>
<th>Type in OpenCL Language</th>
<th>API type for application</th>
</tr>
</thead>
<tbody>
<tr>
<td>bool</td>
<td>n/a</td>
</tr>
<tr>
<td>char</td>
<td>cl_char</td>
</tr>
<tr>
<td>unsigned char,</td>
<td>cl_uchar</td>
</tr>
<tr>
<td>uchar</td>
<td></td>
</tr>
<tr>
<td>short</td>
<td>cl_short</td>
</tr>
<tr>
<td>unsigned short,</td>
<td>cl_ushort</td>
</tr>
<tr>
<td>ushort</td>
<td></td>
</tr>
<tr>
<td>int</td>
<td>cl_int</td>
</tr>
<tr>
<td>unsigned int,</td>
<td>cl_uint</td>
</tr>
<tr>
<td>uint</td>
<td></td>
</tr>
<tr>
<td>long</td>
<td>cl_long</td>
</tr>
<tr>
<td>unsigned long,</td>
<td>cl_ulong</td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### 6.1.1.1 The half data type

The half data type must be IEEE 754-2008 compliant. Half numbers have 1 sign bit, 5 exponent bits, and 10 mantissa bits. The interpretation of the sign, exponent and mantissa is analogous to IEEE 754 floating-point numbers. The exponent bias is 15. The half data type must represent finite and normal numbers, denormalized numbers, infinities and NaN. Denormalized numbers for the half data type which may be generated when converting a float to a half using vstore_half and converting a half to a float using vload_half cannot be flushed to zero. Conversions from float to half correctly round the mantissa to 11 bits of precision. Conversions from half to float are lossless; all half numbers are exactly representable as float values.

The half data type can only be used to declare a pointer to a buffer that contains half values. A few valid examples are given below:

```c
void
bar (__global half *p)
{
  ...
}

__kernel void
foo (__global half *pg, __local half *pl)
{
  __global half *ptr;
  int offset;

  ptr = pg + offset;
  bar(ptr);
}
```

Below are some examples that are not valid usage of the half type:

```c
half a;
half b[100];
```
half *p;
a = *p;  \textit{← not allowed. must use vload_half function}

Loads from a pointer to a half and stores to a pointer to a half can be performed using the vload_half, vload_halfn, vloada_halfn and vstore_half, vstore_halfn, vstorea_halfn functions respectively as described in section 6.1.7. The load functions read scalar or vector half values from memory and convert them to a scalar or vector float value. The store functions take a scalar or vector float value as input, convert it to a half scalar or vector value (with appropriate rounding mode) and write the half scalar or vector value to memory.

6.1.2 Built-in Vector Data Types$^3$

The char, unsigned char, short, unsigned short, integer, unsigned integer, long, unsigned long, float vector data types are supported. The vector data type is defined with the type name i.e. char, uchar, short, ushort, int, uint, float, long, ulong followed by a literal value $n$ that defines the number of elements in the vector. Supported values of $n$ are 2, 3, 4, 8, and 16 for all vector data types.

Table 6.2 describes the list of built-in vector data types.

<table>
<thead>
<tr>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>charn</td>
<td>A vector of $n$ 8-bit signed two’s complement integer values.</td>
</tr>
<tr>
<td>ucharn</td>
<td>A vector of $n$ 8-bit unsigned integer values.</td>
</tr>
<tr>
<td>shortn</td>
<td>A vector of $n$ 16-bit signed two’s complement integer values.</td>
</tr>
<tr>
<td>ushortn</td>
<td>A vector of $n$ 16-bit unsigned integer values.</td>
</tr>
<tr>
<td>intn</td>
<td>A vector of $n$ 32-bit signed two’s complement integer values.</td>
</tr>
<tr>
<td>uintn</td>
<td>A vector of $n$ 32-bit unsigned integer values.</td>
</tr>
<tr>
<td>longn</td>
<td>A vector of $n$ 64-bit signed two’s complement integer values.</td>
</tr>
<tr>
<td>ulongn</td>
<td>A vector of $n$ 64-bit unsigned integer values.</td>
</tr>
<tr>
<td>floatn</td>
<td>A vector of $n$ 32-bit floating-point values.</td>
</tr>
<tr>
<td>doublen$^4$</td>
<td>A vector of $n$ 64-bit floating-point values.</td>
</tr>
</tbody>
</table>

Table 6.2  Built-in Vector Data Types

The built-in vector data types are also declared as appropriate types in the OpenCL API (and header files) that can be used by an application. The following table describes the built-in vector

---

$^3$ Built-in vector data types are supported by the OpenCL implementation even if the underlying compute device does not support any or all of the vector data types. These are to be converted by the device compiler to appropriate instructions that use underlying built-in types supported natively by the compute device. Refer to Appendix B for a description of the order of the components of a vector type in memory.

$^4$ The double vector type is an optional type that is supported if CL_DEVICE_DOUBLE_FP_CONFIG in table 4.3 for a device is not zero.
data type in the OpenCL C programming language and the corresponding data type available to the application:

<table>
<thead>
<tr>
<th>Type in OpenCL Language</th>
<th>API type for application</th>
</tr>
</thead>
<tbody>
<tr>
<td>charn</td>
<td>cl_charn</td>
</tr>
<tr>
<td>ucharn</td>
<td>cl_uchar</td>
</tr>
<tr>
<td>shortn</td>
<td>cl_shortn</td>
</tr>
<tr>
<td>ushortn</td>
<td>cl_ushortn</td>
</tr>
<tr>
<td>intn</td>
<td>cl_intn</td>
</tr>
<tr>
<td>uintn</td>
<td>cl_uintn</td>
</tr>
<tr>
<td>longn</td>
<td>cl_longn</td>
</tr>
<tr>
<td>ulongn</td>
<td>cl_ulongn</td>
</tr>
<tr>
<td>floatn</td>
<td>cl_floatn</td>
</tr>
<tr>
<td>doublen</td>
<td>cl_doublen</td>
</tr>
</tbody>
</table>

6.1.3 Other Built-in Data Types

Table 6.3 describes the list of additional data types supported by OpenCL.

<table>
<thead>
<tr>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>image2d_t</td>
<td>A 2D image. Refer to section 6.13.14 for a detailed description of the built-in functions that use this type.</td>
</tr>
<tr>
<td>image3d_t</td>
<td>A 3D image. Refer to section 6.13.14 for a detailed description of the built-in functions that use this type.</td>
</tr>
<tr>
<td>image2d_array_t</td>
<td>A 2D image array. Refer to section 6.13.14 for a detailed description of the built-in functions that use this type.</td>
</tr>
<tr>
<td>image1d_t</td>
<td>A 1D image. Refer to section 6.13.14 for a detailed description of the built-in functions that use this type.</td>
</tr>
<tr>
<td>image1d_buffer_t</td>
<td>A 1D image created from a buffer object. Refer to section 6.13.14 for a detailed description of the built-in functions that use this type.</td>
</tr>
<tr>
<td>image1d_array_t</td>
<td>A 1D image array. Refer to section 6.13.14 for a detailed description of the built-in functions that use this type.</td>
</tr>
<tr>
<td>image2d_depth_t</td>
<td>A 2D depth image. Refer to section 6.13.14 for a detailed description of the built-in functions that use this type.</td>
</tr>
<tr>
<td>image2d_array_depth_t</td>
<td>A 2D depth image array. Refer to section 6.13.14 for a detailed description of the built-in functions that use this type.</td>
</tr>
<tr>
<td>sampler_t</td>
<td>A sampler type. Refer to section 6.13.14 for a detailed description the built-in functions that use of this type.</td>
</tr>
</tbody>
</table>
queue_t
A device command queue. This queue can only be used to enqueue commands from kernels executing on the device.

ndrange_t
The N-dimensional range over which a kernel executes.

clk_event_t
A device side event that identifies a command enqueue to a device command queue.

reserve_id_t
A reservation ID. This opaque type is used to identify the reservation for reading and writing a pipe. Refer to section 6.13.16.

event_t
An event. This can be used to identify async copies from global to local memory and vice-versa. Refer to section 6.13.10.

cl_mem_fence_flags
This is a bitfield and can be 0 or a combination of the following values ORed together:

- CLK_GLOBAL_MEM_FENCE
- CLK_LOCAL_MEM_FENCE
- CLK_IMAGE_MEM_FENCE

These flags are described in detail in section 6.13.8.

<table>
<thead>
<tr>
<th>Table 6.3</th>
<th>Other Built-in Data Types</th>
</tr>
</thead>
</table>

NOTE: The image2d_t, image3d_t, image2d_array_t, image1d_t, image1d_buffer_t, image1d_array_t, image2d_depth_t, image2d_array_depth_t and sampler_t types are only defined if the device supports images i.e. CL_DEVICE_IMAGE_SUPPORT as described in table 4.3 is CL_TRUE.

The C99 derived types (arrays, structs, unions, functions, and pointers), constructed from the built-in data types described in sections 6.1.1, 6.1.2 and 6.1.3 are supported, with restrictions described in section 6.9.

The following tables describe the other built-in data types in OpenCL described in table 6.3 and the corresponding data type available to the application:

<table>
<thead>
<tr>
<th>Type in OpenCL C</th>
<th>API type for application</th>
</tr>
</thead>
<tbody>
<tr>
<td>queue_t</td>
<td>cl_command_queue</td>
</tr>
<tr>
<td>clk_event_t</td>
<td>cl_event</td>
</tr>
</tbody>
</table>

6.1.4 Reserved Data Types

The data type names described in table 6.4 are reserved and cannot be used by applications as type names. The vector data type names defined in table 6.2, but where n is any value other than 2, 3, 4, 8 and 16, are also reserved.
<table>
<thead>
<tr>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>booln</td>
<td>A boolean vector.</td>
</tr>
<tr>
<td>halfn</td>
<td>A 16-bit floating-point vector.</td>
</tr>
<tr>
<td>quad, quadn</td>
<td>A 128-bit floating-point scalar and vector.</td>
</tr>
<tr>
<td>complex half,</td>
<td>A complex 16-bit floating-point scalar and vector.</td>
</tr>
<tr>
<td>complex halfn</td>
<td></td>
</tr>
<tr>
<td>imaginary half,</td>
<td>An imaginary 16-bit floating-point scalar and</td>
</tr>
<tr>
<td>imaginary halfn</td>
<td>vector.</td>
</tr>
<tr>
<td>complex float,</td>
<td>A complex 32-bit floating-point scalar and vector.</td>
</tr>
<tr>
<td>complex floatn</td>
<td></td>
</tr>
<tr>
<td>imaginary float,</td>
<td>An imaginary 32-bit floating-point scalar and</td>
</tr>
<tr>
<td>imaginary floatn</td>
<td>vector.</td>
</tr>
<tr>
<td>complex double,</td>
<td>A complex 64-bit floating-point scalar and vector.</td>
</tr>
<tr>
<td>complex doublen,</td>
<td></td>
</tr>
<tr>
<td>imaginary double,</td>
<td>An imaginary 64-bit floating-point scalar and</td>
</tr>
<tr>
<td>imaginary doublen</td>
<td>vector.</td>
</tr>
<tr>
<td>complex quad,</td>
<td>A complex 128-bit floating-point scalar and vector.</td>
</tr>
<tr>
<td>complex quadn,</td>
<td></td>
</tr>
<tr>
<td>imaginary quad,</td>
<td>An imaginary 128-bit floating-point scalar and</td>
</tr>
<tr>
<td>imaginary quadn</td>
<td>vector.</td>
</tr>
<tr>
<td>floatnmxm</td>
<td>An n x m matrix of single precision floating-</td>
</tr>
<tr>
<td></td>
<td>point values stored in column-major order.</td>
</tr>
<tr>
<td>doublenmxm</td>
<td>An n x m matrix of double precision floating-</td>
</tr>
<tr>
<td></td>
<td>point values stored in column-major order.</td>
</tr>
<tr>
<td>long double</td>
<td>A floating-point scalar and vector type with at</td>
</tr>
<tr>
<td>long doublen</td>
<td>least as much precision and range as a double</td>
</tr>
<tr>
<td></td>
<td>and no more precision and range than a quad.</td>
</tr>
<tr>
<td>long long, long longn</td>
<td>A 128-bit signed integer scalar and vector.</td>
</tr>
<tr>
<td>ulong long, ulong longn</td>
<td>A 128-bit unsigned integer scalar and vector.</td>
</tr>
</tbody>
</table>

Table 6.4  
Reserved Data Types

### 6.1.5 Alignment of Types

A data item declared to be a data type in memory is always aligned to the size of the data type in bytes. For example, a float4 variable will be aligned to a 16-byte boundary, a char2 variable will be aligned to a 2-byte boundary.

For 3-component vector data types, the size of the data type is $4 \times \text{sizeof(component)}$. This means that a 3-component vector data type will be aligned to a $4 \times \text{sizeof(component)}$ boundary. The `vload3` and `vstore3` built-in functions can be used to read and write, respectively, 3-component vector data types from an array of packed scalar data type.
A built-in data type that is not a power of two bytes in size must be aligned to the next larger power of two. This rule applies to built-in types only, not structs or unions.

The OpenCL compiler is responsible for aligning data items to the appropriate alignment as required by the data type. For arguments to a __kernel function declared to be a pointer to a data type, the OpenCL compiler can assume that the pointee is always appropriately aligned as required by the data type. The behavior of an unaligned load or store is undefined, except for the vloadn, vload_halfn, vstoren, and vstore_halfn functions defined in section 6.13.7. The vector load functions can read a vector from an address aligned to the element type of the vector. The vector store functions can write a vector to an address aligned to the element type of the vector.

### 6.1.6 Vector Literals

Vector literals can be used to create vectors from a list of scalars, vectors or a mixture thereof. A vector literal can be used either as a vector initializer or as a primary expression. A vector literal cannot be used as an L-value.

A vector literal is written as a parenthesized vector type followed by a parenthesized comma delimited list of parameters. A vector literal operates as an overloaded function. The forms of the function that are available is the set of possible argument lists for which all arguments have the same element type as the result vector, and the total number of elements is equal to the number of elements in the result vector. In addition, a form with a single scalar of the same type as the element type of the vector is available. For example, the following forms are available for float4:

```plaintext
(float4)( float, float, float, float )
(float4)( float2, float, float )
(float4)( float, float2, float )
(float4)( float, float, float2 )
(float4)( float2, float2 )
(float4)( float3, float )
(float4)( float, float3 )
(float4)( float )
```

Operands are evaluated by standard rules for function evaluation, except that implicit scalar widening shall not occur. The order in which the operands are evaluated is undefined. The operands are assigned to their respective positions in the result vector as they appear in memory order. That is, the first element of the first operand is assigned to result.x, the second element of the first operand (or the first element of the second operand if the first operand was a scalar) is assigned to result.y, etc. In the case of the form that has a single scalar operand, the operand is replicated across all lanes of the vector.

Examples:
float4  f = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
uint4  u = (uint4)(1);  //</u will be (1, 1, 1, 1).
float4  f = (float4)((float2)(1.0f, 2.0f),
                     (float2)(3.0f, 4.0f));
float4  f = (float4)(1.0f, (float2)(2.0f, 3.0f), 4.0f);
float4  f = (float4)(1.0f, 2.0f);  //</error

6.1.7 Vector Components

The components of vector data types with 1 ... 4 components can be addressed as
<vector_data_type>.xyzw. Vector data types of type char2, uchar2, short2,
ushort2, int2, uint2, long2, ulong2, and float2 can access .xy elements. Vector data types of type
char3, uchar3, short3, ushort3, int3, uint3,
long3, ulong3, and float3 can access .xyz elements. Vector data types of type
char4, uchar4, short4, ushort4, int4, uint4, long4, ulong4,
float4 can access .xyzw elements.

Accessing components beyond those declared for the vector type is an error so, for example:

    float2 pos;
    pos.x = 1.0f;  // is legal
    pos.z = 1.0f;  // is illegal

    float3 pos;
    pos.z = 1.0f;  // is legal
    pos.w = 1.0f;  // is illegal

The component selection syntax allows multiple components to be selected by appending their
names after the period ( . ).

    float4 c;
    c.xyzw = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
    c.z = 1.0f;
    c.xy = (float2)(3.0f, 4.0f);
    c.xyz = (float3)(3.0f, 4.0f, 5.0f);

The component selection syntax also allows components to be permuted or replicated.

    float4 pos = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
```
float4 swiz = pos.wzyx;  // swiz = (4.0f, 3.0f, 2.0f, 1.0f)
float4 dup = pos.xxyy;  // dup = (1.0f, 1.0f, 2.0f, 2.0f)
```

The component group notation can occur on the left hand side of an expression. To form an l-value, swizzling must be applied to an l-value of vector type, contain no duplicate components, and it results in an l-value of scalar or vector type, depending on number of components specified. Each component must be a supported scalar or vector type.

```
float4 pos = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
pos.xw = (float2)(5.0f, 6.0f);  // pos = (5.0f, 2.0f, 3.0f, 6.0f)
pos.wx = (float2)(7.0f, 8.0f);  // pos = (8.0f, 2.0f, 3.0f, 7.0f)
pos.xyz = (float3)(3.0f, 5.0f, 9.0f);  // pos = (3.0f, 5.0f, 9.0f, 4.0f)
pos.xx = (float2)(3.0f, 4.0f);  // illegal - 'x' used twice
```

Elements of vector data types can also be accessed using a numeric index to refer to the appropriate element in the vector. The numeric indices that can be used are given in the table below:

<table>
<thead>
<tr>
<th>Vector Components</th>
<th>Numeric indices that can be used</th>
</tr>
</thead>
<tbody>
<tr>
<td>2-component</td>
<td>0, 1</td>
</tr>
<tr>
<td>3-component</td>
<td>0, 1, 2</td>
</tr>
<tr>
<td>4-component</td>
<td>0, 1, 2, 3</td>
</tr>
<tr>
<td>8-component</td>
<td>0, 1, 2, 3, 4, 5, 6, 7</td>
</tr>
<tr>
<td>16-component</td>
<td>0, 1, 2, 3, 4, 5, 6, 7, 8, 9, a, a, b, B, c, C, d, D, e, E, f, F</td>
</tr>
</tbody>
</table>

Table 6.5  Numeric indices for built-in vector data types

The numeric indices must be preceded by the letter s or S.

In the following example

```
float4 a, b, c, d;
float16 x;
x = (float16)(a.xxxx, b.xyz, c.xyz, d.xyz, a.yzw);
```

// illegal - component a.xxxxxxx is not a valid vector type
```
x = (float16)(a.xxxxxxx, b.xyz, c.xyz, d.xyz);
```
float8     f;

f.s0 refers to the 1st element of the float8 variable f and f.s7 refers to the 8th element of the float8 variable f.

In the following example

    float16     x;

x.sA (or x.sA) refers to the 11th element of the float16 variable x and x.sf (or x.sF) refers to the 16th element of the float16 variable x.

The numeric indices used to refer to an appropriate element in the vector cannot be intermixed with .xyzw notation used to access elements of a 1 .. 4 component vector.

For example

    float4     f, a;

    a = f.x12w;    // illegal use of numeric indices with .xyzw

    a.xyzw = f.s0123;    // valid

Vector data types can use the .lo (or .even) and .hi (or .odd) suffixes to get smaller vector types or to combine smaller vector types to a larger vector type. Multiple levels of .lo (or .even) and .hi (or .odd) suffixes can be used until they refer to a scalar term.

The .lo suffix refers to the lower half of a given vector. The .hi suffix refers to the upper half of a given vector.

The .even suffix refers to the even elements of a vector. The .odd suffix refers to the odd elements of a vector.

Some examples to help illustrate this are given below:

    float4     vf;

    float2    low = vf.lo;    // returns vf.xy
    float2    high = vf.hi;    // returns vf.zw

    float2    even = vf.even;    // returns vf.xz
    float2    odd = vf.odd;    // returns vf.yw

The suffixes .lo (or .even) and .hi (or .odd) for a 3-component vector type operate as if the 3-component vector type is a 4-component vector type with the value in the w component undefined.
Some examples are given below:

```plaintext
float8   vf;
float4    odd = vf.odd;
float4    even = vf.even;
float2    high = vf.even.hi;
float2    low = vf.odd.lo;

// interleave L+R stereo stream
float4    left, right;
float8    interleaved;
interleaved.even = left;
interleaved.odd = right;

// deinterleave
left = interleaved.even;
right = interleaved.odd;

// transpose a 4x4 matrix
void transpose( float4 m[4] )
{
    // read matrix into a float16 vector
    float16 x = (float16)( m[0], m[1], m[2], m[3] );
    float16 t;

    //transpose
    t.even = x.lo;
    t.odd = x.hi;
    x.even = t.lo;
    x.odd = t.hi;

    //write back
    m[0] = x.lo.lo;  // { m[0][0], m[1][0], m[2][0], m[3][0] }
    m[1] = x.lo.hi;  // { m[0][1], m[1][1], m[2][1], m[3][1] }
    m[2] = x.hi.lo;  // { m[0][2], m[1][2], m[2][2], m[3][2] }
    m[3] = x.hi.hi;  // { m[0][3], m[1][3], m[2][3], m[3][3] }
}

float3    vf = (float3)(1.0f, 2.0f, 3.0f);
float2    low = vf.lo;  // (1.0f, 2.0f);
float2    high = vf.hi;  // (3.0f, undefined);
```

It is an error to take the address of a vector element and will result in a compilation error. For example:

```plaintext
float8    vf;
float    *f = &vf.x;  // is illegal
float2   *f2 = &vf.s07;  // is illegal
float4   *odd = &vf.odd;  // is illegal
float4   *even = &vf.even;  // is illegal
float2   *high = &vf.even.hi;  // is illegal
float2   *low = &vf.odd.lo;  // is illegal
```
6.1.8 Aliasing Rules

OpenCL C programs shall comply with the C99 type-based aliasing rules (defined in section 6.5, item 7 of the C99 specification). The OpenCL C built-in vector data types are considered aggregate\(^5\) types for the purpose of applying these aliasing rules.

6.1.9 Keywords

The following names are reserved for use as keywords in OpenCL C and shall not be used otherwise.

- Names reserved as keywords by C99.
- OpenCL C data types defined in tables 6.2, 6.3 and 6.4.
- Address space qualifiers: `__global, global, __local, local, __constant, constant, __private and private`.
- Function qualifiers: `__kernel` and `kernel`.
- Access qualifiers: `__read_only, read_only, __write_only, write_only, __read_write` and `read_write`.
- uniform, pipe.

\(^5\) That is, for the purpose of applying type-based aliasing rules, a built-in vector data type will be considered equivalent to the corresponding array type.
6.2 Conversions and Type Casting

6.2.1 Implicit Conversions

Implicit conversions between scalar built-in types defined in table 6.1 (except void and half) are supported. When an implicit conversion is done, it is not just a re-interpretation of the expression's value but a conversion of that value to an equivalent value in the new type. For example, the integer value 5 will be converted to the floating-point value 5.0.

Implicit conversions from a scalar type to a vector type are allowed. In this case, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector. The scalar type is then widened to the vector.

Implicit conversions between built-in vector data types are disallowed.

Implicit conversions for pointer types follow the rules described in the C99 specification.

6.2.2 Explicit Casts

Standard typecasts for built-in scalar data types defined in table 6.1 will perform appropriate conversion (except void and half). In the example below:

```c
float f = 1.0f;
int i = (int)f;
```

f stores 0x3F800000 and i stores 0x1 which is the floating-point value 1.0f in f converted to an integer value.

Explicit casts between vector types are not legal. The examples below will generate a compilation error.

```c
int4 i;
uint4 u = (uint4) i;  // not allowed

float4 f;
int4 i = (int4) f;   // not allowed

float4 f;
int8 i = (int8) f;  // not allowed
```

Scalar to vector conversions may be performed by casting the scalar to the desired vector data

---

6 Unless the cl_khr_fp16 extension is supported.
7 Unless the cl_khr_fp16 extension is supported.
Type casting will also perform appropriate arithmetic conversion. The round to zero rounding mode will be used for conversions to built-in integer vector types. The default rounding mode will be used for conversions to floating-point vector types. When casting a bool to a vector integer data type, the vector components will be set to -1 (i.e. all bits set) if the bool value is true and 0 otherwise.

Below are some correct examples of explicit casts.

```cpp
float f = 1.0f;
float4 va = (float4)f;
// va is a float4 vector with elements (f, f, f, f).
uchar u = 0xFF;
float4 vb = (float4)u;
// vb is a float4 vector with elements((float)u, (float)u, (float)u, (float)u).
float f = 2.0f;
int2 vc = (int2)f;
// vc is an int2 vector with elements ((int)f, (int)f).
uchar4 vtrue = (uchar4)true;
// vtrue is a uchar4 vector with elements (0xff, 0xff, 0xff, 0xff).
```

### 6.2.3 Explicit Conversions

Explicit conversions may be performed using the

`convert_destType(sourceType)`

suite of functions. These provide a full set of type conversions between supported types (see sections 6.1.1, 6.1.2 and 6.1.3) except for the following types: bool, half, size_t, ptrdiff_t, intptr_t, uintptr_t, and void.

The number of elements in the source and destination vectors must match.

In the example below:

```cpp
uchar4 u;
int4 c = convert_int4(u);
```

`convert_int4` converts a uchar4 vector `u` to an int4 vector `c`. 
float   f;
int     i = convert_int(f);

convert_int converts a float scalar f to an int scalar i.

The behavior of the conversion may be modified by one or two optional modifiers that specify saturation for out-of-range inputs and rounding behavior.

The full form of the scalar convert function is:

```
destType convert_destType<_sat><_roundingMode> (sourceType)
```

The full form of the vector convert function is:

```
destTypeN convert_destTypeN<_sat><_roundingMode> (sourceTypeN)
```

### 6.2.3.1 Data Types

Conversions are available for the following scalar types: char, uchar, short, ushort, int, uint, long, ulong, float, and built-in vector types derived therefrom. The operand and result type must have the same number of elements. The operand and result type may be the same type in which case the conversion has no effect on the type or value of an expression.

Conversions between integer types follow the conversion rules specified in sections 6.3.1.1 and 6.3.1.3 of the C99 specification except for out-of-range behavior and saturated conversions which are described in section 6.2.3.3 below.

### 6.2.3.2 Rounding Modes

Conversions to and from floating-point type shall conform to IEEE-754 rounding rules. Conversions may have an optional rounding mode modifier described in table 6.6.

<table>
<thead>
<tr>
<th>Modifier</th>
<th>Rounding Mode Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>_rte</td>
<td>Round to nearest even</td>
</tr>
<tr>
<td>_rtz</td>
<td>Round toward zero</td>
</tr>
<tr>
<td>_rtp</td>
<td>Round toward positive infinity</td>
</tr>
<tr>
<td>_rtn</td>
<td>Round toward negative infinity</td>
</tr>
<tr>
<td>no modifier specified</td>
<td>Use the default rounding mode for this destination type, _rtz for conversion to integers or the default rounding mode for conversion to floating-point types.</td>
</tr>
</tbody>
</table>

**Table 6.6  Rounding Modes**
By default, conversions to integer type use the _rtz (round toward zero) rounding mode and conversions to floating-point type\(^8\) use the default rounding mode. The only default floating-point rounding mode supported is round to nearest even i.e the default rounding mode will be _rte for floating-point types.

### 6.2.3.3 Out-of-Range Behavior and Saturated Conversions

When the conversion operand is either greater than the greatest representable destination value or less than the least representable destination value, it is said to be out-of-range. The result of out-of-range conversion is determined by the conversion rules specified by the C99 specification in section 6.3. When converting from a floating-point type to integer type, the behavior is implementation-defined.

Conversions to integer type may opt to convert using the optional saturated mode by appending the _sat modifier to the conversion function name. When in saturated mode, values that are outside the representable range shall clamp to the nearest representable value in the destination format. (NaN should be converted to 0).

Conversions to floating-point type shall conform to IEEE-754 rounding rules. The _sat modifier may not be used for conversions to floating-point formats.

### 6.2.3.4 Explicit Conversion Examples

Example 1:

```c
short4 s;

// negative values clamped to 0
ushort4 u = convert_ushort4_sat(s);

// values > CHAR_MAX converted to CHAR_MAX
// values < CHAR_MIN converted to CHAR_MIN
char4 c = convert_char4_sat(s);
```

Example 2:

```c
float4 f;

// values implementation defined for
// f > INT_MAX, f < INT_MIN or NaN
int4 i = convert_int4(f);
```

\(^8\) For conversions to floating-point format, when a finite source value exceeds the maximum representable finite floating-point destination value, the rounding mode will affect whether the result is the maximum finite floating-point value or infinity of same sign as the source value, per IEEE-754 rules for rounding.
// values > INT_MAX clamp to INT_MAX, values < INT_MIN clamp
// to INT_MIN. NaN should produce 0.
// The _rtz rounding mode is used to produce the integer values.
int4 i2 = convert_int4_sat( f );

// similar to convert_int4, except that floating-point values
// are rounded to the nearest integer instead of truncated
int4 i3 = convert_int4_rte( f );

// similar to convert_int4_sat, except that floating-point values
// are rounded to the nearest integer instead of truncated
int4 i4 = convert_int4_sat_rte( f );

Example 3:

int4 i;

// convert ints to floats using the default rounding mode.
float4 f = convert_float4( i );

// convert ints to floats. integer values that cannot
// be exactly represented as floats should round up to the
// next representable float.
float4 f = convert_float4_rtp( i );

6.2.4 Reinterpreting Data As Another Type

It is frequently necessary to reinterpret bits in a data type as another data type in OpenCL. This is typically required when direct access to the bits in a floating-point type is needed, for example to mask off the sign bit or make use of the result of a vector relational operator (see section 6.3.d) on floating-point data. Several methods to achieve this (non-) conversion are frequently practiced in C, including pointer aliasing, unions and memcpy. Of these, only memcpy is strictly correct in C99. Since OpenCL does not provide memcpy, other methods are needed.

6.2.4.1 Reinterpreting Types Using Unions

The OpenCL language extends the union to allow the program to access a member of a union object using a member of a different type. The relevant bytes of the representation of the object

---

9 In addition, some other extensions to the C language designed to support particular vector ISA (e.g. AltiVec™, CELL Broadband Engine™ Architecture) use such conversions in conjunction with swizzle operators to achieve type unconversion. So as to support legacy code of this type, as_typed() allows conversions between vectors of the same size but different numbers of elements, even though the behavior of this sort of conversion is not likely to be portable except to other OpenCL implementations for the same hardware architecture. AltiVec™ is a trademark of Motorola Inc. Cell Broadband Engine is a trademark of Sony Computer Entertainment, Inc.
are treated as an object of the type used for the access. If the type used for access is larger than
the representation of the object, then the value of the additional bytes is undefined.

Examples:

    union{ float f; uint u; double d10;} u;
    u.u = 1; // u.f contains 2**-149. u.d is undefined --
               // depending on endianness the low or high half
               // of d is unknown
    u.f = 1.0f; // u.u contains 0x3f800000, u.d contains an
               // undefined value -- depending on endianness
               // the low or high half of d is unknown
    u.d = 1.0; // u.u contains 0x3ff00000 (big endian) or 0
               // (little endian). u.f contains either 0x1.ep0f
               // (big endian) or 0.0f (little endian)

6.2.4.2 Reinterpreting Types Using as_type() and as_typeen()

All data types described in tables 6.1 and 6.2 (except bool, half11 and void) may be also
reinterpreted as another data type of the same size using the as_type() operator for scalar data
types and the as_typeen() operator12 for vector data types. When the operand and result type
contain the same number of elements, the bits in the operand shall be returned directly without
modification as the new type. The usual type promotion for function arguments shall not be
performed.

For example, as_float(0x3f800000) returns 1.0f, which is the value that the bit pattern
0x3f800000 has if viewed as an IEEE-754 single precision value.

When the operand and result type contain a different number of elements, the result shall be
implementation-defined except if the operand is a 4-component vector and the result is a 3-
component vector. In this case, the bits in the operand shall be returned directly without

---

10 Only if double precision is supported.
11 Unless the cl_khr_fp16 extension is supported.
12 While the union is intended to reflect the organization of data in memory, the as_type() and as_typeen() constructs
   are intended to reflect the organization of data in register. The as_type() and as_typeen() constructs are intended to
   compile to no instructions on devices that use a shared register file designed to operate on both the operand and
   result types. Note that while differences in memory organization are expected to largely be limited to those arising
   from endianness, the register based representation may also differ due to size of the element in register. (For
   example, an architecture may load a char into a 32-bit register, or a char vector into a SIMD vector register with
   fixed 32-bit element size.) If the element count does not match, then the implementation should pick a data
   representation that most closely matches what would happen if an appropriate result type operator was applied to a
   register containing data of the source type. If the number of elements matches, then the as_typeen() should faithfully
   reproduce the behavior expected from a similar data type reinterpretation using memory/unions. So, for example if
   an implementation stores all single precision data as double in register, it should implement as_int( float ) by first
downconverting the double to single precision and then (if necessary) moving the single precision bits to a register
suitable for operating on integer data. If data stored in different address spaces do not have the same endianness,
then the "dominant endianness" of the device should prevail.
modification as the new type. That is, a conforming implementation shall explicitly define a behavior, but two conforming implementations need not have the same behavior when the number of elements in the result and operand types does not match. The implementation may define the result to contain all, some or none of the original bits in whatever order it chooses. It is an error to use as_type() or as_type_n() operator to reinterpret data to a type of a different number of bytes.

Examples:

```c
float f = 1.0f;
uint u = as_uint(f); // Legal. Contains: 0x3f800000

float4 f = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
// Legal. Contains:
// (int4)(0x3f800000, 0x40000000, 0x40400000, 0x40800000)
int4 i = as_int4(f);

float4 f, g;
int4 is_less = f < g;
// Legal. f[i] = f[i] < g[i] ? f[i] : 0.0f
f = as_float4(as_int4(f) & is_less);

int i;
// Legal. Result is implementation-defined.
short2 j = as_short2(i);

int4 i;
// Legal. Result is implementation-defined.
short8 j = as_short8(i);

float4 f;
// Error. Result and operand have different sizes
double4 g = as_double413(f);

float4 f;
// Legal. g.xyz will have same values as f.xyz. g.w is undefined
float3 g = as_float3(f);
```

### 6.2.5 Pointer Casting

Pointers to old and new types may be cast back and forth to each other. Casting a pointer to a new type represents an unchecked assertion that the address is correctly aligned. The developer will also need to know the endianness of the OpenCL device and the endianness of the data to determine how the scalar and vector data elements are stored in memory.

---

13 Only if double precision is supported.
6.2.6 Usual Arithmetic Conversions

Many operators that expect operands of arithmetic type cause conversions and yield result types in a similar way. The purpose is to determine a common real type for the operands and result. For the specified operands, each operand is converted, without change of type domain, to a type whose corresponding real type is the common real type. For this purpose, all vector types shall be considered to have higher conversion ranks than scalars. Unless explicitly stated otherwise, the common real type is also the corresponding real type of the result, whose type domain is the type domain of the operands if they are the same, and complex otherwise. This pattern is called the usual arithmetic conversions. If the operands are of more than one vector type, then an error shall occur. Implicit conversions between vector types are not permitted, per section 6.2.1.

Otherwise, if there is only a single vector type, and all other operands are scalar types, the scalar types are converted to the type of the vector element, then widened into a new vector containing the same number of elements as the vector, by duplication of the scalar value across the width of the new vector. An error shall occur if any scalar operand has greater rank than the type of the vector element. For this purpose, the rank order defined as follows:

1. The rank of a floating-point type is greater than the rank of another floating-point type, if the first floating-point type can exactly represent all numeric values in the second floating-point type. (For this purpose, the encoding of the floating-point value is used, rather than the subset of the encoding usable by the device.)
2. The rank of any floating-point type is greater than the rank of any integer type.
3. The rank of an integer type is greater than the rank of an integer type with less precision.
4. The rank of an unsigned integer type is greater than the rank of a signed integer type with the same precision.\(^{14}\)
5. The rank of the bool type is less than the rank of any other type.
6. The rank of an enumerated type shall equal the rank of the compatible integer type.
7. For all types, \(T_1, T_2 \) and \(T_3\), if \(T_1\) has greater rank than \(T_2\), and \(T_2\) has greater rank than \(T_3\), then \(T_1\) has greater rank than \(T_3\).

Otherwise, if all operands are scalar, the usual arithmetic conversions apply, per section 6.3.1.8 of the C99 standard.

NOTE: Both the standard orderings in sections 6.3.1.8 and 6.3.1.1 of C99 were examined and rejected. Had we used integer conversion rank here, \(\text{int4} \, + \, \text{0U}\) would have been legal and had \(\text{int4}\) return type. Had we used standard C99 usual arithmetic conversion rules for scalars, then the standard integer promotion would have been performed on vector integer element types and \(\text{short8} \, + \, \text{char}\) would either have return type of \(\text{int8}\) or be illegal.

\(^{14}\) This is different from the standard integer conversion rank described in C99 TC2, section 6.3.1.1.
6.3 Operators

a. The arithmetic operators add (+), subtract (-), multiply (*) and divide (/) operate on built-in integer and floating-point scalar, and vector data types. The remainder (%) operates on built-in integer scalar and integer vector data types. All arithmetic operators return result of the same built-in type (integer or floating-point) as the type of the operands, after operand type conversion. After conversion, the following cases are valid:

- The two operands are scalars. In this case, the operation is applied, resulting in a scalar.
- One operand is a scalar, and the other is a vector. In this case, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.
- The two operands are vectors of the same type. In this case, the operation is done component-wise resulting in the same size vector.

All other cases of implicit conversions are illegal. Division on integer types which results in a value that lies outside of the range bounded by the maximum and minimum representable values of the integer type will not cause an exception but will result in an unspecified value. A divide by zero with integer types does not cause an exception but will result in an unspecified value. Division by zero for floating-point types will result in ±infinity or NaN as prescribed by the IEEE-754 standard. Use the built-in functions dot and cross to get, respectively, the vector dot product and the vector cross product.

b. The arithmetic unary operators (+ and −) operate on built-in scalar and vector types.

c. The arithmetic post- and pre-increment and decrement operators (−− and ++) operate on built-in scalar and vector types except the built-in scalar and vector float types. All unary operators work component-wise on their operands. These result with the same type they operated on. For post- and pre-increment and decrement, the expression must be one that could be assigned to (an l-value). Pre-increment and pre-decrement add or subtract 1 to the contents of the expression they operate on, and the value of the pre-increment or pre-decrement expression is the resulting value of that modification. Post-increment and post-decrement expressions add or subtract 1 to the contents of the expression they operate on, but the resulting expression has the expression’s value before the post-increment or post-decrement was executed.

---

15 The pre- and post- increment operators may have unexpected behavior on floating-point values and are therefore not supported for floating-point scalar and vector built-in types. For example, if variable a has type float and holds the value 0x1.0p25f, then a++ returns 0x1.0p25f. Also, (a++)-- is not guaranteed to return a, if a has fractional value. In non-default rounding modes, (a++)-- may produce the same result as a++ or a-- for large a.
d. The relational operators\(^\text{16}\) greater than (>), less than (<), greater than or equal (\(\geq\)), and less than or equal (\(\leq\)) operate on scalar and vector types. All relational operators result in an integer type. After operand type conversion, the following cases are valid:

- The two operands are scalars. In this case, the operation is applied, resulting in an `int` scalar.
- One operand is a scalar, and the other is a vector. In this case, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.
- The two operands are vectors of the same type. In this case, the operation is done component-wise resulting in the same size vector.

All other cases of implicit conversions are illegal.

The result is a scalar signed integer of type `int` if the source operands are scalar and a vector signed integer type of the same size as the source operands if the source operands are vector types. Vector source operands of type `char` and `uchar` return a `char` result; vector source operands of type `short` and `ushort` return a `short` result; vector source operands of type `int`, `uint` and `float` return an `int` result; vector source operands of type `long`, `ulong` and `double` return a `long` result. For scalar types, the relational operators shall return 0 if the specified relation is `false` and 1 if the specified relation is `true`. For vector types, the relational operators shall return 0 if the specified relation is `false` and −1 (i.e. all bits set) if the specified relation is `true`. The relational operators always return 0 if either argument is not a number (NaN).

e. The equality operators\(^\text{17}\) equal (\(==\)), and not equal (\(!=\)) operate on built-in scalar and vector types. All equality operators result in an integer type. After operand type conversion, the following cases are valid:

- The two operands are scalars. In this case, the operation is applied, resulting in a scalar.
- One operand is a scalar, and the other is a vector. In this case, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in

---

\(^{16}\) To test whether any or all elements in the result of a vector relational operator test true, for example to use in the context in an `if ( )` statement, please see the `any` and `all` builtins in section 6.11.6.

\(^{17}\) To test whether any or all elements in the result of a vector equality operator test true, for example to use in the context in an `if ( )` statement, please see the `any` and `all` builtins in section 6.11.6.
The two operands are vectors of the same type. In this case, the operation is done component-wise resulting in the same size vector.

All other cases of implicit conversions are illegal.

The result is a scalar signed integer of type `int` if the source operands are scalar and a vector signed integer type of the same size as the source operands if the source operands are vector types. Vector source operands of type `char n` and `uchar n` return a `char n` result; vector source operands of type `short n` and `ushort n` return a `short n` result; vector source operands of type `int n`, `uint n` and `float n` return an `int n` result; vector source operands of type `long n`, `ulong n` and `double n` return a `long n` result.

For scalar types, the equality operators return 0 if the specified relation is `false` and return 1 if the specified relation is `true`. For vector types, the equality operators shall return 0 if the specified relation is `false` and -1 (i.e. all bits set) if the specified relation is `true`. The equality operator equal (==) returns 0 if one or both arguments are not a number (NaN). The equality operator not equal (!=) returns 1 (for scalar source operands) or -1 (for vector source operands) if one or both arguments are not a number (NaN).

f. The bitwise operators and (&), or (|), exclusive or (^), not (~) operate on all scalar and vector built-in types except the built-in scalar and vector float types. For vector built-in types, the operators are applied component-wise. If one operand is a scalar and the other is a vector, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.

g. The logical operators and (&&), or (||) operate on all scalar and vector built-in types. For scalar built-in types only, and (&&) will only evaluate the right hand operand if the left hand operand compares unequal to 0. For scalar built-in types only, or (||) will only evaluate the right hand operand if the left hand operand compares equal to 0. For built-in vector types, both operands are evaluated and the operators are applied component-wise. If one operand is a scalar and the other is a vector, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.

The logical operator exclusive or (^^) is reserved.

The result is a scalar signed integer of type `int` if the source operands are scalar and a vector signed integer type of the same size as the source operands if the source operands are vector types. Vector source operands of type `char n` and `uchar n` return a `char n` result; vector source operands of type `short n` and `ushort n` return a `short n` result; vector source operands of type `int n`, `uint n` and `float n` return an `int n` result; vector source operands of type `long n`, `ulong n` and `double n` return a `long n` result.
operands of type int\(n\), uint\(n\) and float\(n\) return an int\(n\) result; vector source operands of type long\(n\), ulong\(n\) and double\(n\) return a long\(n\) result.

For scalar types, the logical operators shall return 0 if the result of the operation is false and 1 if the result is true. For vector types, the logical operators shall return 0 if the result of the operation is false and -1 (i.e. all bits set) if the result is true.

h. The logical unary operator not (!) operates on all scalar and vector built-in types. For built-in vector types, the operators are applied component-wise.

The result is a scalar signed integer of type int if the source operands are scalar and a vector signed integer type of the same size as the source operands if the source operands are vector types. Vector source operands of type char\(n\) and uchar\(n\) return a char\(n\) result; vector source operands of type short\(n\) and ushort\(n\) return a short\(n\) result; vector source operands of type int\(n\), uint\(n\) and float\(n\) return an int\(n\) result; vector source operands of type long\(n\), ulong\(n\) and double\(n\) return a long\(n\) result.

For scalar types, the result of the logical unary operator is 0 if the value of its operand compares unequal to 0, and 1 if the value of its operand compares equal to 0. For vector types, the unary operator shall return a 0 if the value of its operand compares unequal to 0, and -1 (i.e. all bits set) if the value of its operand compares equal to 0.

i. The ternary selection operator (?:) operates on three expressions (exp\(1\) ? exp\(2\) : exp\(3\)). This operator evaluates the first expression exp\(1\), which can be a scalar or vector result except float. If the result is a scalar value then it selects to evaluate the second expression if the result compares unequal to 0, otherwise it selects to evaluate the third expression. If the result is a vector value, then this is equivalent to calling select(exp\(3\), exp\(2\), exp\(1\)). The select function is described in table 6.14. The second and third expressions can be any type, as long their types match, or there is a conversion in section 6.2.1 Implicit Conversions that can be applied to one of the expressions to make their types match, or one is a vector and the other is a scalar and the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand and widened to the same type as the vector type. This resulting matching type is the type of the entire expression.

j. The operators right-shift (\(\gg\)), left-shift (\(\ll\)) operate on all scalar and vector built-in types except the built-in scalar and vector float types. For built-in vector types, the operators are applied component-wise. For the right-shift (\(\gg\)), left-shift (\(\ll\)) operators, the rightmost operand must be a scalar if the first operand is a scalar, and the rightmost operand can be a vector or scalar if the first operand is a vector.

The result of \(E1 \ll E2\) is \(E1\) left-shifted by \(\log_2(N)\) least significant bits in \(E2\) viewed as an unsigned integer value, where \(N\) is the number of bits used to represent the data type of \(E1\) after integer promotion\(^{18}\), if \(E1\) is a scalar, or the number of bits used to represent the

\(^{18}\) Integer promotion is described in ISO/IEC 9899:1999 in section 6.3.1.1.
type of $E_1$ elements, if $E_1$ is a vector. The vacated bits are filled with zeros.

The result of $E_1 >> E_2$ is $E_1$ right-shifted by $\log_2(N)$ least significant bits in $E_2$ viewed as an unsigned integer value, where $N$ is the number of bits used to represent the data type of $E_1$ after integer promotion, if $E_1$ is a scalar, or the number of bits used to represent the type of $E_1$ elements, if $E_1$ is a vector. If $E_1$ has an unsigned type or if $E_1$ has a signed type and a nonnegative value, the vacated bits are filled with zeros. If $E_1$ has a signed type and a negative value, the vacated bits are filled with ones.

k. The `sizeof` operator yields the size (in bytes) of its operand, including any padding bytes (refer to section 6.1.5) needed for alignment, which may be an expression or the parenthesized name of a type. The size is determined from the type of the operand. The result is of type `size_t`. If the type of the operand is a variable length array\(^{19}\) type, the operand is evaluated; otherwise, the operand is not evaluated and the result is an integer constant.

When applied to an operand that has type `char`, `uchar`, the result is 1. When applied to an operand that has type `short`, `ushort`, or `half` the result is 2. When applied to an operand that has type `int`, `uint` or `float`, the result is 4. When applied to an operand that has type `long`, `ulong` or `double`, the result is 8. When applied to an operand that is a vector type, the result\(^{20}\) is number of components \(\times\) size of each scalar component. When applied to an operand that has array type, the result is the total number of bytes in the array. When applied to an operand that has structure or union type, the result is the total number of bytes in such an object, including internal and trailing padding. The `sizeof` operator shall not be applied to an expression that has function type or an incomplete type, to the parenthesized name of such a type, or to an expression that designates a bit-field struct member\(^{21}\).

The behavior of applying the `sizeof` operator to the `bool`, `image2d_t`, `image3d_t`, `image2d_array_t`, `image2d_depth_t`, `image2d_array_depth_t`, `image1d_t`, `image1d_buffer_t` or `image1d_array_t`, `sampler_t`, `clk_event_t`, `queue_t` and `event_t` types is implementation-defined.

l. The comma (\(,\)) operator operates on expressions by returning the type and value of the right-most expression in a comma separated list of expressions. All expressions are evaluated, in order, from left to right.

m. The unary (*\(\)) operator denotes indirection. If the operand points to an object, the result is an lvalue designating the object. If the operand has type "pointer to type", the result has type "type". If an invalid value has been assigned to the pointer, the behavior of the unary *

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\(^{19}\) Variable length arrays are not supported in OpenCL 1.1. Refer to section 6.9.d.

\(^{20}\) Except for 3-component vectors whose size is defined as $4 \times$ size of each scalar component.

\(^{21}\) Bit-field struct members are not supported in OpenCL 1.1. Refer to section 6.9.c.
The unary (\&) operator returns the address of its operand. If the operand has type "‘type’", the result has type "‘pointer to type’". If the operand is the result of a unary * operator, neither that operator nor the \& operator is evaluated and the result is as if both were omitted, except that the constraints on the operators still apply and the result is not an lvalue. Similarly, if the operand is the result of a [] operator, neither the \& operator nor the unary * that is implied by the [] is evaluated and the result is as if the \& operator were removed and the [] operator were changed to a + operator. Otherwise, the result is a pointer to the object designated by its operand\(^{23}\).

Assignments of values to variable names are done with the assignment operator ( = ), like

\[ \text{lvalue} = \text{expression} \]

The assignment operator stores the value of expression into lvalue. The expression and lvalue must have the same type, or the expression must have a type in table 6.1, in which case an implicit conversion will be done on the expression before the assignment is done.

If expression is a scalar type and lvalue is a vector type, the scalar is converted to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.

Any other desired type-conversions must be specified explicitly. L-values must be writable. Variables that are built-in types, entire structures or arrays, structure fields, l-values with the field selector ( . ) applied to select components or swizzles without repeated fields, l-values within parentheses, and l-values dereferenced with the array subscript operator ( [ ] ) are all l-values. Other binary or unary expressions, function names, swizzles with repeated fields, and constants cannot be l-values. The ternary operator (?:) is also not allowed as an l-value.

The order of evaluation of the operands is unspecified. If an attempt is made to modify the result of an assignment operator or to access it after the next sequence point, the behavior is undefined. Other assignment operators are the assignments add into (+=), subtract from (=-), multiply into (*=), divide into (/=), modulus into (%=), left shift by (<=), right shift by (>>=), and into (&=), inclusive or into ( |=), and exclusive or into ( ^=).

The expression

\[ \text{lvalue op= expression} \]

\(^{22}\) Among the invalid values for dereferencing a pointer by the unary * operator are a null pointer, an address inappropriately aligned for the type of object pointed to, and the address of an object after the end of its lifetime. If \*P is an lvalue and \( T \) is the name of an object pointer type, \((T)P\) is an lvalue that has a type compatible with that to which \( T \) points.

\(^{23}\) Thus, &*E is equivalent to E (even if E is a null pointer), and &*(E1[E2]) to ((E1)+(E2)). It is always true that if E is an lvalue that is a valid operand of the unary \& operator, \&E is an lvalue equal to E.
is equivalent to

\[ \text{lvalue} = \text{lvalue} \ \text{op} \ \text{expression} \]

and the l-value and expression must satisfy the requirements for both operator \text{op} and assignment (=).

**Note:** Except for the **sizeof** operator, the **half** data type cannot be used with any of the operators described in this section.

### 6.4 Vector Operations

Vector operations are component-wise. Usually, when an operator operates on a vector, it is operating independently on each component of the vector, in a component-wise fashion.

For example,

```c
float4 v, u;
float f;

v = u + f;
```

will be equivalent to

```c
v.x = u.x + f;
v.y = u.y + f;
v.z = u.z + f;
v.w = u.w + f;
```

And

```c
float4 v, u, w;

w = v + u;
```

will be equivalent to

```c
w.x = v.x + u.x;
w.y = v.y + u.y;
w.z = v.z + u.z;
w.w = v.w + u.w;
```

and likewise for most operators and all integer and floating-point vector types.
6.5 Address Space Qualifiers

OpenCL implements the following disjoint named address spaces: __global, __local, __constant and __private. The address space qualifier may be used in variable declarations to specify the region of memory that is used to allocate the object. The C syntax for type qualifiers is extended in OpenCL to include an address space name as a valid type qualifier. If the type of an object is qualified by an address space name, the object is allocated in the specified address name; otherwise, the object is allocated in the generic address space.

The address space names without the __ prefix i.e. global, local, constant and private may be substituted for the corresponding address space names with the __prefix.

The address space name for arguments to a function in a program, or local variables of a function is __private. All function arguments shall be in the __private address space. The address space for a variable at program scope or a static variable inside a function can either be __global or __constant, but defaults to __global if not specified.

Examples:

```
// declares a pointer p in the private address space that
// points to an object in address space global
global int *p;

void foo (...) {
    // declares an array of 4 floats in the private address space
    float x[4];
    ...
}
```

OpenCL 2.0 adds support for an unnamed generic address space. Pointers that are declared without pointing to a named address space point to the generic address space. Before referring to the region pointed to, the pointer must be associated with a named address space. Functions may be written with arguments and return values that point to the generic address space.

kernel function arguments declared to be a pointer or an array of a type must point to one of the named address spaces __global, __local or __constant.

The named address spaces are a subset of the generic address space except for the constant address space.

A pointer to address space A can only be assigned to a pointer to the same address space A or a pointer to the generic address space. Casting a pointer to address space A to a pointer to address space B is illegal if A and B are named address spaces and A is not the same as B.
Examples:

```c
private int f() { ... } // should generate an error
local int *f() { ... } // allowed
local int * private f() { ... }; // should generate an error.
```

The `__global`, `__constant`, `__local`, `__private`, `__generic`, `global`, `constant`, `local`, `private` and `generic` names are reserved for use as address space qualifiers and shall not be used otherwise.

NOTE: The size of pointers to different address spaces may differ. It is not correct to assume that, for example, `sizeof(__global int *)` always equals `sizeof(__local int *)`.

### 6.5.1 `__global` (or global)

The `__global` or `global` address space name is used to refer to memory objects (buffer or image objects) allocated from the global memory pool.

A buffer memory object can be declared as a pointer to a scalar, vector or user-defined struct. This allows the kernel to read and/or write any location in the buffer.

The actual size of the array memory object is determined when the memory object is allocated via appropriate API calls in the host code.

Some examples are:

```c
global float4 *color; // An array of float4 elements
typedef struct {
    float a[3];
    int b[2];
} foo_t;

global foo_t *my_info; // An array of foo_t elements.
```

As image objects are always allocated from the global address space, the `__global` or `global` qualifier should not be specified for image types. The elements of an image object cannot be directly accessed. Built-in functions to read from and write to an image object are provided.

Variables defined at program scope and `static` variables inside a function can also be declared in the global address space. They can be defined with any valid OpenCL C data type except for those in `table 6.3`. In particular, such program scope variables may be of any user-defined type, or a pointer to a user-defined type. In the presence of shared virtual memory, these pointers or pointer members should work as expected as long as they are shared virtual memory pointers.
and the referenced storage has been mapped appropriately. These variables in the `global` address space have the same lifetime as the program, and their values persist between calls to any of the kernels in the program. These variables are not shared across devices. They have distinct storage.

Program scope and static variables in the `global` address space may be initialized, but only with constant expressions.

Examples:

```cpp
global int foo;        // OK.
int foo;               // OK. Declared in the global address space
global uchar buf[512]; // OK.
global int baz = 12;   // OK. Initialization is allowed
static global int bat; // OK. Internal linkage

static int foo;        // OK. Declared in the global address space
static global int foo; // OK.

int *foo;              // OK. foo is allocated in global address space.
                        // pointer to foo in generic address space
void func(...)         
{                       
  int *foo;             // OK. foo is allocated in private address space.
  // foo points to a location in generic address space.
  ...
}

global int * global ptr;  // OK.
int * global ptr;        // OK.
constant int *global ptr=&baz;  // error since baz is in global address
                                 // space.
global int * constant ptr = &baz; // OK

// Pointers work. Also, initialization to a constant known at
// program load time
global int *global baz_ptr = &baz;

global image2d_t im;     // Error. Invalid type for program scope
                         // variables

global event_t ev;       // Error. Invalid type for program scope variables

global int *bad_ptr;     // Error. No implicit address space
```

The `const` qualifier can also be used with the `__global` qualifier to specify a read-only buffer memory object.
6.5.2 __local (or local)

The __local or local address space name is used to describe variables that need to be allocated in local memory and are shared by all work-items of a work-group. Pointers to the __local address space are allowed as arguments to functions (including kernel functions). Variables declared in the __local address space inside a kernel function must occur at kernel function scope.

Some examples of variables allocated in the __local address space inside a kernel function are:

```c
kernel void my_func(...) {
    local float a;  // A single float allocated
    // in local address space

    local float b[10];  // An array of 10 floats
    // allocated in local address space.

    if (...) {
        // example of variable in __local address space but not
        // declared at __kernel function scope.
        local float c;  // not allowed.
    }
}
```

Variables allocated in the __local address space inside a kernel function cannot be initialized.

```c
kernel void my_func(...) {
    local float a = 1;  // not allowed

    local float b;
    b = 1;  // allowed
}
```

NOTE: Variables allocated in the __local address space inside a kernel function are allocated for each work-group executing the kernel and exist only for the lifetime of the work-group executing the kernel.

6.5.3 __constant (or constant)

The __constant or constant address space name is used to describe variables allocated in global memory and which are accessed inside a kernel(s) as read-only variables. These read-only variables can be accessed by all (global) work-items of the kernel during its execution.
Pointers to the \texttt{__constant} address space are allowed as arguments to functions (including kernel functions) and for variables declared inside functions.

All string literal storage shall be in the \texttt{__constant} address space.

NOTE: Each argument to a kernel that is a pointer to the \texttt{__constant} address space is counted separately towards the maximum number of such arguments, defined as CL\_DEVICE\_MAX\_CONSTANT\_ARGS in Table 4.3.

Variables in the program scope can be declared in the \texttt{__constant} address space. Variables in the outermost scope of kernel functions can be declared in the \texttt{__constant} address space. These variables are required to be initialized and the values used to initialize these variables must be a compile time constant. Writing to such a variable results in a compile-time error.

Implementations are not required to aggregate these declarations into the fewest number of constant arguments. This behavior is implementation defined.

Thus portable code must conservatively assume that each variable declared inside a function or in program scope allocated in the \texttt{__constant} address space counts as a separate constant argument.

6.5.4 \texttt{__private} (or private)

Variables inside a kernel function not declared with an address space qualifier, all variables inside non-kernel functions, and all function arguments are in the \texttt{__private} or \texttt{private} address space. Variables declared as pointers are considered to point to the \texttt{__private} address space if an address space qualifier is not specified.

6.5.5 The generic address space

The following rules apply when using pointers that point to the generic address space:

\begin{itemize}
  \item A pointer that points to the \texttt{global}, \texttt{local} or \texttt{private} address space can be implicitly converted to a pointer to the unnamed generic address space but not vice-versa.
  \item Pointer casts can be used to cast a pointer that points to the \texttt{global}, \texttt{local} or \texttt{private} space to the unnamed generic address space and vice-versa.
  \item A pointer that points to the \texttt{constant} address space cannot be cast or implicitly converted to the generic address space.
\end{itemize}

A few examples follow.
This is the canonical example. In this example, function `foo` is declared with an argument that is a pointer with no address space qualifier.

```c
void foo(int *a)
{
    *a = *a + 2;
}

kernel void k1(local int *a)
{
    ...
    foo(a);
    ...
}

kernel void k2(global int *a)
{
    ...
    foo(a);
    ...
}
```

In the example below, `var` is in the unnamed generic address space which gets mapped to the global or local address space depending on the result of the conditional expression.

```c
kernel void bar(global int *g, local int *l)
{
    int *var;
    if (is_even(get_global_id(0))
        var = g;
    else
        var = l;
    *var = 42;
    ...
}
```

The example below is an example with one unnamed generic address space pointer with multiple named address space assignments.

```c
int *ptr;
global int g;
ptr = &g; // legal

local int l;
ptr = &l; // legal
```
private int p;
ptr = &p; // legal

constant int c;
ptr = &c; // illegal

The example below is an example with one unnamed generic address space pointer being assigned to point to several named address spaces.

global int * gp;
local int * lp;
private int * pp;

int * p;
p = gp; // legal
p = lp; // legal
p = pp; // legal

// it is illegal to convert from a generic pointer
// to an explicit address space pointer without a cast:
gp = p; // compile-time error
lp = p; // compile-time error
pp = p; // compile-time error

6.5.6 Changes to ISO/IEC 9899:1999

This section details the modifications to ISO/IEC 9899:1999 needed to incorporate the functionality of named address space and the generic address space:

Clause 6.2.5 – Types, replace paragraph 26 with the following paragraphs:

If type $T$ is qualified by the address space qualifier for address space $A$, then “$T$ is in $A$”. If type $T$ is in address space $A$, a pointer to $T$ is also a “pointer into $A$” and the referenced address space of the pointer is $A$.

A pointer to $\text{void}$ in any address space shall have the same representation and alignment requirements as a pointer to a character type in the same address space. Similarly, pointers to differently access-qualified versions of compatible types shall have the same representation and alignment requirements. All pointers to structure types in the same address space shall have the same representation and alignment requirements as each other. All pointers to union types in the same address space shall have the same representation and alignment requirements as each other.

Clause 6.3.2.3 – Pointers, replace the first two paragraphs with the following paragraphs:

If a pointer into one address space is converted to a pointer into another address space, then
unless the original pointer is a null pointer or the location referred to by the original pointer is within the second address space, the behavior is undefined. (For the original pointer to refer to a location within the second address space, the two address spaces must overlap).

A pointer to `void` in any address space may be converted to or from a pointer to any incomplete or object type. A pointer to any incomplete or object type in some address space may be converted to a pointer to `void` in an enclosing address space and back again; the result shall compare equal to the original pointer.

For any qualifier q, a pointer to a non-q-qualified type may be converted to a pointer to the q-qualified version of the type (but with the same address-space qualifier or the generic address space); the values stored in the original and converted pointers shall compare equal.

**Clause 6.3.2.3 – Pointers, replace the last sentence of paragraph 4 with:**

Conversion of a null pointer to another pointer type yields a null pointer of that type. Any two null pointers whose referenced address spaces overlap shall compare equal.

**Clause 6.5.2.2 – Function calls, change the second bullet of paragraph 6 to:**

both types are pointers to qualified or unqualified versions of a character type or `void` in the same address space or one type is a pointer in a named address space and the other is a pointer in the generic address space.

**Clause 6.5.6 – Additive operators, add another constraint paragraph:**

For subtraction, if the two operands are pointers into different address spaces, the address spaces must overlap.

**Clause 6.5.8 – Relational operators, add another constraint paragraph:**

If the two operands are pointers into different address spaces, the address spaces must overlap.

**Clause 6.5.8 – Relational operators, add a new paragraph between existing paragraphs 3 and 4:**

If the two operands are pointers into different address spaces, one of the address spaces encloses the other. The pointer into the enclosed address space is first converted to a pointer to the same reference type except with any address-space qualifier removed and any address-space qualifier of the other pointer’s reference type added. (After this conversion, both pointers are pointers into the same address space).

Examples:

```c
kernel void test1()
{
```
global int arr[5] = { 0, 1, 2, 3, 4 };
int *p = &arr[1];
global int *q = &arr[3];

// q implicitly converted to the generic address space
// since the generic address space encloses the global
// address space
if (q >= p)
    printf("true\n");

// q implicitly converted to the generic address space
// since the generic address space encloses the global
// address space
if (p <= q)
    printf("true\n");
}

Clause 6.5.9 – Equality operators, add another constraint paragraph:
If the two operands are pointers into different address spaces, the address spaces must overlap.

Clause 6.5.9 – Equality operators, replace paragraph 5 with:
Otherwise, at least one operand is a pointer. If one operand is a pointer and the other is a null pointer constant, the null pointer constant is converted to the type of the pointer. If both operands are pointers, each of the following conversions is performed as applicable:

- If the two operands are pointers into different address spaces, one of the address spaces encloses the other. The pointer into the enclosed address space is first converted to a pointer to the same reference type except with any address-space qualifier removed and any address-space qualifier of the other pointer’s reference type added. (After this conversion, both pointers are pointers into the same address space).

- Then, if one operand is a pointer to an object or incomplete type and the other is a pointer to a qualified or unqualified version of `void`, the former is converted to the type of the latter.

Examples:

```c
int *ptr = NULL;
local int lval = SOME_VAL;
local int *lptr = &lval;
global int gval = SOME_OTHER_VAL;
global int *gptr = &gval;

ptr = lptr;
```
if (ptr == gptr) // legal
{  
  ...
}

if (ptr == lptr) // legal
{  
  ...
}

if (lptr == gptr) // illegal, compiler error
{  
  ...
}

Consider the following example:

```c
bool callee(int *p1, int *p2)
{
    if (p1 == p2)
        return true;
    return false;
}
```

```c
void caller()
{
    global int *gptr = 0xdeadbeef;
    private int *pptr = 0xdeadbeef;

    // behavior of callee is undefined
    bool b = callee(gptr, pptr);
}
```

The behavior of `callee` is undefined as `gptr` and `pptr` are in different address spaces. The example above would have the same undefined behavior if the equality operator is replaced with a relational operator.

Examples:

```c
int *ptr = NULL;
local int *lptr = NULL;
global int *gptr = NULL;

if (ptr == NULL)  // legal
{
  ...
}
```
if (ptr == lptr) // legal
{
    ...
}

if (lptr == gptr) // compile-time error
{
    ...
}

ptr = lptr; // legal

intptr l = (intptr_t)lptr;
if (l == 0) // legal
{
    ...
}

if (l == NULL) // legal
{
    ...
}

Clause 6.5.9 – Equality operators, replace first sentence of paragraph 6 with:

Two pointers compare equal if and only if both are null pointers with overlapping address spaces.

Clause 6.5.15 – Conditional operator, add another constraint paragraph:

If the second and third operands are pointers into different address spaces, the address spaces must overlap.

Examples:

kernel void test1()
{
    global int arr[5] = { 0, 1, 2, 3, 4 };
    int *p = &arr[1];
    global int *q = &arr[3];
    local int *r = NULL;
    int *val = NULL;

    // legal. 2\textsuperscript{nd} and 3\textsuperscript{rd} operands are in address spaces
// that overlap
val = (q >= p) ? q : p;

// compiler error.  2\textsuperscript{nd} and 3\textsuperscript{rd} operands are in disjoint address spaces
val = (q >= p) ? q : r;
}

\textbf{Clause 6.5.16.1 – Simple assignment, change the third and fourth bullets of paragraph 1 to:}

- both operands are pointers to qualified or unqualified versions of compatible types, the referenced address space of the left encloses the referenced address space of the right, and the type pointed to by the left has all the qualifiers of the type pointed to by the right.

- one operand is a pointer to an object or incomplete type and the other is a pointer to a qualified or unqualified version of \texttt{void}, the referenced address space of the left encloses the referenced address space of the right, and the type pointed to by the left has all the qualifiers of the type pointed to by the right.

Examples:

\begin{verbatim}
kernel void f()
{
    int *ptr;
    local int *lptr;
    global int *gptr;
    local int val = 55;

    ptr = &val; // legal: implicit cast to generic, then assign
    lptr = ptr; // illegal: no implicit cast from
                // generic to local
    lptr = gptr; // illegal: no implicit cast from
                 // global to local
    ptr = gptr; // legal: implicit cast from global to generic,
                 // then assign
}
\end{verbatim}

\textbf{Clause 6.7.2.1 – Structure and union specifiers, add a new constraint paragraph:}

Within a structure or union specifier, the type of a member shall not be qualified by an address space qualifier.

\textbf{Clause 6.7.3 – Type qualifiers, add three new constraint paragraphs:}

No type shall be qualified by qualifiers for two or more different address spaces.
6.6 Access Qualifiers

Image objects specified as arguments to a kernel can be declared to be read-only, write-only or read-write. A kernel cannot read from and write to the same image object. The \texttt{__read_only} (or \texttt{read_only}) and \texttt{__write_only} (or \texttt{write_only}) qualifiers must be used with image object arguments to declare if the image object is being read or written by a kernel. The \texttt{__read_write} (or \texttt{read_write}) qualifier must be used with image object arguments to declare if the image object is being both read and written by a kernel. The default qualifier is \texttt{read_only}.

In the following example

```c
kernel void foo (read_only image2d_t imageA,
    write_only image2d_t imageB)
{
    ...
}
```

\texttt{imageA} is a read-only 2D image object, and \texttt{imageB} is a write-only 2D image object.

The sampler-less read image and write image built-ins can be used with image declared with the \texttt{__read_write} (or \texttt{read_write}) qualifier. Calls to built-ins that read from an image using a sampler for images declared with the \texttt{__read_write} (or \texttt{read_write}) qualifier will be a compilation error.

Pipe objects specified as arguments to a kernel can also use these access qualifiers. Refer to section 6.13.16 for a detailed description on how these access qualifiers can be used with pipes.

The \texttt{__read_only}, \texttt{__write_only}, \texttt{__read_write}, \texttt{read_only}, \texttt{write_only} and \texttt{read_write} names are reserved for use as access qualifiers and shall not be used otherwise.
6.7 Function Qualifiers

6.7.1 __kernel (or kernel)

The __kernel (or kernel) qualifier declares a function to be a kernel that can be executed by an application on an OpenCL device(s). The following rules apply to functions that are declared with this qualifier:

- It can be executed on the device only
- It can be called by the host
- It is just a regular function call if a __kernel function is called by another kernel function.

NOTE:

Kernel functions with variables declared inside the function with the __local or local qualifier can be called by the host using appropriate APIs such as clEnqueueNDRangeKernel.

The __kernel and kernel names are reserved for use as functions qualifiers and shall not be used otherwise.

6.7.2 Optional Attribute Qualifiers

The __kernel qualifier can be used with the keyword __attribute__ to declare additional information about the kernel function as described below.

The optional __attribute__((vec_type_hint(<type>)))\(^{24}\) is a hint to the compiler and is intended to be a representation of the computational width of the __kernel, and should serve as the basis for calculating processor bandwidth utilization when the compiler is looking to autovectorize the code. In the __attribute__((vec_type_hint(<type>))) qualifier <type> is one of the built-in vector types listed in table 6.2 or the constituent scalar element types. If vec_type_hint (<type>) is not specified, the kernel is assumed to have the __attribute__((vec_type_hint(int))) qualifier.

For example, where the developer specified a width of float4, the compiler should assume that the computation usually uses up to 4 lanes of a float vector, and would decide to merge work-items or possibly even separate one work-item into many threads to better match the hardware capabilities. A conforming implementation is not required to autovectorize code, but

---

\(^{24}\) Implicit in autovectorization is the assumption that any libraries called from the __kernel must be recompilable at run time to handle cases where the compiler decides to merge or separate workitems. This probably means that such libraries can never be hard coded binaries or that hard coded binaries must be accompanied either by source or some retargetable intermediate representation. This may be a code security question for some.
shall support the hint. A compiler may autovectorize, even if no hint is provided. If an implementation merges \( N \) work-items into one thread, it is responsible for correctly handling cases where the number of global or local work-items in any dimension modulo \( N \) is not zero.

Examples:

```c
// autovectorize assuming float4 as the
// basic computation width
__kernel __attribute__((vec_type_hint(float4)))
void foo(__global float4 *p ) { .... }

// autovectorize assuming double as the
// basic computation width
__kernel __attribute__((vec_type_hint(double)))
void foo(__global float4 *p ) { .... }

// autovectorize assuming int (default)
// as the basic computation width
__kernel
void foo(__global float4 *p ) { .... }
```

If for example, a `__kernel` function is declared with `__attribute__((vec_type_hint(float4)))` (meaning that most operations in the `__kernel` function are explicitly vectorized using `float4`) and the kernel is running using Intel® Advanced Vector Instructions (Intel® AVX) which implements a 8-float-wide vector unit, the autovectorizer might choose to merge two work-items to one thread, running a second work-item in the high half of the 256-bit AVX register.

As another example, a Power4 machine has two scalar double precision floating-point units with an 6-cycle deep pipe. An autovectorizer for the Power4 machine might choose to interleave six kernels declared with the `__attribute__((vec_type_hint(double2)))` qualifier into one hardware thread, to ensure that there is always 12-way parallelism available to saturate the FPUs. It might also choose to merge 4 or 8 work-items (or some other number) if it concludes that these are better choices, due to resource utilization concerns or some preference for divisibility by 2.

The optional `__attribute__((work_group_size_hint(X, Y, Z)))` is a hint to the compiler and is intended to specify the work-group size that may be used i.e. value most likely to be specified by the `local_work_size` argument to `clEnqueueNDRangeKernel`. For example the `__attribute__((work_group_size_hint(1, 1, 1)))` is a hint to the compiler that the kernel will most likely be executed with a work-group size of 1.

The optional `__attribute__((reqd_work_group_size(X, Y, Z)))` is the work-group size that must be used as the `local_work_size` argument to `clEnqueueNDRangeKernel`. This allows the compiler to optimize the generated code appropriately for this kernel.
If \( Z \) is one, the \textit{work\_dim} argument to \texttt{clEnqueueNDRangeKernel} can be 2 or 3. If \( Y \) and \( Z \) are one, the \textit{work\_dim} argument to \texttt{clEnqueueNDRangeKernel} can be 1, 2 or 3.

The optional \texttt{__attribute__((nosvm))} qualifier can be used with a pointer variable to inform the compiler that the pointer does not refer to a shared virtual memory region.
6.8 Storage-Class Specifiers

The typedef, extern and static storage-class specifiers are supported. The auto and register storage-class specifiers are not supported.

The extern storage-class specifier can only be used for functions (kernel and non-kernel functions) and global variables declared in program scope or variables declared inside a function (kernel and non-kernel functions). The static storage-class specifier can only be used for non-kernel functions, global variables declared in program scope and variables inside a function declared in the global or constant address space.

Examples:

extern constant float4 noise_table[256];
static constant float4 color_table[256];

extern kernel void my_foo(image2d_t img);
extern void my_bar(global float *a);

kernel void my_func(image2d_t img, global float *a)
{
extern constant float4 a;
static constant float4 b = (float4)(1.0f); // OK.
static float c; // Error: No implicit address space
global int hurl; // Error: Must be static
...
my_foo(img);
...
my_bar(a);
...
while (1)
{
    static global int inside; // OK.
    ...
}
...
}
6.9 Restrictions\textsuperscript{25}

a. The use of pointers is somewhat restricted. The following rules apply:

- Arguments to kernel functions declared in a program that are pointers must be declared with the \texttt{__global}, \texttt{__constant} or \texttt{__local} qualifier.

- A pointer declared with the \texttt{__constant} qualifier can only be assigned to a pointer declared with the \texttt{__constant} qualifier respectively.

- Pointers to functions are not allowed.

- Arguments to kernel functions in a program cannot be declared as a pointer to a pointer(s). Variables inside a function or arguments to non-kernel functions in a program can be declared as a pointer to a pointer(s).

b. An image type (\texttt{image2d_t}, \texttt{image3d_t}, \texttt{image2d_array_t}, \texttt{image1d_t}, \texttt{image1d_buffer_t} or \texttt{image1d_array_t}) can only be used as the type of a function argument. An image function argument cannot be modified. Elements of an image can only be accessed using built-in functions described in section 6.13.14.

An image type cannot be used to declare a variable, a structure or union field, an array of images, a pointer to an image, or the return type of a function. An image type cannot be used with the \texttt{__global}, \texttt{__private}, \texttt{__local} and \texttt{__constant} address space qualifiers. The \texttt{image3d_t} type cannot be used with the \texttt{__write_only} access qualifier unless the \texttt{cl_khr_3d_image_writes} extension is enabled. An image type cannot be used with the \texttt{__read_write} access qualifier which is reserved for future use.

The sampler type (\texttt{sampler_t}) can only be used as the type of a function argument or a variable declared in the program scope or the outermost scope of a kernel function. The behavior of a sampler variable declared in a non-outermost scope of a kernel function is implementation-defined. A sampler argument or variable cannot be modified.

The sampler type cannot be used to declare a structure or union field, an array of samplers, a pointer to a sampler, or the return type of a function. The sampler type cannot be used with the \texttt{__local} and \texttt{__global} address space qualifiers.

c. Bit-field struct members are currently not supported.

d. Variable length arrays and structures with flexible (or unsized) arrays are not supported.

\textsuperscript{25} Items struckthrough are restrictions in OpenCL 1.0 that are removed in OpenCL 1.1.
e. Variadic macros and functions with the exception of printf and enqueue_kernel are not supported.

f. The library functions defined in the C99 standard headers assert.h, ctype.h, complex.h, errno.h, fenv.h, float.h, inttypes.h, limits.h, locale.h, setjmp.h, signal.h, stdarg.h, stdio.h, stdlib.h, string.h, tgmath.h, time.h, wchar.h and wctype.h are not available and cannot be included by a program.

g. The auto and register storage-class specifiers are not supported.

h. Predefined identifiers are not supported.

i. Recursion is not supported.

j. The return type of a kernel function must be void.

k. Arguments to kernel functions in a program cannot be declared with the built-in scalar types bool, half, size_t, ptdiff_t, intptr_t, and uintptr_t or a struct and/or union that contain fields declared to be one of these built-in scalar types. The size in bytes of these types except half are implementation-defined and in addition can also be different for the OpenCL device and the host processor making it difficult to allocate buffer objects to be passed as arguments to a kernel declared as pointer to these types. half is not supported as half can be used as a storage format only and is not a data type on which floating-point arithmetic can be performed.

l. Whether or not irreducible control flow is illegal is implementation defined.

m. Built-in types that are less than 32 bits in size i.e. char, uchar, char2, uchar2, short, ushort, and half have the following restriction:

   Writes to a pointer (or arrays) of type char, uchar, char2, uchar2, short, ushort, and half or to elements of a struct that are of type char, uchar, char2, uchar2, short and ushort are not supported. Refer to section 9.9 for additional information.

The kernel example below shows what memory operations are not supported on built-in types less than 32-bits in size.

```c
kernel void
do_proc (__global char *pA, short b, __global short *pB)
{
    char x[100];
```

26 Unless the cl_khr_fp16 extension is supported.
private char *px = x;

int id = (int)get_global_id(0);

short f;

f = pB[id] + b; ← is allowed


pB[id] = b; ← error. pB cannot be written

n. The type qualifiers const, restrict and volatile as defined by the C99 specification are supported. These qualifiers cannot be used with image2d_t, image3d_t, image2d_array_t, image2d_depth_t, image2d_array_depth_t, image1d_t, image1d_buffer_t and image1d_array_t types. Types other than pointer types shall not use the restrict qualifier.

o. The event type (event_t) cannot be used as the type of a kernel function argument. The event type cannot be used to declare a program scope variable. The event type cannot be used to declare a structure or union field. The event type cannot be used with the _local, _constant and _global address space qualifiers.

p. The clk_event_t, ndrange_t and reserve_id_t types cannot be used as arguments to kernel functions that get enqueued from the host. The clk_event_t and reserve_id_t types cannot be declared in program scope.

q. The values returned by applying the sizeof operator to the queue_t, clk_event_t, ndrange_t and reserve_id_t types is implementation-defined.

r. Kernels enqueued by the host must continue to have their arguments that are a pointer to a type declared to point to a named address space.

s. A function in an OpenCL program cannot be called main.
6.10 Preprocessor Directives and Macros

The preprocessing directives defined by the C99 specification are supported.

The `#pragma` directive is described as:

```
# pragma pp-tokens opt new-line
```

A `#pragma` directive where the preprocessing token `OPENCL` (used instead of `STDC`) does not immediately follow `pragma` in the directive (prior to any macro replacement) causes the implementation to behave in an implementation-defined manner. The behavior might cause translation to fail or cause the translator or the resulting program to behave in a non-conforming manner. Any such `pragma` that is not recognized by the implementation is ignored. If the preprocessing token `OPENCL` does immediately follow `pragma` in the directive (prior to any macro replacement), then no macro replacement is performed on the directive, and the directive shall have one of the following forms whose meanings are described elsewhere:

```
#pragma OPENCL FP_CONTRACT on-off-switch
    on-off-switch: one of ON OFF DEFAULT

#pragma OPENCL EXTENSION extensionname : behavior
#pragma OPENCL EXTENSION all : behavior
```

The following predefined macro names are available.

```
__FILE__  The presumed name of the current source file (a character string literal).
__LINE__  The presumed line number (within the current source file) of the current source line (an integer constant).
__OPENCL_VERSION__ substitutes an integer reflecting the version number of the OpenCL supported by the OpenCL device. The version of OpenCL described in this document will have __OPENCL_VERSION__ substitute the integer 200.
CL_VERSION_1_0 substitutes the integer 100 reflecting the OpenCL 1.0 version.
CL_VERSION_1_1 substitutes the integer 110 reflecting the OpenCL 1.1 version.
CL_VERSION_1_2 substitutes the integer 120 reflecting the OpenCL 1.2 version.
CL_VERSION_2_0 substitutes the integer 200 reflecting the OpenCL 2.0 version.
__OPENCL_C_VERSION__ substitutes an integer reflecting the OpenCL C version
```
specified by the –cl-std build option (specified in section 5.8.4.5) to clBuildProgram or clCompileProgram. If the –cl-std build option is not specified, the highest OpenCL C 1.x language version supported by each device is used as the version of OpenCL C when compiling the program for each device. The version of OpenCL C described in this document will have __OPENCL_C_VERSION__ substitute the integer 200 if –cl-std=CL2.0 is specified.

__ENDIAN_LITTLE__ is used to determine if the OpenCL device is a little endian architecture or a big endian architecture (an integer constant of 1 if device is little endian and is undefined otherwise). Also refer to CL_DEVICE_ENDIAN_LITTLE specified in table 4.3.

__kernel_exec(X, type\_n) (and kernel\_exec(X, type\_n)) is defined as

__kernel __attribute__((work\_group\_size\_hint(X, 1, 1))) \
__attribute__((vec\_type\_hint(type\_n)))

__IMAGE_SUPPORT__ is used to determine if the OpenCL device supports images. This is an integer constant of 1 if images are supported and is undefined otherwise. Also refer to CLDEVICE_IMAGE_SUPPORT specified in table 4.3.

__FAST_RELAXED_MATH__ is used to determine if the –cl-fast-relaxed-math optimization option is specified in build options given to clBuildProgram or clCompileProgram. This is an integer constant of 1 if the –cl-fast-relaxed-math build option is specified and is undefined otherwise.

The NULL macro expands to a null pointer constant. An integer constant expression with the value 0, or such an expression cast to type void * is called a null pointer constant.

The macro names defined by the C99 specification but not currently supported by OpenCL are reserved for future use.

The predefined identifier __func__ is available.
6.11 Attribute Qualifiers

This section describes the syntax with which __attribute__ may be used, and the constructs to which attribute specifiers bind.

An attribute specifier is of the form __attribute__ ((attribute-list)).

An attribute list is defined as:

```
attribute-list:
  attribute_opt
  attribute-list , attribute_opt

attribute:
  attribute-token attribute-argument-clause_opt

attribute-token:
  identifier

attribute-argument-clause:
  ( attribute-argument-list )

attribute-argument-list:
  attribute-argument
  attribute-argument-list, attribute-argument

attribute-argument:
  assignment-expression
```

This syntax is taken directly from GCC but unlike GCC, which allows attributes to be applied only to functions, types, and variables, OpenCL attributes can be associated with:

- types;
- functions;
- variables;
- blocks; and
- control-flow statements.

In general, the rules for how an attribute binds, for a given context, are non-trivial and the reader is pointed to GCC’s documentation and Maurer and Wong’s paper [See 16. and 17. in section 11 – References] for the details.
6.11.1 Specifying Attributes of Types

The keyword __attribute__ allows you to specify special attributes of enum, struct and union types when you define such types. This keyword is followed by an attribute specification inside double parentheses. Two attributes are currently defined for types: aligned, and packed.

You may specify type attributes in an enum, struct or union type declaration or definition, or for other types in a typedef declaration.

For an enum, struct or union type, you may specify attributes either between the enum, struct or union tag and the name of the type, or just past the closing curly brace of the definition. The former syntax is preferred.

aligned (alignment)

This attribute specifies a minimum alignment (in bytes) for variables of the specified type. For example, the declarations:

```c
struct S { short f[3]; } __attribute__((aligned (8)));
typedef int more_aligned_int __attribute__((aligned (8)));
```

force the compiler to insure (as far as it can) that each variable whose type is struct S or more_aligned_int will be allocated and aligned at least on a 8-byte boundary.

Note that the alignment of any given struct or union type is required by the ISO C standard to be at least a perfect multiple of the lowest common multiple of the alignments of all of the members of the struct or union in question and must also be a power of two. This means that you can effectively adjust the alignment of a struct or union type by attaching an aligned attribute to any one of the members of such a type, but the notation illustrated in the example above is a more obvious, intuitive, and readable way to request the compiler to adjust the alignment of an entire struct or union type.

As in the preceding example, you can explicitly specify the alignment (in bytes) that you wish the compiler to use for a given struct or union type. Alternatively, you can leave out the alignment factor and just ask the compiler to align a type to the maximum useful alignment for the target machine you are compiling for. For example, you could write:

```c
struct S { short f[3]; } __attribute__((aligned));
```

Whenever you leave out the alignment factor in an aligned attribute specification, the compiler automatically sets the alignment for the type to the largest alignment which is ever used for any data type on the target machine you are compiling for. In the example
above, the size of each short is 2 bytes, and therefore the size of the entire struct S type is 6 bytes. The smallest power of two which is greater than or equal to that is 8, so the compiler sets the alignment for the entire struct S type to 8 bytes.

Note that the effectiveness of aligned attributes may be limited by inherent limitations of the OpenCL device and compiler. For some devices, the OpenCL compiler may only be able to arrange for variables to be aligned up to a certain maximum alignment. If the OpenCL compiler is only able to align variables up to a maximum of 8 byte alignment, then specifying aligned(16) in an __attribute__ will still only provide you with 8 byte alignment. See your platform-specific documentation for further information.

The aligned attribute can only increase the alignment; but you can decrease it by specifying packed as well. See below.

packed

This attribute, attached to struct or union type definition, specifies that each member of the structure or union is placed to minimize the memory required. When attached to an enum definition, it indicates that the smallest integral type should be used.

Specifying this attribute for struct and union types is equivalent to specifying the packed attribute on each of the structure or union members.

In the following example struct my_packed_struct's members are packed closely together, but the internal layout of its s member is not packed. To do that, struct my_unpacked_struct would need to be packed, too.

```c
struct my_unpacked_struct
{
    char c;
    int i;
};

struct __attribute__((packed)) my_packed_struct
{
    char c;
    int i;
    struct my_unpacked_struct s;
};
```

You may only specify this attribute on the definition of a enum, struct or union, not on a typedef which does not also define the enumerated type, structure or union.
6.11.2 Specifying Attributes of Functions

Refer to section 6.7 for the function attribute qualifiers currently supported.

6.11.3 Specifying Attributes of Variables

The keyword __attribute__ allows you to specify special attributes of variables or structure fields. This keyword is followed by an attribute specification inside double parentheses. The following attribute qualifiers are currently defined:

aligned (alignment)

This attribute specifies a minimum alignment for the variable or structure field, measured in bytes. For example, the declaration:

```c
int x __attribute__ ((aligned (16))) = 0;
```

causes the compiler to allocate the global variable `x` on a 16-byte boundary. The alignment value specified must be a power of two.

You can also specify the alignment of structure fields. For example, to create double-word aligned `int` pair, you could write:

```c
struct foo { int x[2] __attribute__ ((aligned (8))); }
```

This is an alternative to creating a union with a `double` member that forces the union to be double-word aligned.

As in the preceding examples, you can explicitly specify the alignment (in bytes) that you wish the compiler to use for a given variable or structure field. Alternatively, you can leave out the alignment factor and just ask the compiler to align a variable or field to the maximum useful alignment for the target machine you are compiling for. For example, you could write:

```c
short array[3] __attribute__ ((aligned));
```

Whenever you leave out the alignment factor in an `aligned` attribute specification, the OpenCL compiler automatically sets the alignment for the declared variable or field to the largest alignment which is ever used for any data type on the target device you are compiling for.

When used on a `struct`, or `struct` member, the `aligned` attribute can only increase
the alignment; in order to decrease it, the \texttt{packed} attribute must be specified as well. When used as part of a \texttt{typedef}, the \texttt{aligned} attribute can both increase and decrease alignment, and specifying the \texttt{packed} attribute will generate a warning.

Note that the effectiveness of aligned attributes may be limited by inherent limitations of the OpenCL device and compiler. For some devices, the OpenCL compiler may only be able to arrange for variables to be aligned up to a certain maximum alignment. If the OpenCL compiler is only able to align variables up to a maximum of 8 byte alignment, then specifying \texttt{aligned(16)} in an \texttt{__attribute__} will still only provide you with 8 byte alignment. See your platform-specific documentation for further information.

\texttt{packed}

The \texttt{packed} attribute specifies that a variable or structure field should have the smallest possible alignment—one byte for a variable, unless you specify a larger value with the \texttt{aligned} attribute.

Here is a structure in which the field \texttt{x} is packed, so that it immediately follows \texttt{a}:

\begin{verbatim}
struct foo
{
  char a;
  int x[2] __attribute__ ((packed));
};
\end{verbatim}

An attribute list placed at the beginning of a user-defined type applies to the variable of that type and not the type, while attributes following the type body apply to the type. For example:

\begin{verbatim}
/* a has alignment of 128 */
__attribute__((aligned(128))) struct A {int i;} a;

/* b has alignment of 16 */
__attribute__((aligned(16))) struct B {double d;}
   __attribute__((aligned(32))) b ;

struct A a1; /* a1 has alignment of 4 */
struct B b1; /* b1 has alignment of 32 */
\end{verbatim}

\texttt{endian} \hspace{1cm} \texttt{(endiantype)}

The \texttt{endian} attribute determines the byte ordering of a variable. \texttt{endiantype} can be set to \texttt{host} indicating the variable uses the endianness of the host processor or can be set to \texttt{device} indicating the variable uses the endianness of the device on which the kernel will be executed. The default is \texttt{device}. 

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For example:

```c
    global float4 *p __attribute__((endian(host)));
```

specifies that data stored in memory pointed to by `p` will be in the host endian format.

The `endian` attribute can only be applied to pointer types that are in the `global` or `constant` address space. The `endian` attribute cannot be used for variables that are not a pointer type. The `endian` attribute value for both pointers must be the same when one pointer is assigned to another.

### 6.11.4 Specifying Attributes of Blocks and Control-Flow Statements

For basic blocks and control-flow statements the attribute is placed before the structure in question, for example:

```c
    __attribute__((attr1)) {...}
    for __attribute__((attr2)) (...) __attribute__((attr3)) {...}
```

Here `attr1` applies to the block in braces and `attr2` and `attr3` apply to the loop’s control construct and body, respectively.

No attribute qualifiers for blocks and control-flow statements are currently defined.

### 6.11.5 Specifying Attribute For Unrolling Loops

The `__attribute__((opencl_unroll_hint))` and `__attribute__((opencl_unroll_hint(n)))` attribute qualifiers can be used to specify that a loop (for, while and do loops) can be unrolled. This attribute qualifier can be used to specify full unrolling or partial unrolling by a specified amount. This is a compiler hint and the compiler may ignore this directive.

\( n \) is the loop unrolling factor and must be a positive integral compile time constant expression. An unroll factor of 1 disables unrolling. If \( n \) is not specified, the compiler determines the unrolling factor for the loop.

**NOTE:** The `__attribute__((opencl_unroll_hint(n)))` attribute qualifier must appear immediately before the loop to be affected.

Examples:
__attribute__((opencl_unroll_hint(2)))
while (*s != 0)
    *p++ = *s++;

The tells the compiler to unroll the above while loop by a factor of 2.

__attribute__((opencl_unroll_hint))
for (int i=0; i<2; i++)
{
    ...
}

In the example above, the compiler will determine how much to unroll the loop.

__attribute__((opencl_unroll_hint(1)))
for (int i=0; i<32; i++)
{
    ...
}

The above is an example where the loop should not be unrolled.

Below are some examples of invalid usage of
__attribute__((opencl_unroll_hint(n))).

__attribute__((opencl_unroll_hint(-1)))
while (...)
{
    ...
}

The above example is an invalid usage of the loop unroll factor as the loop unroll factor is negative.

__attribute__((opencl_unroll_hint))
if (...)
{
    ...
}

The above example is invalid because the unroll attribute qualifier is used on a non-loop construct

kernel void
my_kernel( ...)
int x;
__attribute__((opencl_unroll_hint(x)))
for (int i=0; i<x; i++)
{
    ... 
}

The above example is invalid because the loop unroll factor is not a compile-time constant expression.

### 6.11.6 Extending Attribute Qualifiers

The attribute syntax can be extended for standard language extensions and vendor specific extensions. Any extensions should follow the naming conventions outlined in the introduction to section 9 in the OpenCL 1.2 Extension Specification.

Attributes are intended as useful hints to the compiler. It is our intention that a particular implementation of OpenCL be free to ignore all attributes and the resulting executable binary will produce the same result. This does not preclude an implementation from making use of the additional information provided by attributes and performing optimizations or other transformations as it sees fit. In this case it is the programmer’s responsibility to guarantee that the information provided is in some sense correct.
6.12 Blocks

OpenCL C 2.0 adds support for the clang block syntax. Like function types, the Block type is a pair consisting of a result value type and a list of parameter types very similar to a function type. Blocks are intended to be used much like functions with the key distinction being that in addition to executable code they also contain various variable bindings to automatic (stack) or global memory.

6.12.1 Declaring and Using a Block

You use the `^` operator to declare a Block variable and to indicate the beginning of a Block literal. The body of the Block itself is contained within `{ }`, as shown in this example (as usual with C, `;` indicates the end of the statement):

The example is explained in the following illustration:

```
int multiplier = 7;

int (^myBlock)(int) = (^int num) { return num * multiplier; };
```

Notice that the Block is able to make use of variables from the same scope in which it was defined.

If you declare a Block as a variable, you can then use it just as you would a function:

```
int multiplier = 7;
```

---

27 This syntax is already part of the clang source tree on which most vendors have based their OpenCL implementations. Additionally, blocks based closures are supported by the clang open source C compiler as well as Mac OS X’s C and Objective C compilers. Specifically, Mac OS X’s Grand Central Dispatch allows applications to queue tasks as a block.
int (^myBlock)(int) = ^(int num) {
    return num * multiplier;
};

printf("%d\n", myBlock(3));  // prints 21

### 6.12.2 Declaring a Block Reference

Block variables hold references to Blocks. You declare them using syntax similar to that you use to declare a pointer to a function, except that you use ^ instead of *. The Block type fully interoperates with the rest of the C type system. The following are valid Block variable declarations:

```c
void (^blockReturningVoidWithVoidArgument)(void);
int (^blockReturningIntWithIntAndCharArguments)(int, char);
```

A Block that takes no arguments must specify void in the argument list. A Block reference may not be dereferenced via the pointer dereference operation *, and thus a Block’s size may not be computed at compile time.

Blocks are designed to be fully type safe by giving the compiler a full set of metadata to use to validate use of Blocks, parameters passed to blocks, and assignment of the return value.

You can also create types for Blocks—doing so is generally considered to be best practice when you use a block with a given signature in multiple places:

```c
typedef float (^MyBlockType)(float, float);
```

```c
MyBlockType myFirstBlock = // ...;
MyBlockType mySecondBlock = // ...;
```

### 6.12.3 Block Literal Expressions

A Block literal expression produces a reference to a Block. It is introduced by the use of the ^ token as a unary operator.

```c
Block_literal_expression ::= ^ block_decl compound_statement_body
block_decl ::= block_decl ::= parameter_list
block_decl ::= type_expression
```
where type expression is extended to allow `^` as a Block reference where `*` is allowed as a function reference.

The following Block literal:

```
^ void (void) { printf("hello world\n"); }
```

produces a reference to a Block with no arguments with no return value.

The return type is optional and is inferred from the return statements. If the return statements return a value, they all must return a value of the same type. If there is no value returned the inferred type of the Block is void; otherwise it is the type of the return statement value. If the return type is omitted and the argument list is ( void ), the ( void ) argument list may also be omitted.

So:

```
^ ( void ) { printf("hello world\n"); }
```

and:

```
^ { printf("hello world\n"); }
```

are exactly equivalent constructs for the same expression.

The compound statement body establishes a new lexical scope within that of its parent. Variables used within the scope of the compound statement are bound to the Block in the normal manner with the exception of those in automatic (stack) storage. Thus one may access functions and global variables as one would expect, as well as static local variables.

Local automatic (stack) variables referenced within the compound statement of a Block are imported and captured by the Block as const copies. The capture (binding) is performed at the time of the Block literal expression evaluation.

The compiler is not required to capture a variable if it can prove that no references to the variable will actually be evaluated.

The lifetime of variables declared in a Block is that of a function..

Block literal expressions may occur within Block literal expressions (nested) and all variables captured by any nested blocks are implicitly also captured in the scopes of their enclosing Blocks.

A Block literal expression may be used as the initialization value for Block variables at global or local static scope.
You can also declare a Block as a global literal in program scope.

```c
int GlobalInt = 0;

int (^getGlobalInt)(void) = ^( return GlobalInt; );
```

6.12.4 **Control Flow**

The compound statement of a Block is treated much like a function body with respect to control flow in that continue, break and goto do not escape the Block.

6.12.5 **Restrictions**

The following Blocks features are currently not supported in OpenCL C.

- The __block storage type.
- The Block_copy() and Block_release() functions that copy and release Blocks.
- Blocks with variadic arguments.
- Arrays of Blocks.
- Blocks as structures and union members.

Block literals are assumed to allocate memory at the point of definition and to be destroyed at the end of the same scope. To support these behaviors, additional restrictions\(^{28}\) in addition to the above feature restrictions are:

- Block variables must be defined and used in a way that allows them to be statically determinable at build or “link to executable” time. In particular:
  - Block variables assigned in one scope must be used only with the same or any nested scope.
  - The “extern” storage-class specified cannot be used with program scope block variables.
  - Block variable declarations must be qualified with const. Therefore all block variables must be initialized at declaration time and may not be reassigned.
  - A block cannot be the return value of a function.
- The unary operators (* and &) cannot be used with a Block.
- Blocks cannot be used as expressions of the ternary selection operator (? :).
- A Block cannot reference another Block.

---

\(^{28}\) OpenCL C does not allow function pointers (see section 6.9) primarily because it is difficult or expensive to implement generic indirections to executable code in many hardware architectures that OpenCL targets. OpenCL C’s design of Blocks is intended to respect that same condition, yielding the restrictions listed here. As such, Blocks allow a form of dynamically enqueued function scheduling without providing a form of runtime synchronous dynamic dispatch analogous to function pointers.
Some examples that describe legal and illegal issue of Blocks in OpenCL C are described below.

Example 1:

```c
void foo(int *x, int (^bar)(int, int))
{
    *x = bar(*x, *x);
}

kernel
void k(global int *x, global int *z)
{
    if (some expression)
        *x = foo(x,
            ^int(int x, int y){return x+y+*z;});  // legal
    else
        *x = foo(x,
            ^int(int x, int y){return (x*y)-*z;});  // legal
}
```

Example 2:

```c
kernel
void k(global int *x, global int *z)
{
    int ^(tmp)(int, int);
    if (some expression)
        { tmp = ^int(int x, int y){return x+y+*z;}); // illegal
    } *x = foo(x, tmp);
}
```

Example 3:

```c
int GlobalInt = 0;
int (^getGlobalInt)(void) = ^{ return GlobalInt; }; // legal
int (^getAnotherGlobalInt)(void); // illegal
extern int (^getExternGlobalInt)(void); // illegal

void foo()
{
    ...
    getGlobalInt = ^{ return 0; }; // illegal - cannot assign to
        // a global block variable
    ...
}
6.13 Built-in Functions

The OpenCL C programming language provides a rich set of built-in functions for scalar and vector operations. Many of these functions are similar to the function names provided in common C libraries but they support scalar and vector argument types. Applications should use the built-in functions wherever possible instead of writing their own version.

User defined OpenCL C functions, behave per C standard rules for functions (C99, TC2, Section 6.9.1). On entry to the function, the size of each variably modified parameter is evaluated and the value of each argument expression is converted to the type of the corresponding parameter as per usual arithmetic conversion rules described in section 6.2.6. Built-in functions described in this section behave similarly, except that in order to avoid ambiguity between multiple forms of the same built-in function, implicit scalar widening shall not occur. Note that some built-in functions described in this section do have forms that operate on mixed scalar and vector types, however.

6.13.1 Work-Item Functions

Table 6.7 describes the list of built-in work-item functions that can be used to query the number of dimensions, the global and local work size specified to clEnqueueNDRangeKernel, and the global and local identifier of each work-item when this kernel is being executed on a device.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>uint get_work_dim ()</td>
<td>Returns the number of dimensions in use. This is the value given to the work_dim argument specified in clEnqueueNDRangeKernel.</td>
</tr>
<tr>
<td>size_t get_global_size (uint dimindx)</td>
<td>Returns the number of global work-items specified for dimension identified by dimindx. This value is given by the global_work_size argument to clEnqueueNDRangeKernel. Valid values of dimindx are 0 to get_work_dim() – 1. For other values of dimindx, get_global_size() returns 1.</td>
</tr>
<tr>
<td>size_t get_global_id (uint dimindx)</td>
<td>Returns the unique global work-item ID value for dimension identified by dimindx. The global work-item ID specifies the work-item ID based on the number of global work-items specified to execute the kernel. Valid values of dimindx are 0 to get_work_dim() – 1. For other values of dimindx, get_global_id() returns 0.</td>
</tr>
<tr>
<td>size_t get_local_size (uint dimindx)</td>
<td>Returns the number of local work-items specified in dimension identified by dimindx. This value is at most the value given by the local_work_size argument to clEnqueueNDRangeKernel if local_work_size is not NULL; otherwise the OpenCL implementation chooses</td>
</tr>
</tbody>
</table>
an appropriate *local_work_size* value which is returned by this function. If the kernel is executed with a non-uniform work-group size\(^{29}\), calls to this built-in from some work-groups may return different values than calls to this built-in from other work-groups.

Valid values of *dimindx* are 0 to \(\text{get_work_dim}() - 1\). For other values of *dimindx*, \(\text{get_local_size}()\) returns 1.

### size_t get_enqueued_local_size (uint *dimindx*)

Returns the same value as that returned by \(\text{get_local_size}(\text{dimindx})\) if the kernel is executed with a uniform work-group size.

If the kernel is executed with a non-uniform work-group size, returns the number of local work-items in each of the work-groups that make up the uniform region of the global range in the dimension identified by *dimindx*. If the *local_work_size* argument to \(\text{clEnqueueNDRangeKernel}\) is not NULL, this value will match the value specified in *local_work_size[dimindx]*. If *local_work_size* is NULL, this value will match the local size that the implementation determined would be most efficient at implementing the uniform region of the global range.

Valid values of *dimindx* are 0 to \(\text{get_work_dim}() - 1\). For other values of *dimindx*, \(\text{get_enqueued_local_size}()\) returns 1.

### size_t get_local_id (uint *dimindx*)

Returns the unique local work-item ID i.e. a work-item within a specific work-group for dimension identified by *dimindx*. Valid values of *dimindx* are 0 to \(\text{get_work_dim}() - 1\). For other values of *dimindx*, \(\text{get_local_id}()\) returns 0.

### size_t get_num_groups (uint *dimindx*)

Returns the number of work-groups that will execute a kernel for dimension identified by *dimindx*. Valid values of *dimindx* are 0 to \(\text{get_work_dim}() - 1\). For other values of *dimindx*, \(\text{get_num_groups}()\) returns 1.

### size_t get_group_id (uint *dimindx*)

\(\text{get_group_id}\) returns the work-group ID which is a number from 0 .. \(\text{get_num_groups}(\text{dimindx}) - 1\). Valid values of *dimindx* are 0 to \(\text{get_work_dim}() - 1\). For other values, \(\text{get_group_id}()\) returns 0.

### size_t get_global_offset (uint *dimindx*)

\(\text{get_global_offset}\) returns the offset values specified in *global_work_offset* argument to \(\text{clEnqueueNDRangeKernel}\). Valid values of *dimindx* are 0 to \(\text{get_work_dim}() - 1\). For other values,
<table>
<thead>
<tr>
<th>Function Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>get_global_offset()</code></td>
<td>Returns 0.</td>
</tr>
<tr>
<td><code>size_t get_global_linear_id()</code></td>
<td>Returns the work-items 1-dimensional global ID. For 1D work-groups, it is computed as <code>get_global_id(0) – get_global_offset(0)</code>. For 2D work-groups, it is computed as <code>(get_global_id(1) – get_global_offset(1)) * get_global_size(0) + (get_global_id(0) – get_global_offset(0))</code>. For 3D work-groups, it is computed as <code>((get_global_id(2) – get_global_offset(2)) * get_global_size(1) * get_global_size(0)) + ((get_global_id(1) – get_global_offset(1)) * get_global_size(0)) + (get_global_id(0) – get_global_offset(0))</code>.</td>
</tr>
<tr>
<td><code>size_t get_local_linear_id()</code></td>
<td>Returns the work-items 1-dimensional local ID. For 1D work-groups, it is the same value as <code>get_local_id(0)</code>. For 2D work-groups, it is computed as <code>get_local_id(1) * get_local_size(0) + get_local_id(0)</code>. For 3D work-groups, it is computed as <code>(get_local_id(2) * get_local_size(1) * get_local_size(0)) + (get_local_id(1) * get_local_size(0)) + get_local_id(0)</code>.</td>
</tr>
</tbody>
</table>

Table 6.7  Work-Item Functions Table
6.13.2 Math Functions

The list of built-in math functions is described in table 6.8. The built-in math functions are categorized into the following:

- A list of built-in functions that have scalar or vector argument versions, and,
- A list of built-in functions that only take scalar float arguments.

The vector versions of the math functions operate component-wise. The description is per-component.

The built-in math functions are not affected by the prevailing rounding mode in the calling environment, and always return the same value as they would if called with the round to nearest even rounding mode.

Table 6.8 describes the list of built-in math functions that can take scalar or vector arguments. We use the generic type name *gentype* to indicate that the function can take float, float2, float3, float4, float8, float16, double, double2, double3, double4, double8 or double16 as the type for the arguments. We use the generic type name *gentypef* to indicate that the function can take float, float2, float3, float4, float8, or float16 as the type for the arguments. We use the generic type name *gentyped* to indicate that the function can take double, double2, double3, double4, double8 or double16 as the type for the arguments. For any specific use of a function, the actual type has to be the same for all arguments and the return type, unless otherwise specified.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>gentype acos (gentype)</td>
<td>Arc cosine function.</td>
</tr>
<tr>
<td>gentype acosh (gentype)</td>
<td>Inverse hyperbolic cosine.</td>
</tr>
<tr>
<td>gentype acospi (gentype x)</td>
<td>Compute $\acos(x) / \pi$.</td>
</tr>
<tr>
<td>gentype asin (gentype)</td>
<td>Arc sine function.</td>
</tr>
<tr>
<td>gentype asinh (gentype)</td>
<td>Inverse hyperbolic sine.</td>
</tr>
<tr>
<td>gentype asinpi (gentype x)</td>
<td>Compute $\asin(x) / \pi$.</td>
</tr>
<tr>
<td>gentype atan (gentype y over x)</td>
<td>Arc tangent function.</td>
</tr>
<tr>
<td>gentype atan2 (gentype y, gentype x)</td>
<td>Arc tangent of $y / x$.</td>
</tr>
<tr>
<td>gentype atanh (gentype)</td>
<td>Hyperbolic arc tangent.</td>
</tr>
<tr>
<td>gentype atanpi (gentype x)</td>
<td>Compute $\atan(x) / \pi$.</td>
</tr>
<tr>
<td>gentype atan2pi (gentype y, gentype x)</td>
<td>Compute $\atan2(y, x) / \pi$.</td>
</tr>
<tr>
<td>gentype cbrt (gentype)</td>
<td>Compute cube-root.</td>
</tr>
<tr>
<td>gentype ceil (gentype)</td>
<td>Round to integral value using the round to positive infinity rounding mode.</td>
</tr>
<tr>
<td>gentype copysign (gentype x, gentype y)</td>
<td>Returns x with its sign changed to match the sign of y.</td>
</tr>
<tr>
<td>Function</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>------------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>gentype cos (gentype)</code></td>
<td>Compute cosine.</td>
</tr>
<tr>
<td><code>gentype cosh (gentype)</code></td>
<td>Compute hyperbolic cosine.</td>
</tr>
<tr>
<td><code>gentype cospi (gentype)</code></td>
<td>Compute (\cos (\pi x)).</td>
</tr>
<tr>
<td><code>gentype erfc (gentype)</code></td>
<td>Complementary error function.</td>
</tr>
<tr>
<td><code>gentype erf (gentype)</code></td>
<td>Error function encountered in integrating the normal distribution.</td>
</tr>
<tr>
<td><code>gentype exp (gentype x)</code></td>
<td>Compute the base-(e) exponential of (x).</td>
</tr>
<tr>
<td><code>gentype exp2 (gentype)</code></td>
<td>Exponential base 2 function.</td>
</tr>
<tr>
<td><code>gentype exp10 (gentype)</code></td>
<td>Exponential base 10 function.</td>
</tr>
<tr>
<td><code>gentype expm1 (gentype)</code></td>
<td>Compute (e^x - 1.0).</td>
</tr>
<tr>
<td><code>gentype fabs (gentype)</code></td>
<td>Compute absolute value of a floating-point number.</td>
</tr>
<tr>
<td><code>gentype fdim (gentype x, gentype y)</code></td>
<td>(x - y) if (x &gt; y), +0 if (x) is less than or equal to (y).</td>
</tr>
<tr>
<td><code>gentype floor (gentype)</code></td>
<td>Round to integral value using the round to negative infinity rounding mode.</td>
</tr>
<tr>
<td><code>gentype fma (gentype a, gentype b, gentype c)</code></td>
<td>Returns the correctly rounded floating-point representation of the sum of (c) with the infinitely precise product of (a) and (b). Rounding of intermediate products shall not occur. Edge case behavior is per the IEEE 754-2008 standard.</td>
</tr>
<tr>
<td><code>gentype fmax (gentype x, gentype y)</code></td>
<td>Returns (y) if (x &lt; y), otherwise it returns (x). If one argument is a NaN, <code>fmax()</code> returns the other argument. If both arguments are NaNs, <code>fmax()</code> returns a NaN.</td>
</tr>
<tr>
<td><code>gentypef fmax (gentypef x, float y)</code></td>
<td>Returns (y) if (x &lt; y), otherwise it returns (x). If one argument is a NaN, <code>fmax()</code> returns the other argument. If both arguments are NaNs, <code>fmax()</code> returns a NaN.</td>
</tr>
<tr>
<td><code>gentyped fmax (gentyped x, double y)</code></td>
<td>Returns (y) if (y &lt; x), otherwise it returns (x). If one argument is a NaN, <code>fmax()</code> returns the other argument. If both arguments are NaNs, <code>fmax()</code> returns a NaN.</td>
</tr>
<tr>
<td><code>gentype fmin (gentype x, gentype y)</code></td>
<td>Returns (y) if (y &lt; x), otherwise it returns (x). If one argument is a NaN, <code>fmin()</code> returns the other argument. If both arguments are NaNs, <code>fmin()</code> returns a NaN.</td>
</tr>
<tr>
<td><code>gentypef fmin (gentypef x, float y)</code></td>
<td>Returns (y) if (y &lt; x), otherwise it returns (x). If one argument is a NaN, <code>fmin()</code> returns the other argument. If both arguments are NaNs, <code>fmin()</code> returns a NaN.</td>
</tr>
<tr>
<td><code>gentyped fmin (gentyped x, double y)</code></td>
<td>Returns (y) if (y &lt; x), otherwise it returns (x). If one argument is a NaN, <code>fmin()</code> returns the other argument. If both arguments are NaNs, <code>fmin()</code> returns a NaN.</td>
</tr>
<tr>
<td><code>gentype fmod (gentype x, gentype y)</code></td>
<td>Modulus. Returns (x - y * \text{trunc}(x/y)).</td>
</tr>
<tr>
<td><code>gentype fract (gentype x, gentype *iptr)</code></td>
<td>Returns (\text{fmin}(x - \text{floor}(x), 0x1.fffffep-1f)). \text{floor}(x) is returned in \text{iptr}.</td>
</tr>
<tr>
<td><code>float frexp (float x, intptr *exp)</code></td>
<td>Extract mantissa and exponent from (x). For each component the mantissa returned is a float with magnitude in the interval [1/2, 1) or 0. Each component of (x) equals mantissa returned (\times 2^{exp}).</td>
</tr>
<tr>
<td><code>float frexp (float x, intptr *exp)</code></td>
<td>Extract mantissa and exponent from (x). For each component the mantissa returned is a float with magnitude in the interval [1/2, 1) or 0. Each component of (x) equals mantissa returned (\times 2^{exp}).</td>
</tr>
<tr>
<td><code>double frexp (double x, intptr *exp)</code></td>
<td>Extract mantissa and exponent from (x). For each component the mantissa returned is a float with magnitude in the interval [1/2, 1) or 0. Each component of (x) equals mantissa returned (\times 2^{exp}).</td>
</tr>
</tbody>
</table>

---

30 `fmin` and `fmax` behave as defined by C99 and may not match the IEEE 754-2008 definition for `minNum` and `maxNum` with regard to signaling NaNs. Specifically, signaling NaNs may behave as quiet NaNs.  
31 The `min()` operator is there to prevent `fract(-small)` from returning 1.0. It returns the largest positive floating-point number less than 1.0.
<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>gentype hypot (gentype x, gentype y)</code></td>
<td>Compute the value of the square root of $x^2 + y^2$ without undue overflow or underflow.</td>
</tr>
<tr>
<td><code>int n ilogb (float n x)</code></td>
<td>Return the exponent as an integer value.</td>
</tr>
<tr>
<td><code>int n ilogb (double n x)</code></td>
<td></td>
</tr>
<tr>
<td><code>float n ldexp (float n x, int n k)</code></td>
<td>Multiply $x$ by 2 to the power $k$.</td>
</tr>
<tr>
<td><code>float n ldexp (double n x, int n k)</code></td>
<td></td>
</tr>
<tr>
<td><code>gentype lgamma (gentype x)</code></td>
<td>Log gamma function. Returns the natural logarithm of the absolute value of the gamma function. The sign of the gamma function is returned in the <code>signp</code> argument of <code>lgamma_r</code>.</td>
</tr>
<tr>
<td><code>float n lgamma_r (float n x, int n *signp)</code></td>
<td></td>
</tr>
<tr>
<td><code>float n lgamma_r (double n x, int n *signp)</code></td>
<td></td>
</tr>
<tr>
<td><code>double n lgamma_r (double n x, int n *signp)</code></td>
<td></td>
</tr>
<tr>
<td><code>gentype log (gentype)</code></td>
<td>Compute natural logarithm.</td>
</tr>
<tr>
<td><code>gentype log2 (gentype)</code></td>
<td>Compute a base 2 logarithm.</td>
</tr>
<tr>
<td><code>gentype log10 (gentype)</code></td>
<td>Compute a base 10 logarithm.</td>
</tr>
<tr>
<td><code>gentype log1p (gentype x)</code></td>
<td>Compute $\log_e(1.0 + x)$.</td>
</tr>
<tr>
<td><code>gentype logb (gentype x)</code></td>
<td>Compute the exponent of $x$, which is the integral part of $\log_e</td>
</tr>
<tr>
<td><code>gentype mad (gentype a, gentype b, gentype c)</code></td>
<td><code>mad</code> approximates $a \times b + c$. Whether or how the product of $a \times b$ is rounded and how supernormal or subnormal intermediate products are handled is not defined. <code>mad</code> is intended to be used where speed is preferred over accuracy.</td>
</tr>
<tr>
<td><code>gentype maxmag (gentype x, gentype y)</code></td>
<td>Returns $x$ if $</td>
</tr>
<tr>
<td><code>gentype minmag (gentype x, gentype y)</code></td>
<td>Returns $x$ if $</td>
</tr>
<tr>
<td><code>gentype modf (gentype x, gentype *iptr)</code></td>
<td>Decompose a floating-point number. The <code>modf</code> function breaks the argument $x$ into integral and fractional parts, each of which has the same sign as the argument. It stores the integral part in the object pointed to by <code>iptr</code>.</td>
</tr>
<tr>
<td><code>float n nan (uint n nancode)</code></td>
<td>Returns a quiet NaN. The <code>nancode</code> may be placed.</td>
</tr>
<tr>
<td>Function</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>float nan</code></td>
<td><code>uint nancode</code></td>
</tr>
<tr>
<td><code>double nan</code></td>
<td><code>ulong nancode</code></td>
</tr>
<tr>
<td><code>double nan</code></td>
<td><code>ulong nancode</code></td>
</tr>
<tr>
<td><strong>in the significand of the resulting NaN.</strong></td>
<td></td>
</tr>
<tr>
<td><code>gentype nextafter</code></td>
<td><code>gentype x</code>&lt;br&gt;<code>gentype y</code>&lt;br&gt;Computes the next representable single-precision floating-point value following <code>x</code> in the direction of <code>y</code>. Thus, if <code>y</code> is less than <code>x</code>, <code>nextafter()</code> returns the largest representable floating-point number less than <code>x</code>.</td>
</tr>
<tr>
<td><code>gentype pow</code></td>
<td><code>gentype x</code>&lt;br&gt;<code>gentype y</code>&lt;br&gt;Compute <code>x</code> to the power <code>y</code>.</td>
</tr>
<tr>
<td><code>float pown</code></td>
<td><code>float x</code>&lt;br&gt;<code>int y</code>&lt;br&gt;Compute <code>x</code> to the power <code>y</code>, where <code>y</code> is an integer.</td>
</tr>
<tr>
<td><code>double pown</code></td>
<td><code>double x</code>&lt;br&gt;<code>int y</code>&lt;br&gt;Compute <code>x</code> to the power <code>y</code>, where <code>x</code> is &gt;= 0.</td>
</tr>
<tr>
<td><code>gentype powr</code></td>
<td><code>gentype x</code>&lt;br&gt;<code>gentype y</code>&lt;br&gt;Compute <code>x</code> to the power <code>y</code>, where <code>x</code> is &gt;= 0.</td>
</tr>
<tr>
<td><code>gentype remainder</code></td>
<td><code>gentype x</code>&lt;br&gt;<code>gentype y</code>&lt;br&gt;Compute the value <code>r</code> such that <code>r = x - n*y</code>, where <code>n</code> is the integer nearest the exact value of <code>x/y</code>. If there are two integers closest to <code>x/y</code>, <code>n</code> shall be the even one. If <code>r</code> is zero, it is given the same sign as <code>x</code>.</td>
</tr>
<tr>
<td><code>float remquo</code></td>
<td><code>float x</code>&lt;br&gt;<code>float y</code>&lt;br&gt;<code>int *quo</code>&lt;br&gt;The <code>remquo</code> function computes the value <code>r</code> such that <code>r = x - k*y</code>, where <code>k</code> is the integer nearest the exact value of <code>x/y</code>. If there are two integers closest to <code>x/y</code>, <code>k</code> shall be the even one. If <code>r</code> is zero, it is given the same sign as <code>x</code>. This is the same value that is returned by the <code>remainder</code> function. <code>remquo</code> also calculates the lower seven bits of the integral quotient <code>x/y</code>, and gives that value the same sign as <code>x/y</code>. It stores this signed value in the object pointed to by <code>quo</code>.</td>
</tr>
<tr>
<td><code>double remquo</code></td>
<td><code>double x</code>&lt;br&gt;<code>double y</code>&lt;br&gt;<code>int *quo</code>&lt;br&gt;The <code>remquo</code> function computes the value <code>r</code> such that <code>r = x - k*y</code>, where <code>k</code> is the integer nearest the exact value of <code>x/y</code>. If there are two integers closest to <code>x/y</code>, <code>k</code> shall be the even one. If <code>r</code> is zero, it is given the same sign as <code>x</code>. This is the same value that is returned by the <code>remainder</code> function. <code>remquo</code> also calculates the lower seven bits of the integral quotient <code>x/y</code>, and gives that value the same sign as <code>x/y</code>. It stores this signed value in the object pointed to by <code>quo</code>.</td>
</tr>
<tr>
<td><code>gentype rint</code></td>
<td><code>gentype</code>&lt;br&gt;Round to integral value (using round to nearest even rounding mode) in floating-point format. Refer to section 7.1 for description of rounding modes.</td>
</tr>
<tr>
<td><code>float rootn</code></td>
<td><code>float x</code>&lt;br&gt;<code>int y</code>&lt;br&gt;Compute <code>x</code> to the power <code>1/y</code>.</td>
</tr>
<tr>
<td>Function</td>
<td>Description</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
</tr>
<tr>
<td><code>double rootn (double x, int y)</code></td>
<td>Return the integral value nearest to ( x ) rounding halfway cases away from zero, regardless of the current rounding direction.</td>
</tr>
<tr>
<td><code>double rootn (double x, int y)</code></td>
<td></td>
</tr>
<tr>
<td><code>gentype round (gentype x)</code></td>
<td>Return the integral value nearest to ( x ) rounding halfway cases away from zero, regardless of the current rounding direction.</td>
</tr>
<tr>
<td><code>gentype rsqrt (gentype)</code></td>
<td>Compute inverse square root.</td>
</tr>
<tr>
<td><code>gentype sin (gentype)</code></td>
<td>Compute sine.</td>
</tr>
<tr>
<td><code>gentype sincos (gentype x, gentype *cosval)</code></td>
<td>Compute sine and cosine of ( x ). The computed sine is the return value and computed cosine is returned in ( \cosval ).</td>
</tr>
<tr>
<td><code>gentype sinh (gentype)</code></td>
<td>Compute hyperbolic sine.</td>
</tr>
<tr>
<td><code>gentype sinpi (gentype x)</code></td>
<td>Compute ( \sin (\pi x) ).</td>
</tr>
<tr>
<td><code>gentype sqrt (gentype)</code></td>
<td>Compute square root.</td>
</tr>
<tr>
<td><code>gentype tan (gentype)</code></td>
<td>Compute tangent.</td>
</tr>
<tr>
<td><code>gentype tanh (gentype x)</code></td>
<td>Compute hyperbolic tangent.</td>
</tr>
<tr>
<td><code>gentype tanpi (gentype x)</code></td>
<td>Compute ( \tan (\pi x) ).</td>
</tr>
<tr>
<td><code>gentype tgamma (gentype)</code></td>
<td>Compute the gamma function.</td>
</tr>
<tr>
<td><code>gentype trunc (gentype)</code></td>
<td>Round to integral value using the round to zero rounding mode.</td>
</tr>
</tbody>
</table>

**Table 6.8**  **Scalar and Vector Argument Built-in Math Function Table**

*Table 6.9* describes the following functions:

- A subset of functions from *table 6.8* that are defined with the `half_` prefix. These functions are implemented with a minimum of 10-bits of accuracy i.e. an ULP value \( \leq 8192 \text{ ulp} \).

- A subset of functions from *table 6.8* that are defined with the `native_` prefix. These functions may map to one or more native device instructions and will typically have better performance compared to the corresponding functions (without the `native_` prefix) described in *table 6.8*. The accuracy (and in some cases the input range(s)) of these functions is implementation-defined.

- `half_` and `native_` functions for following basic operations: divide and reciprocal.

We use the generic type name `gentype` to indicate that the functions in *table 6.9* can take `float`, `float2`, `float3`, `float4`, `float8` or `float16` as the type for the arguments.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>gentype half_cos (gentype x)</code></td>
<td>Compute cosine. ( x ) must be in the range (-2^{16} \ldots +2^{16}).</td>
</tr>
<tr>
<td><code>gentype half_divide (gentype x, ...)</code></td>
<td>Compute ( x / y ).</td>
</tr>
<tr>
<td>Function</td>
<td>Description</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
</tr>
<tr>
<td><code>gentype half_exp (gentype x)</code></td>
<td>Compute the base- e exponential of x.</td>
</tr>
<tr>
<td><code>gentype half_exp2 (gentype x)</code></td>
<td>Compute the base- 2 exponential of x.</td>
</tr>
<tr>
<td><code>gentype half_exp10 (gentype x)</code></td>
<td>Compute the base- 10 exponential of x.</td>
</tr>
<tr>
<td><code>gentype half_log (gentype x)</code></td>
<td>Compute natural logarithm.</td>
</tr>
<tr>
<td><code>gentype half_log2 (gentype x)</code></td>
<td>Compute a base 2 logarithm.</td>
</tr>
<tr>
<td><code>gentype half_log10 (gentype x)</code></td>
<td>Compute a base 10 logarithm.</td>
</tr>
<tr>
<td><code>gentype half_powr (gentype x, gentype y)</code></td>
<td>Compute x to the power y, where x is &gt;= 0.</td>
</tr>
<tr>
<td><code>gentype half_recip (gentype x)</code></td>
<td>Compute reciprocal.</td>
</tr>
<tr>
<td><code>gentype half_rsqrt (gentype x)</code></td>
<td>Compute inverse square root.</td>
</tr>
<tr>
<td><code>gentype half_sin (gentype x)</code></td>
<td>Compute sine. x must be in the range (-2^{16} \ldots +2^{16}).</td>
</tr>
<tr>
<td><code>gentype half_sqrt (gentype x)</code></td>
<td>Compute square root.</td>
</tr>
<tr>
<td><code>gentype half_tan (gentype x)</code></td>
<td>Compute tangent. x must be in the range (-2^{16} \ldots +2^{16}).</td>
</tr>
<tr>
<td><code>gentype native_cos (gentype x)</code></td>
<td>Compute cosine over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_divide (gentype x, gentype y)</code></td>
<td>Compute x / y over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_exp (gentype x)</code></td>
<td>Compute the base- e exponential of x over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_exp2 (gentype x)</code></td>
<td>Compute the base- 2 exponential of x over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_exp10 (gentype x)</code></td>
<td>Compute the base- 10 exponential of x over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_log (gentype x)</code></td>
<td>Compute natural logarithm over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_log2 (gentype x)</code></td>
<td>Compute a base 2 logarithm over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_log10 (gentype x)</code></td>
<td>Compute a base 10 logarithm over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_powr (gentype x, gentype y)</code></td>
<td>Compute x to the power y, where x is &gt;= 0. The range of x and y are implementation-defined. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_recip (gentype x)</code></td>
<td>Compute reciprocal over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
<tr>
<td><code>gentype native_rsqrt (gentype x)</code></td>
<td>Compute inverse square root over an implementation-defined range. The maximum error is implementation-defined.</td>
</tr>
</tbody>
</table>
Table 6.9  Scalar and Vector Argument Built-in half__ and native__ Math Functions

Support for denormal values is optional for half__ functions. The half__ functions may return any result allowed by section 7.5.3, even when -cl-denorms-are-zero (see section 5.8.4.2) is not in force. Support for denormal values is implementation-defined for native__ functions.

The following symbolic constants are available. Their values are of type float and are accurate within the precision of a single precision floating-point number.

<table>
<thead>
<tr>
<th>Constant Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MAXFLOATE</td>
<td>Value of maximum non-infinite single-precision floating-point number.</td>
</tr>
<tr>
<td>HUGE_VALF</td>
<td>A positive float constant expression. HUGE_VALF evaluates to +infinity. Used as an error value returned by the built-in math functions.</td>
</tr>
<tr>
<td>INFINITY</td>
<td>A constant expression of type float representing positive or unsigned infinity.</td>
</tr>
<tr>
<td>NAN</td>
<td>A constant expression of type float representing a quiet NaN.</td>
</tr>
</tbody>
</table>

If double precision is supported by the device, the following symbolic constant will also be available:

<table>
<thead>
<tr>
<th>Constant Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>HUGE_VAL</td>
<td>A positive double constant expression. HUGE_VAL evaluates to +infinity. Used as an error value returned by the built-in math functions.</td>
</tr>
</tbody>
</table>

6.13.2.1  Floating-point macros and pragmas

The FP_CONTRACT pragma can be used to allow (if the state is on) or disallow (if the state is off) the implementation to contract expressions. Each pragma can occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect from its occurrence until another FP_CONTRACT pragma is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect from its occurrence until another FP_CONTRACT pragma is encountered (including within a nested compound statement), or until the end of the
compound statement; at the end of a compound statement the state for the pragma is restored to its condition just before the compound statement. If this pragma is used in any other context, the behavior is undefined.

The pragma definition to set `FP_CONTRACT` is:

```c
#pragma OPENCL FP_CONTRACT on-off-switch
```

`on-off-switch` is one of:

- **ON**, **OFF** or **DEFAULT**.
- The **DEFAULT** value is **ON**.

The `FP_FAST_FMAF` macro indicates whether the `fma` function is fast compared with direct code for single precision floating-point. If defined, the `FP_FAST_FMAF` macro shall indicate that the `fma` function generally executes about as fast as, or faster than, a multiply and an add of `float` operands.

The macro names given in the following list must use the values specified. These constant expressions are suitable for use in `#if` preprocessing directives.

```c
#define FLT_DIG 6
#define FLT_MANT_DIG 24
#define FLT_MAX_10_EXP +38
#define FLT_MAX_EXP +128
#define FLT_MIN_10_EXP -37
#define FLT_MIN_EXP -125
#define FLT_RADIX 2
#define FLT_MAX 0x1.fffffep127f
#define FLT_MIN 0x1.0p-126f
#define FLT_EPSILON 0x1.0p-23f
```

The following table describes the built-in macro names given above in the OpenCL C programming language and the corresponding macro names available to the application.

<table>
<thead>
<tr>
<th>Macro in OpenCL Language</th>
<th>Macro for application</th>
</tr>
</thead>
<tbody>
<tr>
<td>FLT_DIG</td>
<td>CL_FLT_DIG</td>
</tr>
<tr>
<td>FLT_MANT_DIG</td>
<td>CL_FLT_MANT_DIG</td>
</tr>
<tr>
<td>FLT_MAX_10_EXP</td>
<td>CL_FLT_MAX_10_EXP</td>
</tr>
<tr>
<td>FLT_MAX_EXP</td>
<td>CL_FLT_MAX</td>
</tr>
<tr>
<td>FLT_MIN_10_EXP</td>
<td>CL_FLT_MIN_10_EXP</td>
</tr>
<tr>
<td>FLT_MIN_EXP</td>
<td>CL_FLT_MIN</td>
</tr>
<tr>
<td>FLT_RADIX</td>
<td>CL_FLT_RADIX</td>
</tr>
<tr>
<td>FLT_MAX</td>
<td>CL_FLT_MAX</td>
</tr>
<tr>
<td>FLT_MIN</td>
<td>CL_FLT_MIN</td>
</tr>
<tr>
<td>FLT_EPSILON</td>
<td>CL_FLT_EPSILON</td>
</tr>
</tbody>
</table>
The following macros shall expand to integer constant expressions whose values are returned by \texttt{ilogb}(x) if \(x\) is zero or NaN, respectively. The value of \texttt{FP_ILOGB0} shall be either \{\texttt{INT_MIN}\} or \{-\{\texttt{INT_MAX}\}\}. The value of \texttt{FP_ILOGBNAN} shall be either \{\texttt{INT_MAX}\} or \{\texttt{INT_MIN}\}.

The following constants are also available. They are of type \texttt{float} and are accurate within the precision of the \texttt{float} type.

<table>
<thead>
<tr>
<th>Constant</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{M_E}</td>
<td>Value of e</td>
</tr>
<tr>
<td>\texttt{M_LOG2E}</td>
<td>Value of (\log_2 e)</td>
</tr>
<tr>
<td>\texttt{M_LOG10E}</td>
<td>Value of (\log_{10} e)</td>
</tr>
<tr>
<td>\texttt{M_LN2}</td>
<td>Value of (\log_2)</td>
</tr>
<tr>
<td>\texttt{M_LN10}</td>
<td>Value of (\log_{10})</td>
</tr>
<tr>
<td>\texttt{M_PI}</td>
<td>Value of (\pi)</td>
</tr>
<tr>
<td>\texttt{M_PI2}</td>
<td>Value of (\pi / 2)</td>
</tr>
<tr>
<td>\texttt{M_PI4}</td>
<td>Value of (\pi / 4)</td>
</tr>
<tr>
<td>\texttt{M_1_PI}</td>
<td>Value of (1 / \pi)</td>
</tr>
<tr>
<td>\texttt{M_2_PI}</td>
<td>Value of (2 / \pi)</td>
</tr>
<tr>
<td>\texttt{M_2_SQRTPI}</td>
<td>Value of (2 / \sqrt{\pi})</td>
</tr>
<tr>
<td>\texttt{M_SQRT2}</td>
<td>Value of (\sqrt{2})</td>
</tr>
<tr>
<td>\texttt{M_SQRT1_2}</td>
<td>Value of (1 / \sqrt{2})</td>
</tr>
</tbody>
</table>

If double precision is supported by the device, the following macros and constants are also available:

The \texttt{FP_FAST_FMA} macro indicates whether the \texttt{fma()} family of functions are fast compared with direct code for double precision floating-point. If defined, the \texttt{FP_FAST_FMA} macro shall indicate that the \texttt{fma()} function generally executes about as fast as, or \(\tilde{\text{faster}}\) than, a multiply and an add of \texttt{double} operands.

The macro names given in the following list must use the values specified. These constant expressions are suitable for use in \#if preprocessing directives.

```c
#define DBL_DIG 15
#define DBL_MANT_DIG 53
#define DBL_MAX_10_EXP +308
#define DBL_MAX_EXP +1024
#define DBL_MIN_10_EXP -307
#define DBL_MIN_EXP -1021
#define DBL_MAX 0x1.fffffffffffffp1023
#define DBL_MIN 0x1.0p-1022
#define DBL_EPSILON 0x1.0p-52
```

The following table describes the built-in macro names given above in the OpenCL C programming language and the corresponding macro names available to the application.
The following constants are also available. They are of type `double` and are accurate within the precision of the `double` type.

<table>
<thead>
<tr>
<th>Constant</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>M_E</td>
<td>Value of e</td>
</tr>
<tr>
<td>M_LOG2E</td>
<td>Value of log₂e</td>
</tr>
<tr>
<td>M_LOG10E</td>
<td>Value of log₁₀e</td>
</tr>
<tr>
<td>M_LN2</td>
<td>Value of log₂</td>
</tr>
<tr>
<td>M_LN10</td>
<td>Value of log₁₀</td>
</tr>
<tr>
<td>M_PI</td>
<td>Value of π</td>
</tr>
<tr>
<td>M_PI_2</td>
<td>Value of π / 2</td>
</tr>
<tr>
<td>M_PI_4</td>
<td>Value of π / 4</td>
</tr>
<tr>
<td>M_1_PI</td>
<td>Value of 1 / π</td>
</tr>
<tr>
<td>M_2_PI</td>
<td>Value of 2 / π</td>
</tr>
<tr>
<td>M_2_SQRTP1</td>
<td>Value of 2 / √π</td>
</tr>
<tr>
<td>M_SQRT2</td>
<td>Value of √2</td>
</tr>
<tr>
<td>M_SQRT1_2</td>
<td>Value of 1 / √2</td>
</tr>
</tbody>
</table>
6.13.3 Integer Functions

Table 6.10 describes the built-in integer functions that take scalar or vector arguments. The vector versions of the integer functions operate component-wise. The description is per-component.

We use the generic type name gentype to indicate that the function can take char, char{2|3|4|8|16}, uchar, uchar{2|3|4|8|16}, short, short{2|3|4|8|16}, ushort, ushort{2|3|4|8|16}, int, int{2|3|4|8|16}, uint, uint{2|3|4|8|16}, long, long{2|3|4|8|16} ulong, or ulong{2|3|4|8|16} as the type for the arguments. We use the generic type name ugentype to refer to unsigned versions of gentype. For example, if gentype is char4, ugentype is uchar4. We also use the generic type name sgentype to indicate that the function can take a scalar data type i.e. char, uchar, short, ushort, int, uint, long, or ulong as the type for the arguments. For built-in integer functions that take gentype and sgentype arguments, the gentype argument must be a vector or scalar version of the sgentype argument. For example, if sgentype is uchar, gentype must be uchar or uchar{2|3|4|8|16}. For vector versions, sgentype is implicitly widened to gentype as described in section 6.3a.

For any specific use of a function, the actual type has to be the same for all arguments and the return type unless otherwise specified.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ugentype abs (gentype x)</td>
<td>Returns</td>
</tr>
<tr>
<td>ugentype abs_diff (gentype x, gentype y)</td>
<td>Returns</td>
</tr>
<tr>
<td>gentype add_sat (gentype x, gentype y)</td>
<td>Returns x + y and saturates the result.</td>
</tr>
<tr>
<td>gentype hadd (gentype x, gentype y)</td>
<td>Returns (x + y) &gt;&gt; 1. The intermediate sum does not modulo overflow.</td>
</tr>
<tr>
<td>gentype rhadd (gentype x, gentype y)</td>
<td>Returns (x + y + 1) &gt;&gt; 1. The intermediate sum does not modulo overflow.</td>
</tr>
<tr>
<td>gentype clamp (gentype x, gentype minval, gentype maxval)</td>
<td>Returns min(max(x, minval), maxval).</td>
</tr>
<tr>
<td>gentype clamp (gentype x, sgentype minval, sgentype maxval)</td>
<td>Results are undefined if minval &gt; maxval.</td>
</tr>
<tr>
<td>gentype clz (gentype x)</td>
<td>Returns the number of leading 0-bits in x, starting at the most significant bit position. If x is 0, returns the size in bits of the type of x or</td>
</tr>
</tbody>
</table>

---

33 Frequently vector operations need n + 1 bits temporarily to calculate a result. The rhadd instruction gives you an extra bit without needing to upsample and downsample. This can be a profound performance win.
Table 6.10  Scalar and Vector Integer Argument Built-in Functions

*Table 6.11* describes fast integer functions that can be used for optimizing performance of kernels. We use the generic type name `gentype` to indicate that the function can take `int, int2, int3, int4, int8, int16, uint, uint2, uint3, uint4, uint8` or `uint16` as the type for the arguments.
<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>gentype <strong>mad24</strong> (gentype x, gentype y, gentype z)</td>
<td>Multiply two 24-bit integer values x and y and add the 32-bit integer result to the 32-bit integer z. Refer to definition of <strong>mul24</strong> to see how the 24-bit integer multiplication is performed.</td>
</tr>
<tr>
<td>gentype <strong>mul24</strong> (gentype x, gentype y)</td>
<td>Multiply two 24-bit integer values x and y. x and y are 32-bit integers but only the low 24-bits are used to perform the multiplication. <strong>mul24</strong> should only be used when values in x and y are in the range ([-2^{23}, 2^{23}-1]) if x and y are signed integers and in the range ([0, 2^{24}-1]) if x and y are unsigned integers. If x and y are not in this range, the multiplication result is implementation-defined.</td>
</tr>
</tbody>
</table>

**Table 6.11  Fast Integer Built-in Functions**

The macro names given in the following list must use the values specified. The values shall all be constant expressions suitable for use in \#if preprocessing directives.

```c
#define CHAR_BIT 8
#define CHAR_MAX SCHAR_MAX
#define CHAR_MIN SCHAR_MIN
#define INT_MAX 2147483647
#define INT_MIN (-2147483647 - 1)
#define LONG_MAX 0x7fffffffffffffffL
#define LONG_MIN (-0x7fffffffffffffffL - 1)
#define SCHAR_MAX 127
#define SCHAR_MIN (-127 - 1)
#define SHRT_MAX 32767
#define SHRT_MIN (-32767 - 1)
#define UCHAR_MAX 255
#define USHRT_MAX 65535
#define UINT_MAX 0xffffffff
#define ULONG_MAX 0xffffffffffffffffUL
```

The following table describes the built-in macro names given above in the OpenCL C programming language and the corresponding macro names available to the application.

<table>
<thead>
<tr>
<th>Macro in OpenCL Language</th>
<th>Macro for application</th>
</tr>
</thead>
<tbody>
<tr>
<td>CHAR_BIT</td>
<td>CL_CHAR_BIT</td>
</tr>
<tr>
<td>CHAR_MAX</td>
<td>CL_CHAR_MAX</td>
</tr>
<tr>
<td>CHAR_MIN</td>
<td>CL_CHAR_MIN</td>
</tr>
<tr>
<td>INT_MAX</td>
<td>CL_INT_MAX</td>
</tr>
<tr>
<td>INT_MIN</td>
<td>CL_INT_MIN</td>
</tr>
<tr>
<td>LONG_MAX</td>
<td>CL_LONG_MAX</td>
</tr>
<tr>
<td>Variable</td>
<td>Definition</td>
</tr>
<tr>
<td>--------------</td>
<td>----------------</td>
</tr>
<tr>
<td>LONG_MIN</td>
<td>CL_LONG_MIN</td>
</tr>
<tr>
<td>SCHAR_MAX</td>
<td>CL_SCHAR_MAX</td>
</tr>
<tr>
<td>SCHAR_MIN</td>
<td>CL_SCHAR_MIN</td>
</tr>
<tr>
<td>SHRT_MAX</td>
<td>CL_SHRT_MAX</td>
</tr>
<tr>
<td>SHRT_MIN</td>
<td>CL_SHRT_MIN</td>
</tr>
<tr>
<td>UCHAR_MAX</td>
<td>CL_UCHAR_MAX</td>
</tr>
<tr>
<td>USHRT_MAX</td>
<td>CL_USHRT_MAX</td>
</tr>
<tr>
<td>UINT_MAX</td>
<td>CL_UINT_MAX</td>
</tr>
<tr>
<td>ULONG_MAX</td>
<td>CLULONG_MAX</td>
</tr>
</tbody>
</table>
6.13.4 Common Functions

Table 6.12 describes the list of built-in common functions. These all operate component-wise. The description is per-component. We use the generic type name `gentype` to indicate that the function can take `float`, `float2`, `float3`, `float4`, `float8`, `float16`, `double`, `double2`, `double3`, `double4`, `double8` or `double16` as the type for the arguments. We use the generic type name `gentypef` to indicate that the function can take `float`, `float2`, `float3`, `float4`, `float8`, or `float16` as the type for the arguments. We use the generic type name `gentyped` to indicate that the function can take `double`, `double2`, `double3`, `double4`, `double8` or `double16` as the type for the arguments.

The built-in common functions are implemented using the round to nearest even rounding mode.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>gentype clamp</code></td>
<td>Returns ( f\min(f\max(x, \text{minval}), \text{maxval}) ). Results are undefined if ( \text{minval} &gt; \text{maxval} ).</td>
</tr>
<tr>
<td><code>gentypef clamp</code></td>
<td></td>
</tr>
<tr>
<td><code>gentyped clamp</code></td>
<td></td>
</tr>
<tr>
<td><code>gentype degrees</code></td>
<td>Converts ( \text{radians} ) to degrees, i.e. ((180 / \pi) \times \text{radians}).</td>
</tr>
<tr>
<td><code>gentype max</code></td>
<td>Returns ( y ) if ( x &lt; y ), otherwise it returns ( x ). If ( x ) or ( y ) are infinite or NaN, the return values are undefined.</td>
</tr>
<tr>
<td><code>gentypef max</code></td>
<td></td>
</tr>
<tr>
<td><code>gentyped max</code></td>
<td></td>
</tr>
<tr>
<td><code>gentype min</code></td>
<td>Returns ( y ) if ( y &lt; x ), otherwise it returns ( x ). If ( x ) or ( y ) are infinite or NaN, the return values are undefined.</td>
</tr>
<tr>
<td><code>gentypef min</code></td>
<td></td>
</tr>
<tr>
<td><code>gentyped min</code></td>
<td></td>
</tr>
<tr>
<td><code>gentype mix</code></td>
<td>Returns the linear blend of ( x ) &amp; ( y ) implemented as: ( x + (y - x) \times a )</td>
</tr>
</tbody>
</table>

34 The `mix` and `smoothstep` functions can be implemented using contractions such as `mad` or `fma`. 
<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>gentype mix</code> (<code>gentype x</code>, <code>gentype y</code>, <code>float a</code>)</td>
<td><code>a</code> must be a value in the range 0.0 ... 1.0. If <code>a</code> is not in the range 0.0 ... 1.0, the return values are undefined.</td>
</tr>
<tr>
<td><code>gentyped mix</code> (<code>gentyped x</code>, <code>gentyped y</code>, <code>double a</code>)</td>
<td></td>
</tr>
<tr>
<td><code>gentype radians</code> (<code>gentype degrees</code>)</td>
<td>Converts <code>degrees</code> to radians, i.e. <code>(\pi / 180) * degrees</code>.</td>
</tr>
<tr>
<td><code>gentype step</code> (<code>gentype edge</code>, <code>gentype x</code>)</td>
<td>Returns 0.0 if <code>x &lt; edge</code>, otherwise it returns 1.0.</td>
</tr>
<tr>
<td><code>gentypef step</code> (<code>float edge</code>, <code>gentypef x</code>)</td>
<td></td>
</tr>
<tr>
<td><code>gentyped step</code> (<code>double edge</code>, <code>gentyped x</code>)</td>
<td>Returns 0.0 if <code>x &lt;= edge0</code> and 1.0 if <code>x &gt;= edge1</code> and performs smooth Hermite interpolation between 0 and 1 when <code>edge0 &lt; x &lt; edge1</code>. This is useful in cases where you would want a threshold function with a smooth transition. This is equivalent to:</td>
</tr>
<tr>
<td><code>gentype smoothstep</code> (<code>gentype edge0</code>, <code>gentype edge1</code>, <code>gentype x</code>)</td>
<td></td>
</tr>
<tr>
<td><code>gentypef smoothstep</code> (<code>float edge0</code>, <code>float edge1</code>, <code>gentypef x</code>)</td>
<td></td>
</tr>
<tr>
<td><code>gentyped smoothstep</code> (<code>double edge0</code>, <code>double edge1</code>, <code>gentyped x</code>)</td>
<td></td>
</tr>
<tr>
<td><code>gentype sign</code> (<code>gentype x</code>)</td>
<td>Returns 1.0 if <code>x &gt; 0</code>, -0.0 if <code>x = -0.0</code>, +0.0 if <code>x = +0.0</code>, or -1.0 if <code>x &lt; 0</code>. Returns 0.0 if <code>x</code> is a NaN.</td>
</tr>
</tbody>
</table>

**Table 6.12**  
*Scalar and Vector Argument Built-in Common Function Table*
### 6.13.5 Geometric Functions

Table 6.13 describes the list of built-in geometric functions. These all operate component-wise. The description is per-component. float\(n\) is float, float2, float3, or float4 and double\(n\) is double, double2, double3, or double4. The built-in geometric functions are implemented using the round to nearest even rounding mode.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>float4 <strong>cross</strong></td>
<td>Computes the cross product of (p0.xyz) and (p1.xyz). The (w) component of float4 result returned will be 0.0.</td>
</tr>
<tr>
<td>float3 <strong>cross</strong></td>
<td>Computes the cross product of (p0.xyz) and (p1.xyz).</td>
</tr>
<tr>
<td>double4 <strong>cross</strong></td>
<td>Computes the cross product of (p0.xyz) and (p1.xyz).</td>
</tr>
<tr>
<td>double3 <strong>cross</strong></td>
<td>Computes the cross product of (p0.xyz) and (p1.xyz).</td>
</tr>
<tr>
<td>float <strong>dot</strong></td>
<td>Computes the dot product.</td>
</tr>
<tr>
<td>double <strong>dot</strong></td>
<td>Computes the dot product.</td>
</tr>
<tr>
<td>float <strong>distance</strong></td>
<td>Computes the distance between (p0) and (p1). This is calculated as (\text{length}(p0 - p1)).</td>
</tr>
<tr>
<td>double <strong>distance</strong></td>
<td>Computes the distance between (p0) and (p1). This is calculated as (\text{length}(p0 - p1)).</td>
</tr>
<tr>
<td>float <strong>length</strong></td>
<td>Computes the length of vector (p), i.e., (\sqrt{p.x^2 + p.y^2 + \ldots}).</td>
</tr>
<tr>
<td>double <strong>length</strong></td>
<td>Computes the length of vector (p), i.e., (\sqrt{p.x^2 + p.y^2 + \ldots}).</td>
</tr>
<tr>
<td>floatn <strong>normalize</strong></td>
<td>Returns a vector in the same direction as (p) but with a length of 1.</td>
</tr>
<tr>
<td>doublen <strong>normalize</strong></td>
<td>Returns a vector in the same direction as (p) but with a length of 1.</td>
</tr>
<tr>
<td>float <strong>fast_distance</strong></td>
<td>Computes the distance between (p0) and (p1) as (\text{fast_length}(p0 - p1)).</td>
</tr>
<tr>
<td>float <strong>fast_length</strong></td>
<td>Computes the length of vector (p) as (\text{half_sqrt}(p.x^2 + p.y^2 + \ldots)).</td>
</tr>
<tr>
<td>floatn <strong>fast_normalize</strong></td>
<td>Returns a vector in the same direction as (p) but with a length of 1. computed as: (p * \text{half_rsqrt}(p.x^2 + p.y^2 + \ldots)).</td>
</tr>
</tbody>
</table>

The result shall be within 8192 ulps error from the infinitely precise result of...
if (all(p == 0.0f))
    result = p;
else
    result = p / sqrt(p.x^2 + p.y^2 + ...);

with the following exceptions:

1) If the sum of squares is greater than FLT_MAX then the value of the floating-point values in the result vector are undefined.

2) If the sum of squares is less than FLT_MIN then the implementation may return back p.

3) If the device is in “denoms are flushed to zero” mode, individual operand elements with magnitude less than sqrt(FLT_MIN) may be flushed to zero before proceeding with the calculation.

Table 6.13  Scalar and Vector Argument Built-in Geometric Function Table
6.13.6 Relational Functions

The relational and equality operators (<, <=, >, >=, !=, ==) can be used with scalar and vector built-in types and produce a scalar or vector signed integer result respectively as described in section 6.3.

The functions described in table 6.14 can be used with built-in scalar or vector types as arguments and return a scalar or vector integer result. The argument type gentype refers to the following built-in types: char, charn, uchar, ucharn, short, shortn, ushort, ushortn, int, intn, uint, uintn, long, longn, ulong, ulongn, float, floatn, double, and doublen. The argument type igentype refers to the built-in signed integer types i.e. char, charn, short, shortn, int, intn, long and longn. The argument type ugentype refers to the built-in unsigned integer types i.e. uchar, ucharn, ushort, ushortn, uint, uintn, ulong and ulongn. n is 2, 3, 4, 8, or 16.

The functions isequal, isnotequal, isgreater, isgreateerequal, isless, islessequal, islessgreater, isfinite, isnf, isnan, isnormal, isordered, isunordered and signbit described in table 6.14 shall return a 0 if the specified relation is false and a 1 if the specified relation is true for scalar argument types. These functions shall return a 0 if the specified relation is false and a −1 (i.e. all bits set) if the specified relation is true for vector argument types.

The relational functions isequal, isgreater, isgreateerequal, isless, islessequal, and islessgreater always return 0 if either argument is not a number (NaN). isnotequal returns 1 if one or both arguments are not a number (NaN) and the argument type is a scalar and returns -1 if one or both arguments are not a number (NaN) and the argument type is a vector.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>int isequal (float x, float y)</td>
<td>Returns the component-wise compare of ( x == y ).</td>
</tr>
<tr>
<td>intn isequal (floatn x, floatn y)</td>
<td></td>
</tr>
<tr>
<td>int isequal (double x, double y)</td>
<td></td>
</tr>
<tr>
<td>longn isequal (doublen x, doublen y)</td>
<td></td>
</tr>
<tr>
<td>int isnotequal (float x, float y)</td>
<td>Returns the component-wise compare of ( x \neq y ).</td>
</tr>
<tr>
<td>intn isnotequal (floatn x, floatn y)</td>
<td></td>
</tr>
<tr>
<td>int isnotequal (double x, double y)</td>
<td></td>
</tr>
<tr>
<td>longn isnotequal (doublen x, doublen y)</td>
<td></td>
</tr>
<tr>
<td>int isgreater (float x, float y)</td>
<td>Returns the component-wise compare of ( x &gt; y ).</td>
</tr>
</tbody>
</table>

36 If an implementation extends this specification to support IEEE-754 flags or exceptions, then all builtin functions defined in table 6.14 shall proceed without raising the invalid floating-point exception when one or more of the operands are NaNs.
<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>int n isgreater (float n, float y)</code></td>
<td>Returns the component-wise compare of ( x \geq y ).</td>
</tr>
<tr>
<td><code>int isgreater (double x, double y)</code></td>
<td>Returns the component-wise compare of ( x \geq y ).</td>
</tr>
<tr>
<td><code>long n isgreater (double n x, double n y)</code></td>
<td>Returns the component-wise compare of ( x \geq y ).</td>
</tr>
<tr>
<td><code>int isgreaterequal (float x, float y)</code></td>
<td>Returns the component-wise compare of ( x \geq y ).</td>
</tr>
<tr>
<td><code>int isgreaterequal (double x, double y)</code></td>
<td>Returns the component-wise compare of ( x \geq y ).</td>
</tr>
<tr>
<td><code>long n isgreaterequal (double n x, double n y)</code></td>
<td>Returns the component-wise compare of ( x \geq y ).</td>
</tr>
<tr>
<td><code>int isless (float x, float y)</code></td>
<td>Returns the component-wise compare of ( x &lt; y ).</td>
</tr>
<tr>
<td><code>int isless (double x, double y)</code></td>
<td>Returns the component-wise compare of ( x &lt; y ).</td>
</tr>
<tr>
<td><code>long n isless (double n x, double n y)</code></td>
<td>Returns the component-wise compare of ( x &lt; y ).</td>
</tr>
<tr>
<td><code>int islessequal (float x, float y)</code></td>
<td>Returns the component-wise compare of ( x \leq y ).</td>
</tr>
<tr>
<td><code>int islessequal (double x, double y)</code></td>
<td>Returns the component-wise compare of ( x \leq y ).</td>
</tr>
<tr>
<td><code>long n islessequal (double n x, double n y)</code></td>
<td>Returns the component-wise compare of ( x \leq y ).</td>
</tr>
<tr>
<td><code>int islessgreater (float x, float y)</code></td>
<td>Returns the component-wise compare of ( (x &lt; y) \lor (x &gt; y) ).</td>
</tr>
<tr>
<td><code>int islessgreater (double x, double y)</code></td>
<td>Returns the component-wise compare of ( (x &lt; y) \lor (x &gt; y) ).</td>
</tr>
<tr>
<td><code>long n islessgreater (double n x, double n y)</code></td>
<td>Returns the component-wise compare of ( (x &lt; y) \lor (x &gt; y) ).</td>
</tr>
<tr>
<td><code>int isfinite (float)</code></td>
<td>Test for finite value.</td>
</tr>
<tr>
<td><code>int isfinite (float n)</code></td>
<td>Test for finite value.</td>
</tr>
<tr>
<td><code>int isfinite (double)</code></td>
<td>Test for finite value.</td>
</tr>
<tr>
<td><code>long n isfinite (double n)</code></td>
<td>Test for finite value.</td>
</tr>
<tr>
<td><code>int isinf (float)</code></td>
<td>Test for infinity value (positive or negative).</td>
</tr>
<tr>
<td><code>int isinf (float n)</code></td>
<td>Test for infinity value (positive or negative).</td>
</tr>
<tr>
<td><code>int isinf (double)</code></td>
<td>Test for infinity value (positive or negative).</td>
</tr>
<tr>
<td><code>long n isinf (double n)</code></td>
<td>Test for infinity value (positive or negative).</td>
</tr>
<tr>
<td><code>int isnan (float)</code></td>
<td>Test for a NaN.</td>
</tr>
<tr>
<td><code>int isnan (float n)</code></td>
<td>Test for a NaN.</td>
</tr>
<tr>
<td><code>int isnan (double)</code></td>
<td>Test for a NaN.</td>
</tr>
<tr>
<td><code>long n isnan (double n)</code></td>
<td>Test for a NaN.</td>
</tr>
<tr>
<td><code>int isnormal (float)</code></td>
<td>Test for a normal value.</td>
</tr>
<tr>
<td><code>int isnormal (float n)</code></td>
<td>Test for a normal value.</td>
</tr>
<tr>
<td><code>int isnormal (double)</code></td>
<td>Test for a normal value.</td>
</tr>
<tr>
<td>Function</td>
<td>Description</td>
</tr>
<tr>
<td>-----------</td>
<td>-------------</td>
</tr>
<tr>
<td>long int isnormal (double n)</td>
<td>Test if arguments are ordered. isordered() takes arguments x and y, and returns the result isequal(x, x) &amp;&amp; isequal(y, y).</td>
</tr>
<tr>
<td>int isordered (float x, float y)</td>
<td>int isordered (float n, float n)</td>
</tr>
<tr>
<td>int isordered (double x, double y)</td>
<td>int isunordered (float x, float y)</td>
</tr>
<tr>
<td>int isunordered (float n, float n)</td>
<td>int isunordered (double x, double y)</td>
</tr>
<tr>
<td>int isunordered (double n x, double n y)</td>
<td>int signbit (float)</td>
</tr>
<tr>
<td>int signbit (float n)</td>
<td>int signbit (double)</td>
</tr>
<tr>
<td>int signbit (double n)</td>
<td>long int signbit (double n)</td>
</tr>
<tr>
<td>int any (igentype x)</td>
<td>Returns 1 if the most significant bit in any component of x is set; otherwise returns 0.</td>
</tr>
<tr>
<td>int all (igentype x)</td>
<td>Returns 1 if the most significant bit in all components of x is set; otherwise returns 0.</td>
</tr>
<tr>
<td>gentype bitselect (gentype a, gentype b, gentype c)</td>
<td>Each bit of the result is the corresponding bit of a if the corresponding bit of c is 0. Otherwise it is the corresponding bit of b.</td>
</tr>
<tr>
<td>gentype select (gentype a, gentype b, igentype c)</td>
<td>For each component of a vector type, result[i] = if MSB of c[i] is set ? b[i] : a[i].</td>
</tr>
<tr>
<td>gentype select (gentype a, gentype b, ugentype c)</td>
<td>For a scalar type, result = c ? b : a.</td>
</tr>
</tbody>
</table>

igentype and ugentype must have the same number of elements and bits as gentype.

Table 6.14  Scalar and Vector Relational Functions
### 6.13.7 Vector Data Load and Store Functions

Table 6.15 describes the list of supported functions that allow you to read and write vector types from a pointer to memory. We use the generic type `gentype` to indicate the built-in data types `char, uchar, short, ushort, int, uint, long, ulong, float or double`. We use the generic type name `gentypen` to represent `n`-element vectors of `gentype` elements. We use the type name `halfn` to represent `n`-element vectors of `half` elements. The suffix `n` is also used in the function names (i.e. `vloadn, vstoren` etc.), where `n` = 2, 3, 4, 8 or 16.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>gentypen vloadn</code> (size_t <code>offset</code>, const <code>gentype</code> *<code>p</code>)</td>
<td>Return <code>sizeof (gentypen)</code> bytes of data read from address ((p + (\text{offset} \times n))). The address computed as ((p + (\text{offset} \times n))) must be 8-bit aligned if <code>gentype</code> is <code>char, uchar</code>; 16-bit aligned if <code>gentype</code> is <code>short, ushort</code>; 32-bit aligned if <code>gentype</code> is <code>int, uint, float</code>; 64-bit aligned if <code>gentype</code> is <code>long, ulong</code>.</td>
</tr>
<tr>
<td><code>gentypen vloadn</code> (size_t <code>offset</code>, const <code>constant gentype</code> *<code>p</code>)</td>
<td></td>
</tr>
<tr>
<td><code>void vstoren</code> (gentypen <code>data</code>, size_t <code>offset</code>, <code>gentype</code> *<code>p</code>)</td>
<td>Write <code>sizeof (gentypen)</code> bytes given by <code>data</code> to address ((p + (\text{offset} \times n))). The address computed as ((p + (\text{offset} \times n))) must be 8-bit aligned if <code>gentype</code> is <code>char, uchar</code>; 16-bit aligned if <code>gentype</code> is <code>short, ushort</code>; 32-bit aligned if <code>gentype</code> is <code>int, uint, float</code>; 64-bit aligned if <code>gentype</code> is <code>long, ulong</code>.</td>
</tr>
<tr>
<td><code>float vload\_half</code> (size_t <code>offset</code>, const <code>half</code> *<code>p</code>)</td>
<td>Read <code>sizeof (half)</code> bytes of data from address ((p + \text{offset})). The data read is interpreted as a half value. The half value is converted to a float value and the float value is returned. The read address computed as ((p + \text{offset})) must be 16-bit aligned.</td>
</tr>
<tr>
<td><code>float vload\_half</code> (size_t <code>offset</code>, const <code>constant half</code> *<code>p</code>)</td>
<td></td>
</tr>
<tr>
<td><code>floatn vload\_halfn</code> (size_t <code>offset</code>, const <code>half</code> *<code>p</code>)</td>
<td>Read <code>sizeof (halfn)</code> bytes of data from address ((p + (\text{offset} \times n))). The data read is interpreted as a halfn value. The halfn value read is converted to a floatn value and the floatn value is returned. The read address computed as ((p + (\text{offset} \times n))) must be 16-bit aligned.</td>
</tr>
<tr>
<td><code>floatn vload\_halfn</code> (size_t <code>offset</code>, const <code>constant half</code> *<code>p</code>)</td>
<td></td>
</tr>
</tbody>
</table>

---

37 The `halfn` type is only defined by the `cl_khr_fp16` extension described in section 9.5 of the OpenCL 1.2 Extension Specification.
<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>void vstore_half (float data, size_t offset, half *p)</code></td>
<td>The float value given by <code>data</code> is first converted to a half value using the appropriate rounding mode. The half value is then written to address computed as <code>(p + offset)</code>. The address computed as <code>(p + offset)</code> must be 16-bit aligned. <code>vstore_half</code> uses the default rounding mode. The default rounding mode is round to nearest even.</td>
</tr>
<tr>
<td><code>void vstore_half_rte (float data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_half_rtz (float data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_half_rtp (float data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_half_rtn (float data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_halfn (floatn data, size_t offset, half *p)</code></td>
<td>The floatn value given by <code>data</code> is converted to a halfn value using the appropriate rounding mode. The halfn value is then written to address computed as <code>(p + (offset * n))</code>. The address computed as <code>(p + (offset * n))</code> must be 16-bit aligned. <code>vstore_halfn</code> uses the default rounding mode. The default rounding mode is round to nearest even.</td>
</tr>
<tr>
<td><code>void vstore_halfn_rte (floatn data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_halfn_rtz (floatn data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_halfn_rtp (floatn data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_halfn_rtn (floatn data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_half (double data, size_t offset, half *p)</code></td>
<td>The double value given by <code>data</code> is first converted to a half value using the appropriate rounding mode. The half value is then written to address computed as <code>(p + offset)</code>. The address computed as <code>(p + offset)</code> must be 16-bit aligned. <code>vstore_half</code> uses the default rounding mode. The default rounding mode is round to nearest even.</td>
</tr>
<tr>
<td><code>void vstore_half_rte (double data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_half_rtz (double data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_half_rtp (double data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_half_rtn (double data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_halfn (doublen data, size_t offset, half *p)</code></td>
<td>The doublen value given by <code>data</code> is converted to a halfn value using the appropriate rounding mode. The halfn value is then written to address computed as <code>(p + (offset * n))</code>. The address computed as <code>(p + (offset * n))</code> must be 16-bit aligned. <code>vstore_halfn</code> uses the default rounding mode. The default rounding mode is round to nearest even.</td>
</tr>
<tr>
<td><code>void vstore_halfn_rte (doublen data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_halfn_rtz (doublen data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_halfn_rtp (doublen data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>void vstore_halfn_rtn (doublen data, size_t offset, half *p)</code></td>
<td></td>
</tr>
<tr>
<td><code>floatn vloada_halfn (size_t offset, ...)</code></td>
<td>For n = 2, 4, 8 and 16 read sizeof (halfn)</td>
</tr>
</tbody>
</table>
```
const half *p)
floatn vloada_halfn (size_t offset,
  const constant half *p)
bytes of data from address \((p + (\text{offset} * n))\). The data read is interpreted as a half\(n\) value. The half\(n\) value read is converted to a float\(n\) value and the float\(n\) value is returned.

The address computed as \((p + (\text{offset} * n))\) must be aligned to sizeof (half\(n\)) bytes.

For \(n = 3\), \textbf{vloada_half3} reads a half3 from address \((p + (\text{offset} * 4))\) and returns a float3. The address computed as \((p + (\text{offset} * 4))\) must be aligned to sizeof (half) * 4 bytes.

void vstorea_halfn (float \(n\text{data}\),
  size_t offset,
  half *p)
void vstorea_halfn_rte (float \(n\text{data}\),
  size_t offset,
  half *p)
void vstorea_halfn_rtz (float \(n\text{data}\),
  size_t offset,
  half *p)
void vstorea_halfn_rtp (float \(n\text{data}\),
  size_t offset,
  half *p)
void vstorea_halfn_rtn (float \(n\text{data}\),
  size_t offset,
  half *p)
The float\(n\) value given by \(\text{data}\) is converted to a half\(n\) value using the appropriate rounding mode.

For \(n = 2, 4, 8\) and 16, the half\(n\) value is written to the address computed as \((p + (\text{offset} * n))\). The address computed as \((p + (\text{offset} * n))\) must be aligned to sizeof (half\(n\)) bytes.

For \(n = 3\), the half3 value is written to the address computed as \((p + (\text{offset} * 4))\). The address computed as \((p + (\text{offset} * 4))\) must be aligned to sizeof (half) * 4 bytes.

\textbf{vstorea_halfn} uses the default rounding mode. The default rounding mode is round to nearest even.

void vstorea_halfn (double \(n\text{data}\),
  size_t offset,
  half *p)
void vstorea_halfn_rte (double \(n\text{data}\),
  size_t offset,
  half *p)
void vstorea_halfn_rtz (double \(n\text{data}\),
  size_t offset,
  half *p)
void vstorea_halfn_rtp (double \(n\text{data}\),
  size_t offset,
  half *p)
void vstorea_halfn_rtn (double \(n\text{data}\),
  size_t offset,
  half *p)
The double\(n\) value is converted to a half\(n\) value using the appropriate rounding mode.

For \(n = 2, 4, 8\) or 16, the half\(n\) value is written to the address computed as \((p + (\text{offset} * n))\). The address computed as \((p + (\text{offset} * n))\) must be aligned to sizeof (half\(n\)) bytes.

For \(n = 3\), the half3 value is written to the address computed as \((p + (\text{offset} * 4))\). The address computed as \((p + (\text{offset} * 4))\) must be aligned to sizeof (half) * 4 bytes.
```
vstorea_halfn uses the default rounding mode. The default rounding mode is round to nearest even.

<table>
<thead>
<tr>
<th>Table 6.15</th>
<th>Vector Data Load and Store Functions(^{38})</th>
</tr>
</thead>
<tbody>
<tr>
<td>The results of vector data load and store functions are undefined if the address being read from or written to is not correctly aligned as described in table 6.15. The pointer argument p can be a pointer to global, local or private memory for store functions described in table 6.15. The pointer argument p can be a pointer to global, local, constant or private memory for load functions described in table 6.15. NOTE: The vector data load and store functions variants that take pointer arguments which point to the generic address space are also supported.</td>
<td></td>
</tr>
</tbody>
</table>

\(^{38}\) vload\(^3\), and vload\(_{half}\)\(^3\) read x, y, z components from address \((p + (offset \times 3))\) into a 3-component vector. vstore\(^3\), and vstore\(_{half}\)\(^3\) write x, y, z components from a 3-component vector to address \((p + (offset \times 3))\). In addition vloada\(_{half}\)\(^3\) reads x, y, z components from address \((p + (offset \times 4))\) into a 3-component vector and vstorea\(_{half}\)\(^3\) writes x, y, z components from a 3-component vector to address \((p + (offset \times 4))\).
6.13.8 Synchronization Functions

The OpenCL C programming language implements the following synchronization functions.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>void <code>work_group_barrier</code> (cl_mem_fence_flags flags)</td>
<td>All work-items in a work-group executing the kernel on a processor must execute this function before any are allowed to continue execution beyond the <code>work_group_barrier</code>. This function must be encountered by all work-items in a work-group executing the kernel. These rules apply to ND-ranges implemented with uniform and non-uniform work-groups.</td>
</tr>
<tr>
<td>void <code>work_group_barrier</code> (cl_mem_fence_flags flags, memory_scope scope)</td>
<td>If <code>work_group_barrier</code> is inside a conditional statement, then all work-items must enter the conditional if any work-item enters the conditional statement and executes the <code>work_group_barrier</code>.</td>
</tr>
</tbody>
</table>

If `work_group_barrier` is inside a loop, all work-items must execute the `work_group_barrier` for each iteration of the loop before any are allowed to continue execution beyond the `work_group_barrier`.

The `work_group_barrier` function also supports a variant that specifies the memory scope. For the `work_group_barrier` variant that does not take a memory scope, the `scope` is `memory_scope_work_group`.

The `scope` argument specifies whether the memory accesses of work-items in the work-group to memory address space(s) identified by `flags` become visible to all work-items in the work-group, the device or all SVM devices.

The `work_group_barrier` function can also be used to specify which memory operations i.e. to global memory, local memory or images become visible to

---

39 The built-in function `barrier` has been renamed `work_group_barrier`. For backward compatibility, `barrier` is also supported.

40 Refer to section 6.13.11 for description of memory_scope.
the appropriate memory scope identified by \textit{scope}. The \textit{flags} argument specifies the memory address spaces. This is a bitfield and can be set to 0 or a combination of the following values ORed together. When these flags are OR’ed together the \textbf{work\_group\_ barrier} acts as a combined barrier for all address spaces specified by the flags ordering memory accesses both within and across the specified address spaces.

\textbf{CLK\_LOCAL\_MEM\_FENCE} - The \textbf{work\_group\_ barrier} function will ensure that all local memory accesses become visible to all work-items in the work-group. Note that the value of \textit{scope} is ignored as the memory scope is always \texttt{memory\_scope\_work\_group}.

\textbf{CLK\_GLOBAL\_MEM\_FENCE} – The \textbf{work\_group\_ barrier} function ensure that all global memory accesses become visible to the appropriate scope as given by \textit{scope}.

\textbf{CLK\_IMAGE\_MEM\_FENCE} – The \textbf{work\_group\_ barrier} function will ensure that all image memory accesses become visible to the appropriate scope as given by \textit{scope}. The value of \textit{scope} must be \texttt{memory\_scope\_work\_group} or \texttt{memory\_scope\_device}.

The values of \textit{flags} and \textit{scope} must be the same for all work-items in the work-group.

\begin{table}[h]
\centering
\begin{tabular}{|c|c|}
\hline
\textbf{Table 6.1} & \textit{Built-in Synchronization Functions} \\
\hline
\end{tabular}
\end{table}
### 6.13.9 Address Space Qualifier Functions

The OpenCL C programming language implements the following address space qualifier functions. We use the generic type name gentype to indicate any of the built-in data types supported by OpenCL C or a user defined type.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>global gentype *</td>
<td>Returns a pointer that points to a region in the global address space if to_global can cast ptr to the global address space. Otherwise it returns NULL.</td>
</tr>
<tr>
<td>to_global (gentype *ptr)</td>
<td></td>
</tr>
<tr>
<td>const global gentype *</td>
<td></td>
</tr>
<tr>
<td>to_global (const gentype *ptr)</td>
<td></td>
</tr>
<tr>
<td>local gentype *</td>
<td>Returns a pointer that points to a region in the local address space if to_local can cast ptr to the local address space. Otherwise it returns NULL.</td>
</tr>
<tr>
<td>to_local (gentype *ptr)</td>
<td></td>
</tr>
<tr>
<td>const local gentype *</td>
<td></td>
</tr>
<tr>
<td>to_local (const gentype *ptr)</td>
<td></td>
</tr>
<tr>
<td>private gentype *</td>
<td>Returns a pointer that points to a region in the private address space if to_private can cast ptr to the private address space. Otherwise it returns NULL.</td>
</tr>
<tr>
<td>to_private (gentype *ptr)</td>
<td></td>
</tr>
<tr>
<td>const private gentype *</td>
<td></td>
</tr>
<tr>
<td>to_private (const gentype *ptr)</td>
<td></td>
</tr>
<tr>
<td>cl_mem_fence_flags</td>
<td>Returns a valid memory fence value for ptr.</td>
</tr>
<tr>
<td>get_fence (gentype *ptr)</td>
<td></td>
</tr>
<tr>
<td>cl_mem_fence_flags</td>
<td></td>
</tr>
<tr>
<td>get_fence (const gentype *ptr)</td>
<td></td>
</tr>
</tbody>
</table>

**Table 6.17  Built-in Address Space Qualifier Functions**
Async Copies from Global to Local Memory, Local to Global Memory, and Prefetch

The OpenCL C programming language implements the following functions that provide asynchronous copies between global and local memory and a prefetch from global memory.

We use the generic type name `gentype` to indicate the built-in data types `char, char[2|3|4|8|16], uchar, uchar[2|3|4|8|16], short, short[2|3|4|8|16], ushort, ushort[2|3|4|8|16], int, int[2|3|4|8|16], uint, uint[2|3|4|8|16], long, long[2|3|4|8|16], ulong, ulong[2|3|4|8|16], float, float[2|3|4|8|16], or double, double[2|3|4|8|16]` as the type for the arguments unless otherwise stated.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>event_t async_work_group_copy ( __local gentype *dst, const __global gentype *src, size_t num_gentypes, event_t event)</code></td>
<td>Perform an async copy of <code>num_gentypes</code> gentype elements from <code>src</code> to <code>dst</code>. The async copy is performed by all work-items in a work-group and this built-in function must therefore be encountered by all work-items in a work-group executing the kernel with the same argument values; otherwise the results are undefined. This rule applies to ND-ranges implemented with uniform and non-uniform work-groups. Returns an event object that can be used by <code>wait_group_events</code> to wait for the async copy to finish. The <code>event</code> argument can also be used to associate the <code>async_work_group_copy</code> with a previous async copy allowing an event to be shared by multiple async copies; otherwise <code>event</code> should be zero. If <code>event</code> argument is non-zero, the event object supplied in <code>event</code> argument will be returned. This function does not perform any implicit synchronization of source data</td>
</tr>
</tbody>
</table>

---

41 `async_work_group_copy` and `async_work_group_strided_copy` for 3-component vector types behave as `async_work_group_copy` and `async_work_group_strided_copy` respectively for 4-component vector types.
such as using a **barrier** before performing the copy.

<table>
<thead>
<tr>
<th>Declaration</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>event_t async_work_group_strided_copy ( _local gentype *dst, const _global gentype *src, size_t num_gentypes, size_t src_stride, event_t event)</code></td>
<td>Perform an async gather of <code>num_gentypes</code> <code>gentype</code> elements from <code>src</code> to <code>dst</code>. The <code>src_stride</code> is the stride in elements for each <code>gentype</code> element read from <code>src</code>. The <code>dst_stride</code> is the stride in elements for each <code>gentype</code> element written to <code>dst</code>. The async gather is performed by all work-items in a work-group. This built-in function must therefore be encountered by all work-items in a work-group executing the kernel with the same argument values; otherwise the results are undefined. This rule applies to ND-ranges implemented with uniform and non-uniform work-groups. Returns an event object that can be used by <code>wait_group_events</code> to wait for the async copy to finish. The <code>event</code> argument can also be used to associate the <code>async_work_group_strided_copy</code> with a previous async copy allowing an event to be shared by multiple async copies; otherwise <code>event</code> should be zero. If <code>event</code> argument is non-zero, the event object supplied in <code>event</code> argument will be returned. This function does not perform any implicit synchronization of source data such as using a <strong>barrier</strong> before performing the copy. The behavior of <code>async_work_group_strided_copy</code> is undefined if <code>src_stride</code> or <code>dst_stride</code> is 0, or if the <code>src_stride</code> or <code>dst_stride</code> values cause the <code>src</code> or <code>dst</code> pointers to exceed the upper bounds of the address space during the copy.</td>
</tr>
</tbody>
</table>
| `void wait_group_events (int num_events, event_t *event_list)` | Wait for events that identify the `async_work_group_copy` operations to
complete. The event objects specified in \textit{event\_list} will be released after the wait is performed.

This function must be encountered by all work-items in a work-group executing the kernel with the same \textit{num\_events} and event objects specified in \textit{event\_list}; otherwise the results are undefined. This rule applies to ND-ranges implemented with uniform and non-uniform work-groups.

\begin{table}
\centering
\begin{tabular}{|l|l|}
\hline
\textbf{void \texttt{prefetch}} (const \_global gen\texttt{type} *\texttt{p}, \\
size_t \texttt{num\_gentypes}) & Prefetch \texttt{num\_gentypes}\,* \\
& sizeof(gen\texttt{type}) bytes into the global cache. The prefetch instruction is applied \\
& to a work-item in a work-group and does not affect the functional behavior of the \\
& kernel. \\
\hline
\end{tabular}
\end{table}

\textbf{Table 6.18} \quad \textit{Built-in Async Copy and Prefetch Functions}

NOTE: The kernel must wait for the completion of all async copies using the \texttt{wait\_group\_events} built-in function before exiting; otherwise the behavior is undefined.
6.13.11 Atomic Functions

The OpenCL C programming language implements a subset of the C11 atomics (refer to section 7.17 of the C11 specification) and synchronization operations. These operations play a special role in making assignments in one work-item visible to another. A synchronization operation on one or more memory locations is either an acquire operation, a release operation, or both an acquire and release operation\textsuperscript{42}. A synchronization operation without an associated memory location is a fence and can be either an acquire fence, a release fence or both an acquire and release fence. In addition, there are relaxed atomic operations, which are not synchronization operations, and atomic read-modify-write operations which have special characteristics.

The types include

\begin{verbatim}
memory_order
\end{verbatim}

which is an enumerated type whose enumerators identify memory ordering constraints;

\begin{verbatim}
memory_scope
\end{verbatim}

which is an enumerated type whose enumerators identify scope of memory ordering constraints;

\begin{verbatim}
atomic_flag
\end{verbatim}

which is a 32-bit integer type representing a lock-free, primitive atomic flag; and several atomic analogs of integer types.

In the following operation definitions:

- An A refers to one of the atomic types.
- A C refers to its corresponding non-atomic type.
- An M refers to the type of the other argument for arithmetic operations. For atomic integer types, M is C.
- The functions not ending in explicit have the same semantics as the corresponding explicit function with memory_order_seq_cst for the memory_order argument.
- The functions that do not have memory_scope argument have the same semantics as the corresponding functions with the memory_scope argument set to memory_scope_device.

\textsuperscript{42} The C11 consume operation is not supported.
NOTE: With fine-grained system SVM, sharing happens at the granularity of individual loads and stores anywhere in host memory. Memory consistency is always guaranteed at synchronization points, but to obtain finer control over consistency, the OpenCL atomics functions may be used to ensure that the updates to individual data values made by one unit of execution are visible to other execution units. In particular, when a host thread needs fine control over the consistency of memory that is shared with one or more OpenCL devices, it must use atomic and fence operations that are compatible with the C11 atomic operations\textsuperscript{43}.

6.13.11.1 The ATOMIC_VAR_INIT macro

The ATOMIC_VAR_INIT macro expands to a token sequence suitable for initializing an atomic object of a type that is initialization-compatible with value. An atomic object with automatic storage duration that is not explicitly initialized using ATOMIC_VAR_INIT is initially in an indeterminate state; however, the default (zero) initialization for objects with static storage duration is guaranteed to produce a valid state.

\[
\text{#define ATOMIC_VAR_INIT(C \ value)}
\]

This macro can only be used to initialize atomic objects that are declared in program scope in the global address space.

Examples:

\[
\text{global atomic_int guide = ATOMIC_VAR_INIT(42);} \\
\]

Concurrent access to the variable being initialized, even via an atomic operation, constitutes a data-race.

6.13.11.2 The atomic_init function

The atomic_init function non-atomically initializes the atomic object pointed to by obj to the value value.

\[
\text{void atomic_init(volatile A *obj, C value)}
\]

Examples:

\[
\text{local atomic_int local_guide;}
\text{if (get_local_id(0) == 0)}
\text{atomic_init(&guide, 42);} \\
\]

\textsuperscript{43} We can’t require C11 atomics since host programs can be implemented in other programming languages and versions of C or C++, but we do require that the host programs use atomics and that those atomics be compatible with those in C11.
work_group_barrier(CLK_LOCAL_MEM_FENCE);

6.13.11.3 Order and Consistency

The enumerated type memory_order specifies the detailed regular (non-atomic) memory synchronization operations as defined in section 5.1.2.4 of the C11 specification and may provide for operation ordering. Its enumeration constants are as follows:

    memory_order_relaxed
    memory_order_acquire
    memory_order_release
    memory_order_acq_rel
    memory_order_seq_cst

The memory_order can be used when performing atomic operations to global or local memory.

6.13.11.4 Memory Scope

The enumerated type memory_scope specifies whether the memory ordering constraints given by memory_order apply to work-items in a work-group or work-items of a kernel(s) executing on the device or across devices (in the case of shared virtual memory). Its enumeration constants are as follows:

    memory_scope_work_item
    memory_scope_work_group
    memory_scope_device
    memory_scope_all_svm_devices

The memory scope should only be used when performing atomic operations to global memory. Atomic operations to local memory only guarantee memory ordering in the work-group not across work-groups and therefore ignore the memory_scope value.

6.13.11.5 Fences

The following fence operations are supported.

6.13.11.5.1 The atomic_work_item_fence function

    void atomic_work_item_fence(cl_mem_fence_flags flags, flags,
memory_order order,
memory_scope scope)

flags must be set to CLK_GLOBAL_MEM_FENCE, CLK_LOCAL_MEM_FENCE,
CLK_IMAGE_MEM_FENCE or a combination of these values ORed together; otherwise the
behavior is undefined. The behavior of calling atomic_work_item_fence with
CLK_GLOBAL_MEM_FENCE, CLK_LOCAL_MEM_FENCE or CLK_IMAGE_MEM_FENCE
ORed together is equivalent to calling atomic_work_item_fence individually for each of
the fence values set in flags.

Depending on the value of order, this operation:

- has no effects, if order == memory_order_relaxed.
- is an acquire fence, if order == memory_order_acquire.
- is a release fence, if order == memory_order_release.
- is both an acquire fence and a release fence, if order == memory_order_acq_rel.
- is a sequentially consistent acquire and release fence, if order ==
  memory_order_seq_cst.

For images declared with the read_write qualifier, the atomic_work_item_fence
must be called to make sure that writes to the image by a work-item become visible to that work-
item on subsequent reads to that image by that work-item.

6.13.11.6 Atomic integer and floating-point types

The list of supported atomic type names are:

- atomic_int
- atomic_uint
- atomic_long\(^{45}\)
- atomic_ulong
- atomic_float
- atomic_double\(^{46}\)
- atomicintptr_t\(^{47}\)
- atomicuintptr_t
- atomic_size_t
- atomicptrdiff_t

\(^{45}\) The atomic_long and atomic_ulong types are supported if the \texttt{cl\_khr\_int64\_base\_atomics} and
\texttt{cl\_khr\_int64\_extended\_atomics} extensions are supported.

\(^{46}\) The atomic_double type is only supported if double precision is supported and the
\texttt{cl\_khr\_int64\_base\_atomics} and \texttt{cl\_khr\_int64\_extended\_atomics} extensions are supported.

\(^{47}\) If the device address space is 64-bits, the data types atomic intptr_t, atomicuintptr_t,
atomic_size_t and atomicptrdiff_t are supported if the \texttt{cl\_khr\_int64\_base\_atomics} and
\texttt{cl\_khr\_int64\_extended\_atomics} extensions are supported.
Arguments to a kernel can be declared to be a pointer to the above atomic types or the 
atomic_flag type.

The representation of atomic integer, floating-point and pointer types have the same size as their 
corresponding regular types. The atomic_flag type must be implemented as a 32-bit integer.

6.13.11.7 Operations on atomic types

There are only a few kinds of operations on atomic types, though there are many instances of 
those kinds. This section specifies each general kind.

6.13.11.7.1 The atomic_store functions

    void atomic_store(volatile A *object, C desired)

    void atomic_store_explicit(volatile A *object,
                          C desired,
                          memory_order order)

    void atomic_store_explicit(volatile A *object,
                          C desired,
                          memory_order order,
                          memory_scope scope)

The order argument shall not be memory_order_acquire, nor 
memory_order_acq_rel. Atomically replace the value pointed to by object with the 
value of desired. Memory is affected according to the value of order.

6.13.11.7.2 The atomic_load functions

    C atomic_load(volatile A *object)

    C atomic_load_explicit(volatile A *object,
                          memory_order order)

    C atomic_load_explicit(volatile A *object,
                          memory_order order,
                          memory_scope scope)

The order argument shall not be memory_order_release nor 
memory_order_acq_rel. Memory is affected according to the value of order.
Atomically returns the value pointed to by object.

### 6.13.11.7.3 The atomic_exchange functions

```c
C atomic_exchange(volatile A *object, C desired)
```

```c
C atomic_exchange_explicit(volatile A *object,
    C desired,
    memory_order order)
```

```c
C atomic_exchange_explicit(volatile A *object,
    C desired,
    memory_order order,
    memory_scope scope)
```

Atomically replace the value pointed to by object with desired. Memory is affected according to the value of `order`. These operations are read-modify-write operations (as defined by section 5.1.2.4 of the C11 specification). Atomically returns the value pointed to by object immediately before the effects.

### 6.13.11.7.4 The atomic_compare_exchange functions

```c
bool atomic_compare_exchange_strong(volatile A *object,
    C *expected, C desired)
```

```c
bool atomic_compare_exchange_strong_explicit(
    volatile A *object,
    C *expected,
    C desired,
    memory_order success,
    memory_order failure)
```

```c
bool atomic_compare_exchange_strong_explicit(
    volatile A *object,
    C *expected,
    C desired,
    memory_order success,
    memory_order failure,
    memory_scope scope)
```

```c
bool atomic_compare_exchange_weak(volatile A *object,
    C *expected, C desired)
```

```c
bool atomic_compare_exchange_weak_explicit(
    volatile A *object,
    C *expected,
    C desired,
    memory_scope scope)
```

bool atomic_compare_exchange_weak_explicit(  
  volatile A *object,  
  C *expected,  
  C desired,  
  memory_order success,  
  memory_order failure,  
  memory_scope scope)

The failure argument shall not be memory_order_release nor memory_order_acq_rel. The failure argument shall be no stronger than the success argument. Atomically, compares the value pointed to by object for equality with that in expected, and if true, replaces the value pointed to by object with desired, and if false, updates the value in expected with the value pointed to by object. Further, if the comparison is true, memory is affected according to the value of success, and if the comparison is false, memory is affected according to the value of failure. These operations are atomic read-modify-write operations (as defined by section 5.1.2.4 of the C11 specification).

NOTE: The effect of the compare-and-exchange operations is

\[
\begin{align*}
\text{if} \ (\text{memcmp}(\text{object}, \text{expected}, \text{sizeof}(*\text{object}) == 0) \\
\text{memcpy}(\text{object}, &\text{desired}, \text{sizeof}(*\text{object})); \ \\
\text{else} \ \\
\text{memcpy}(\text{expected}, \text{object}, \text{sizeof}(*\text{object}));
\end{align*}
\]

The weak compare-and-exchange operations may fail spuriously. That is, even when the contents of memory referred to by expected and object are equal, it may return zero and store back to expected the same memory contents that were originally there.

These generic functions return the result of the comparison.

6.13.11.7.5 The atomic_fetch and modify functions

The following operations perform arithmetic and bitwise computations. All of these operations are applicable to an object of any atomic integer type. The key, operator, and computation correspondence is given in table below:

<table>
<thead>
<tr>
<th>key</th>
<th>op</th>
<th>computation</th>
</tr>
</thead>
<tbody>
<tr>
<td>add</td>
<td>+</td>
<td>addition</td>
</tr>
</tbody>
</table>

This spurious failure enables implementation of compare-and-exchange on a broader class of machines, e.g. load-locked store-conditional machines.
<table>
<thead>
<tr>
<th>sub</th>
<th>-</th>
<th>subtraction</th>
</tr>
</thead>
<tbody>
<tr>
<td>or</td>
<td></td>
<td>bitwise inclusive or</td>
</tr>
<tr>
<td>xor</td>
<td>^</td>
<td>bitwise exclusive or</td>
</tr>
<tr>
<td>and</td>
<td>&amp;</td>
<td>bitwise and</td>
</tr>
<tr>
<td>min</td>
<td>min</td>
<td>compute min</td>
</tr>
<tr>
<td>max</td>
<td>max</td>
<td>compute max</td>
</tr>
</tbody>
</table>

NOTE: For atomic_fetch and modify functions with key = add or sub on atomic types `atomic_intptr_t` and `atomic_uintptr_t`, M is `ptrdiff_t`. For atomic_fetch and modify functions with key = or, xor, and, min and max on atomic types `atomic_intptr_t` and `atomic_uintptr_t`, M is `intptr_t` and `uintptr_t`.

```c
C atomic_fetch_key(volatile A *object, M operand)
C atomic_fetch_key_explicit(volatile A *object, M operand, memory_order order)
C atomic_fetch_key_explicit(volatile A *object, M operand, memory_order order, memory_scope scope)
```

Atomically replaces the value pointed to by `object` with the result of the computation applied to the value pointed to by `object` and the given operand. Memory is affected according to the value of `order`. These operations are atomic read-modify-write operations (as defined by section 5.1.2.4 of the C11 specification). For signed integer types, arithmetic is defined to use two’s complement representation with silent wrap-around on overflow; there are no undefined results. For address types, the result may be an undefined address, but the operations otherwise have no undefined behavior. Returns atomically, the value pointed to by `object` immediately before the effects.

6.13.11.7.6 Atomic flag type and operations

The `atomic_flag` type provides the classic test-and-set functionality. It has two states, set (value is non-zero) and clear (value is 0). Operations on an object of type `atomic_flag` shall be lock free.

The macro `ATOMIC_FLAG_INIT` may be used to initialize an `atomic_flag` to the clear state. An `atomic_flag` that is not explicitly initialized with `ATOMIC_FLAG_INIT` is initially in an indeterminate state.

This macro can only be used for atomic objects that are declared in program scope in the global address space with the `atomic_flag` type.
Example:

```c
global atomic_flag guard = ATOMIC_FLAG_INIT;
```

### 6.13.11.7.7 The atomic_flag_test_and_set functions

```c
bool atomic_flag_test_and_set(
    volatile atomic_flag *object)

bool atomic_flag_test_and_set_explicit(
    volatile atomic_flag *object,
    memory_order order)

bool atomic_flag_test_and_set_explicit(
    volatile atomic_flag *object,
    memory_order order,
    memory_scope scope)
```

Atomically sets the value pointed to by `object` to true. Memory is affected according to the value of `order`. These operations are atomic read-modify-write operations (as defined by section 5.1.2.4 of the C11 specification). Returns atomically, the value of the object immediately before the effects.

### 6.13.11.7.8 The atomic_flag_clear functions

```c
void atomic_flag_clear(volatile atomic_flag *object)

void atomic_flag_clear_explicit(
    volatile atomic_flag *object,
    memory_order order)

void atomic_flag_clear_explicit(
    volatile atomic_flag *object,
    memory_order order,
    memory_scope scope)
```

The `order` argument shall not be `memory_order_acquire` nor `memory_order_acq_rel`. Atomically sets the value pointed to by `object` to false. Memory is affected according to the value of `order`.

### 6.13.11.8 Restrictions

> All operations on atomic types must be performed using the built-in atomic functions. C11 and C++11 support operators on atomic types. OpenCL C does not support operators with atomic types. Using atomic types with operators should result in a
compilation error.

- The `atomic_bool, atomic_char, atomic_uchar, atomic_short, atomic_ushort, atomic_intmax_t` and `atomic_uintmax_t` types are not supported by OpenCL C.

- OpenCL C requires that the built-in atomic functions on atomic types are lock-free.

- The `_Atomic type specifier and _Atomic type qualifier are not supported by OpenCL C.`

- The behavior of atomic operations where pointer arguments to the atomic functions refers to an atomic type in the private address space is undefined.
## 6.13.12 Miscellaneous Vector Functions

The OpenCL C programming language implements the following additional built-in vector functions. We use the generic type name `gentypen` (or `gentypem`) to indicate the built-in data types `char{2|4|8|16}`, `uchar{2|4|8|16}`, `short{2|4|8|16}`, `ushort{2|4|8|16}`, `half{2|4|8|16}`, `int{2|4|8|16}`, `uint{2|4|8|16}`, `long{2|4|8|16}`, `ulong{2|4|8|16}`, `float{2|4|8|16}` or `double{2|4|8|16}`\(^{49}\) as the type for the arguments unless otherwise stated. We use the generic name `ugentypen` to indicate the built-in unsigned integer data types.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>int vec_step (gentypen a)</code></td>
<td>The <code>vec_step</code> built-in function takes a built-in scalar or vector data type argument and returns an integer value representing the number of elements in the scalar or vector. For all scalar types, <code>vec_step</code> returns 1. The <code>vec_step</code> built-in functions that take a 3-component vector return 4. <code>vec_step</code> may also take a pure type as an argument, e.g. <code>vec_step(float2)</code></td>
</tr>
<tr>
<td><code>gentypen shuffle (gentypen x, ugentypen mask)</code></td>
<td>The <code>shuffle</code> and <code>shuffle2</code> built-in functions construct a permutation of elements from one or two input vectors respectively that are of the same type, returning a vector with the same element type as the input and length that is the same as the shuffle mask. The size of each element in the <code>mask</code> must match the size of each element in the result. For <code>shuffle</code>, only the <code>ilogb(2m-1)</code> least significant bits of each <code>mask</code> element are considered. For <code>shuffle2</code>, only the <code>ilogb(2m-1)+1</code> least significant bits of each <code>mask</code> element are considered. Other bits in the mask shall be ignored. The elements of the input vectors are numbered from left to right across one or both of the vectors.</td>
</tr>
<tr>
<td><code>gentypen shuffle2 (gentypen x, gentypen y, ugentypen mask)</code></td>
<td></td>
</tr>
</tbody>
</table>

\(^{49}\) Only if the `cl_khr_fp16` extension is supported.  
\(^{50}\) Only if double precision is supported.
For this purpose, the number of elements in a vector is given by \texttt{vec\_step(gentype m)}. The shuffle \textit{mask} operand specifies, for each element of the result vector, which element of the one or two input vectors the result element gets.

Examples:

```c
uint4 mask = (uint4)(3, 2, 1, 0);
float4 a;
float4 r = shuffle(a, mask);
// r.s0123 = a.wzyx

uint8 mask = (uint8)
(0, 1, 2, 3,
 4, 5, 6, 7);
float4 a, b;
float8 r = shuffle2(a, b, mask);
// r.s0123 = a.xyzw
// r.s4567 = b.xyzw

uint4 mask;
float8 a;
float4 b;
b = shuffle(a, mask);
```

Examples that are not valid are:

```c
uint8 mask;
short16 a;
short8 b;
b = shuffle(a, mask);  \leftarrow \text{not valid}
```

| Table 6.20 | Built-in Miscellaneous Vector Functions |
6.13.13 `printf`

The OpenCL C programming language implements the `printf` function.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>int printf(constant char * restrict format, …)</code></td>
<td>The <code>printf</code> built-in function writes output to an implementation-defined stream such as stdout under control of the string pointed to by <code>format</code> that specifies how subsequent arguments are converted for output. If there are insufficient arguments for the format, the behavior is undefined. If the format is exhausted while arguments remain, the excess arguments are evaluated (as always) but are otherwise ignored. The <code>printf</code> function returns when the end of the format string is encountered. <code>printf</code> returns 0 if it was executed successfully and -1 otherwise.</td>
</tr>
</tbody>
</table>

Table 6.21  Built-in `printf` Function

6.13.13.1 `printf` output synchronization

When the event that is associated with a particular kernel invocation is completed, the output of all `printf()` calls executed by this kernel invocation is flushed to the implementation-defined output stream. Calling `clFinish` on a command queue flushes all pending output by `printf` in previously enqueued and completed commands to the implementation-defined output stream. In the case that `printf` is executed from multiple work-items concurrently, there is no guarantee of ordering with respect to written data. For example, it is valid for the output of a work-item with a global id (0,0,1) to appear intermixed with the output of a work-item with a global id (0,0,4) and so on.

6.13.13.2 `printf` format string

The format shall be a character sequence, beginning and ending in its initial shift state. The format is composed of zero or more directives: ordinary characters (not %), which are copied unchanged to the output stream; and conversion specifications, each of which results in fetching zero or more subsequent arguments, converting them, if applicable, according to the corresponding conversion specifier, and then writing the result to the output stream. The format is in the `constant` address space and must be resolvable at compile time i.e. cannot be dynamically created by the executing program, itself.
Each conversion specification is introduced by the character \%. After the \%, the following appear in sequence:

- **Zero or more flags** (in any order) that modify the meaning of the conversion specification.

- An optional minimum **field width**. If the converted value has fewer characters than the field width, it is padded with spaces (by default) on the left (or right, if the left adjustment flag, described later, has been given) to the field width. The field width takes the form of a nonnegative decimal integer.\(^{51}\)

- An optional **precision** that gives the minimum number of digits to appear for the \d, i, o, u, x, and X conversions, the number of digits to appear after the decimal-point character for a, A, e, E, f, and F conversions, the maximum number of significant digits for the g and G conversions, or the maximum number of bytes to be written for s conversions. The precision takes the form of a period (.) followed by an optional decimal integer; if only the period is specified, the precision is taken as zero. If a precision appears with any other conversion specifier, the behavior is undefined.

- An optional **vector specifier**.

- A **length modifier** that specifies the size of the argument. The length modifier is required with a vector specifier and together specifies the vector type. Implicit conversions between vector types are disallowed (as per section 6.2.1). If the vector specifier is not specified, the length modifier is optional.

- A **conversion specifier** character that specifies the type of conversion to be applied.

The flag characters and their meanings are:

- **-** The result of the conversion is left-justified within the field. (It is right-justified if this flag is not specified.)

- **+** The result of a signed conversion always begins with a plus or minus sign. (It begins with a sign only when a negative value is converted if this flag is not specified.)\(^{52}\)

- **space** If the first character of a signed conversion is not a sign, or if a signed conversion results in no characters, a space is prefixed to the result. If the space and + flags both appear, the space flag is ignored.

- **#** The result is converted to an “alternative form”. For o conversion, it increases the precision, if and only if necessary, to force the first digit of the result to be a zero (if the value and precision are both 0, a single 0 is printed). For x (or X) conversion, a nonzero value...
result has \(0x\) (or \(0X\)) prefixed to it. For \(a, A, e, E, f, F, g,\) and \(G\) conversions, the result of converting a floating-point number always contains a decimal-point character, even if no digits follow it. (Normally, a decimal-point character appears in the result of these conversions only if a digit follows it.) For \(g\) and \(G\) conversions, trailing zeros are not removed from the result. For other conversions, the behavior is undefined.

0

For \(d, i, o, u, x, X, a, A, e, E, f, F, g,\) and \(G\) conversions, leading zeros (following any indication of sign or base) are used to pad to the field width rather than performing space padding, except when converting an infinity or NaN. If the 0 and - flags both appear, the 0 flag is ignored. For \(d, i, o, u, x,\) and \(X\) conversions, if a precision is specified, the 0 flag is ignored. For other conversions, the behavior is undefined.

The vector specifier and its meaning is:

\[vn\]

Specifies that a following \(a, A, e, E, f, F, g, d, i, o, u, x,\) or \(X\) conversion specifier applies to a vector argument, where \(n\) is the size of the vector and must be 2, 3, 4, 8 or 16.

The vector value is displayed in the following general form:

\[\text{value1} \ C \ text{value2} \ C \ldots \ C \ text{valuen}\]

where \(C\) is a separator character. The value for this separator character is a comma.

If the vector specifier is not used, the length modifiers and their meanings are:

\[hh\]

Specifies that a following \(d, i, o, u, x,\) or \(X\) conversion specifier applies to a \text{char} or \text{uchar} argument (the argument will have been promoted according to the integer promotions, but its value shall be converted to \text{char} or \text{uchar} before printing).

\[h\]

Specifies that a following \(d, i, o, u, x,\) or \(X\) conversion specifier applies to a \text{short} or \text{ushort} argument (the argument will have been promoted according to the integer promotions, but its value shall be converted to \text{short} or \text{unsigned short} before printing).

\[(ell)\]

Specifies that a following \(d, i, o, u, x,\) or \(X\) conversion specifier applies to a \text{long} or \text{ulong} argument. The \(l\) modifier is supported by the full profile. For the embedded profile, the \(l\) modifier is supported only if 64-bit integers are supported by the device.

If the vector specifier is used, the length modifiers and their meanings are:

\[hh\]

Specifies that a following \(d, i, o, u, x,\) or \(X\) conversion specifier applies to a \text{char} or \text{uchar} argument (the argument will not be promoted).

\[h\]

Specifies that a following \(d, i, o, u, x,\) or \(X\) conversion specifier applies to a \text{short} or \text{ushort} argument (the argument will not be promoted); that
a following \texttt{a, A, e, E, f, F, g, or G} conversion specifier applies to a \texttt{halfn} \footnote{\textit{Only if the cl\_khr\_fp16} extension is supported.} argument.

\textbf{hl} \hspace{0.3cm} This modifier can only be used with the vector specifier. Specifies that a following \texttt{d, i, o, u, x, or X} conversion specifier applies to a \texttt{intn} or \texttt{uintn} argument; that a following \texttt{a, A, e, E, f, F, g, or G} conversion specifier applies to a \texttt{floatn} argument.

\textbf{l (ell)} \hspace{0.3cm} Specifies that a following \texttt{d, i, o, u, x, or X} conversion specifier applies to a \texttt{longn} or \texttt{ulongn} argument; that a following \texttt{a, A, e, E, f, F, g, or G} conversion specifier applies to a \texttt{doublen} argument. The \textbf{l} modifier is supported by the full profile. For the embedded profile, the \textbf{l} modifier is supported only if 64-bit integers or double-precision floating-point are supported by the device.

If a vector specifier appears without a length modifier, the behavior is undefined. The vector data type described by the vector specifier and length modifier must match the data type of the argument; otherwise the behavior is undefined.

If a length modifier appears with any conversion specifier other than as specified above, the behavior is undefined.

The conversion specifiers and their meanings are:

\textbf{d,i} \hspace{0.3cm} The \texttt{int, charn, shortn, intn} or \texttt{longn} argument is converted to signed decimal in the style [−]dddd. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it is expanded with leading zeros. The default precision is 1. The result of converting a zero value with a precision of zero is no characters.

\textbf{o,u, x,X} \hspace{0.3cm} The \texttt{unsigned int, ucharn, ushortn, uintn} or \texttt{ulongn} argument is converted to unsigned octal (\texttt{o}), unsigned decimal (\texttt{u}), or unsigned hexadecimal notation (\texttt{x} or \texttt{X}) in the style \texttt{dddd}; the letters \texttt{abcdef} are used for \texttt{x} conversion and the letters \texttt{ABCDEF} for \texttt{X} conversion. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it is expanded with leading zeros. The default precision is 1. The result of converting a zero value with a precision of zero is no characters.

\textbf{f,F} \hspace{0.3cm} A \texttt{double, halfn, floatn} or \texttt{doublen} argument representing a floating-point number is converted to decimal notation in the style [−]ddd.ddd, where the number of digits after the decimal-point character is equal to the precision specification. If the precision is missing, it is taken as 6; if the precision is zero and the \# flag is not specified, no decimal-point character appears. If a decimal-point character appears, at least one digit appears before it. The value is rounded to the appropriate number of digits. A \texttt{double, halfn, floatn} or \texttt{doublen} argument representing an infinity is converted in one of the styles [−]\texttt{inf} or [−]\texttt{infinity} — which style is implementation-defined. A \texttt{double, halfn},
floatn or doublen argument representing a NaN is converted in one of the styles [-]nan or [-]nan(n-char-sequence)—which style, and the meaning of any n-char-sequence, is implementation-defined. The F conversion specifier produces INF, INFINITY, or NaN instead of inf, infinity, or nan, respectively.\(^{54}\)

e,E

A double, halfn, floatn or doublen argument representing a floating-point number is converted in the style [-]d.ddd e±dd, where there is one digit (which is nonzero if the argument is nonzero) before the decimal-point character and the number of digits after it is equal to the precision; if the precision is missing, it is taken as 6; if the precision is zero and the # flag is not specified, no decimal-point character appears. The value is rounded to the appropriate number of digits. The E conversion specifier produces a number with E instead of e introducing the exponent. The exponent always contains at least two digits, and only as many more digits as necessary to represent the exponent. If the value is zero, the exponent is zero. A double, halfn, floatn or doublen argument representing an infinity or NaN is converted in the style of an f or F conversion specifier.

g,G

A double, halfn, floatn or doublen argument representing a floating-point number is converted in style f or e (or in style F or E in the case of a G conversion specifier), depending on the value converted and the precision. Let \(P\) equal the precision if nonzero, 6 if the precision is omitted, or 1 if the precision is zero. Then, if a conversion with style E would have an exponent of \(X\): — if \(P > X \geq -4\), the conversion is with style f (or F) and precision \(P - (X + 1)\). — otherwise, the conversion is with style e (or E) and precision \(P - 1\). Finally, unless the # flag is used, any trailing zeros are removed from the fractional portion of the result and the decimal-point character is removed if there is no fractional portion remaining. A double, halfn, floatn or doublen e argument representing an infinity or NaN is converted in the style of an f or F conversion specifier.

a,A

A double, halfn, floatn or doublen argument representing a floating-point number is converted in the style [-]0xh.hhhh p±d, where there is one hexadecimal digit (which is nonzero if the argument is a normalized floating-point number and is otherwise unspecified) before the decimal-point character\(^{55}\) and the number of hexadecimal digits after it is equal to the precision; if the precision is missing, then the precision is sufficient for an exact representation of the value; if the precision is zero and the # flag is not specified, no decimal point character appears. The letters abcdef are used for a conversion and the letters ABCDEF for A conversion. The A conversion specifier produces a number with X and P instead of x and p. The exponent always contains at least one digit, and only as many more digits as necessary to represent the decimal exponent of 2. If the value is zero, the exponent is zero. A double, halfn, floatn or doublen argument representing an infinity or NaN is converted in the style of an f or F conversion specifier.

\(^{54}\) When applied to infinite and NaN values, the -, +, and space flag characters have their usual meaning; the # and 0 flag characters have no effect.

\(^{55}\) Binary implementations can choose the hexadecimal digit to the left of the decimal-point character so that subsequent digits align to nibble (4-bit) boundaries.
NOTE: The conversion specifiers \texttt{e,E,g,G,a,A} convert a float or half argument that is a scalar type to a double only if the double data type is supported. If the double data type is not supported, the argument will be a \texttt{float} instead of a \texttt{double} and the half type will be converted to a float.

\textbf{c} The \texttt{int} argument is converted to an \texttt{unsigned char}, and the resulting character is written.

\textbf{s} The argument shall be a literal string.\footnote{No special provisions are made for multibyte characters. The behavior of \texttt{printf} with the \texttt{s} conversion specifier is undefined if the argument value is not a pointer to a literal string.} Characters from the literal string array are written up to (but not including) the terminating null character. If the precision is specified, no more than that many bytes are written. If the precision is not specified or is greater than the size of the array, the array shall contain a null character.

\textbf{p} The argument shall be a pointer to \texttt{void}. The pointer can refer to a memory region in the \texttt{global}, \texttt{constant}, \texttt{local}, \texttt{private} or \texttt{generic} address space. The value of the pointer is converted to a sequence of printing characters in an implementation-defined manner.

\textbf{\%} A \% character is written. No argument is converted. The complete conversion specification shall be \texttt{\%\%}.

If a conversion specification is invalid, the behavior is undefined. If any argument is not the correct type for the corresponding conversion specification, the behavior is undefined.

In no case does a nonexistent or small field width cause truncation of a field; if the result of a conversion is wider than the field width, the field is expanded to contain the conversion result.

For \texttt{a} and \texttt{A} conversions, the value is correctly rounded to a hexadecimal floating number with the given precision.

A few examples of \texttt{printf} are given below:

```c
float4 f = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
uchar4 uc = (uchar4)(0xFA, 0xFB, 0xFC, 0xFD);
printf("f4 = %2.2v4hlf\n", f);
printf("uc = %#v4hx\n", uc);
```

The above two \texttt{printf} calls print the following:

```
f4 = 1.00,2.00,3.00,4.00
uc = 0xfa,0xfb,0xfc,0xfd
```
A few examples of valid use cases of printf for the conversion specifier s are given below. The argument value must be a pointer to a literal string.

```
kernel void my_kernel( ... )
{
    printf("%s\n", "this is a test string\n");
}
```

A few examples of invalid use cases of printf for the conversion specifier s are given below:

```
kernel void my_kernel(global char *s, ... )
{
    printf("%s\n", s);

    constant char *p = "this is a test string\n";
    printf("%s\n", p);
    printf("%s\n", &p[3]);
}
```

A few examples of invalid use cases of printf where data types given by the vector specifier and length modifier do not match the argument type are given below:

```
kernelf void my_kernel(global char *s, ... )
{
    uint2 ui = (uint2)(0x12345678, 0x87654321);
    printf("unsigned short value = (%#v2hx)\n", ui)
    printf("unsigned char value = (%#v2hhx)\n", ui)
}
```

### 6.13.13.3 Differences between OpenCL C and C99 printf

- The l modifier followed by a c conversion specifier or s conversion specifier is not supported by OpenCL C.
- The ll, j, z, t, and L length modifiers are not supported by OpenCL C but are reserved.
- The n conversion specifier is not supported by OpenCL C but is reserved.
- OpenCL C adds the optional vn vector specifier to support printing of vector types.
- The conversion specifiers f, F, e, E, g, G, a, A convert a float argument to a double only if the double data type is supported. Refer to the description of `CL_DEVICE_DOUBLE_FP_CONFIG` in table 4.3. If the double data type is not supported, the argument will be a float instead of a double.
- For the embedded profile, the l length modifier is supported only if 64-bit integers are
In OpenCL C, `printf` returns 0 if it was executed successfully and -1 otherwise vs. C99 where `printf` returns the number of characters printed or a negative value if an output or encoding error occurred.

In OpenCL C, the conversion specifier `s` can only be used for arguments that are literal strings.
6.13.14 Image Read and Write Functions

The built-in functions defined in this section can only be used with image memory objects. An image memory object can be accessed by specific function calls that read from and/or write to specific locations in the image.

Image memory objects that are being read by a kernel should be declared with the `read_only` qualifier. `write_image` calls to image memory objects declared with the `read_only` qualifier will generate a compilation error. Image memory objects that are being written to by a kernel should be declared with the `write_only` qualifier. `read_image` calls to image memory objects declared with the `write_only` qualifier will generate a compilation error.

`read_image` and `write_image` calls to the same image memory object in a kernel are supported. Image memory objects that are being read and written by a kernel should be declared with the `read_write` qualifier.

The `read_image` calls returns a four component floating-point, integer or unsigned integer color value. The color values returned by `read_image` are identified as $x$, $y$, $z$, $w$ where $x$ refers to the red component, $y$ refers to the green component, $z$ refers to the blue component and $w$ refers to the alpha component.

6.13.14.1 Samplers

The image read functions take a sampler argument. The sampler can be passed as an argument to the kernel using `clSetKernelArg`, or can be declared in the outermost scope of `kernel` functions, or it can be a constant variable of type `sampler_t` declared in the program source.

Sampler variables in a program are declared to be of type `sampler_t`. A variable of `sampler_t` type declared in the program source must be initialized with a 32-bit unsigned integer constant, which is interpreted as a bit-field specifying the following properties:

- Addressing Mode
- Filter Mode
- Normalized Coordinates

These properties control how elements of an image object are read by `read_image{f|i|ui}`.

Samplers can also be declared as global constants in the program source using the following syntax.

```
const sampler_t <sampler name> = <value>
```

or

```
constant sampler_t <sampler name> = <value>
```

or
__constant sampler_t  <sampler_name> = <value>

Note that samplers declared using the **constant** qualifier are not counted towards the maximum number of arguments pointing to the constant address space or the maximum size of the constant address space allowed per device (i.e. **CL_DEVICE_MAX_CONSTANT_ARGS** and **CL_DEVICE_MAX_CONSTANT_BUFFER_SIZE** as described in table 4.3).

The sampler fields are described in table 6.22.

<table>
<thead>
<tr>
<th>Sampler State</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;normalized coords&gt;</td>
<td>Specifies whether the x, y and z coordinates are passed in as normalized or unnormalized values. This must be a literal value and can be one of the following predefined enums:</td>
</tr>
<tr>
<td></td>
<td>CLK_NORMALIZED_COORDS_TRUE or CLK_NORMALIZED_COORDS_FALSE.</td>
</tr>
<tr>
<td>The samplers used with an image in multiple calls to <strong>read_image</strong> declared in a kernel must use the same value for &lt;normalized coords&gt;.</td>
<td></td>
</tr>
<tr>
<td>&lt;addressing mode&gt;</td>
<td>Specifies the image addressing-mode i.e. how out-of-range image coordinates are handled. This must be a literal value and can be one of the following predefined enums:</td>
</tr>
<tr>
<td></td>
<td>CLK_ADDRESS_MIRRORED_REPEAT - Flip the image coordinate at every integer junction. This addressing mode can only be used with normalized coordinates. If normalized coordinates are not used, this addressing mode may generate image coordinates that are undefined.</td>
</tr>
<tr>
<td></td>
<td>CLK_ADDRESS_REPEAT – out-of-range image coordinates are wrapped to the valid range. This addressing mode can only be used with normalized coordinates. If normalized coordinates are not used, this addressing mode may generate image coordinates that are undefined.</td>
</tr>
<tr>
<td></td>
<td>CLK_ADDRESS_CLAMP_TO_EDGE – out-of-range image coordinates are clamped to the extent.</td>
</tr>
</tbody>
</table>
CLK_ADDRESS_CLAMP\textsuperscript{57} – out-of-range image coordinates will return a border color.

CLK_ADDRESS_NONE – for this addressing mode the programmer guarantees that the image coordinates used to sample elements of the image refer to a location inside the image; otherwise the results are undefined.

For 1D and 2D image arrays, the addressing mode applies only to the $x$ and $(x, y)$ coordinates. The addressing mode for the coordinate which specifies the array index is always CLK_ADDRESS_CLAMP_TO_EDGE.

<filter mode> Specifies the filter mode to use. This must be a literal value and can be one of the following predefined enums: CLK_FILTER_NEAREST or CLK_FILTER_LINEAR.

Refer to section 8.2 for a description of these filter modes.

<table>
<thead>
<tr>
<th>Table 6.22</th>
<th>Sampler Descriptor</th>
</tr>
</thead>
<tbody>
<tr>
<td>Examples:</td>
<td></td>
</tr>
</tbody>
</table>

```
const sampler_t samplerA = CLK_NORMALIZED_COORDS_TRUE |
CLK_ADDRESS_REPEAT                | 
CLK_FILTER_NEAREST;
```

samplerA specifies a sampler that uses normalized coordinates, the repeat addressing mode and a nearest filter.

The maximum number of samplers that can be declared in a kernel can be queried using the CL_DEVICE_MAX_SAMPLERS token in clGetDeviceInfo.

6.13.14.1.1 Determining the border color or value

If <addressing mode> in sampler is CLK_ADDRESS_CLAMP, then out-of-range image coordinates return the border color. The border color selected depends on the image channel order and can be one of the following values:

\textsuperscript{57} This is similar to the GL_ADDRESS_CLAMP_TO_BORDER addressing mode.
If the image channel order is CL_A, CL_INTENSITY, CL_Rx, CL_RA, CL_RGx, CL_RGBx, CL_sRGBx, CL_ARGB, CL_BGRA, CL_ABGR, CL_RGBA, CL_sRGBA or CL_sBGRA, the border color is (0.0f, 0.0f, 0.0f, 0.0f).

If the image channel order is CL_R, CL_RG, CL_RGB, or CL_LUMINANCE, the border color is (0.0f, 0.0f, 0.0f, 1.0f).

If the image channel order is CLDEPTH, the border value is 1.0f.

6.13.14.1.2 sRGB Images

The built-in image read functions will perform sRGB to linear RGB conversions if the image is an sRGB image. Writing to sRGB images from a kernel is an optional extension. The cl_khr_srgb_image_writes extension will be reported in the CL_DEVICE_EXTENSIONS string if a device supports writing to sRGB images using write_imagef. clGetSupportedImageFormats will return the supported sRGB images if CL_MEM_READ_WRITE or CL_MEM_WRITE_ONLY is specified in flags argument and the device supports writing to an sRGB image. If cl_khr_srgb_image_writes is supported, the built-in image write functions will perform the linear to sRGB conversion.

Only the R, G and B components are converted from linear to sRGB and vice-versa. The alpha component is returned as is.
### 6.13.14.2 Built-in Image Read Functions

The following built-in function calls to read images with a sampler\(^{58}\) are supported.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>float4 read_imagef (image2d_t image, sampler_t sampler, int2 coord)</code></td>
<td>Use the coordinate <code>(coord.x, coord.y)</code> to do an element lookup in the 2D image object specified by <code>image</code>. <code>read_imagef</code> returns floating-point values in the range ([0.0 \ldots 1.0]) for image objects created with <code>image_channel_data_type</code> set to one of the pre-defined packed formats or <code>CL_UNORM_INT8</code>, or <code>CL_UNORM_INT16</code>. <code>read_imagef</code> returns floating-point values in the range ([-1.0 \ldots 1.0]) for image objects created with <code>image_channel_data_type</code> set to <code>CL_SNORM_INT8</code>, or <code>CL_SNORM_INT16</code>. <code>read_imagef</code> returns floating-point values for image objects created with <code>image_channel_data_type</code> set to <code>CL_HALF_FLOAT</code> or <code>CL_FLOAT</code>. The <code>read_imagef</code> calls that take integer coordinates must use a sampler with filter mode set to <code>CLK_FILTER_NEAREST</code>, normalized coordinates set to <code>CLK_NORMALIZED_COORDS_FALSE</code> and addressing mode set to <code>CLK_ADDRESS_CLAMP_TO_EDGE</code>, <code>CLK_ADDRESS_CLAMP</code> or <code>CLK_ADDRESS_NONE</code>; otherwise the values returned are undefined. Values returned by <code>read_imagef</code> for image objects with <code>image_channel_data_type</code> values not specified in the description above are undefined.</td>
</tr>
<tr>
<td><code>int4 read_imagei (image2d_t image, sampler_t sampler, int2 coord)</code></td>
<td>Use the coordinate <code>(coord.x, coord.y)</code> to do an element lookup in the 2D image object specified by <code>image</code>. <code>read_imagei</code> and <code>read_imageui</code> return</td>
</tr>
</tbody>
</table>

\(^{58}\) The built-in function calls to read images with a sampler are not supported for `image1d_buffer_t` image types.
sampler_t sampler,
float2 coord)

uint4 **read_imageui** (  
*image2d_t image,
sampler_t sampler,
int2 coord)

**read_imageui** can only be used with image objects created with **image_channel_data_type** set to one of the following values:  
CL_SIGNED_INT8,  
CL_SIGNED_INT16 and  
CL_SIGNED_INT32.  
If the **image_channel_data_type** is not one of the above values, the values returned by **read_imageui** are undefined.

**read_imageui** can only be used with image objects created with **image_channel_data_type** set to one of the following values:  
CL_UNSIGNED_INT8,  
CL_UNSIGNED_INT16 and  
CL_UNSIGNED_INT32.  
If the **image_channel_data_type** is not one of the above values, the values returned by **read_imageui** are undefined.

The **read_image{i|ui}** calls support a nearest filter only. The **filter_mode** specified in **sampler** must be set to CLK_FILTER_NEAREST; otherwise the values returned are undefined.

Furthermore, the **read_image{i|ui}** calls that take integer coordinates must use a sampler with normalized coordinates set to  
CLK_NORMALIZED_COORDS_FALSE and addressing mode set to  
CLK_ADDRESS_CLAMP_TO_EDGE,  
CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined.

| float4 **read_imagef** (image3d_t image, sampler_t sampler, int4 coord )  | Use the coordinate (coord.x, coord.y, coord.z) to do an element lookup in the 3D image object specified by image. coord.w is ignored.  |
| float4 **read_imagef** (image3d_t image, sampler_t sampler, float4 coord)  | **read_imagef** returns floating-point values in the range [0.0 … 1.0] for image objects created with **image_channel_data_type** set to one of the pre-defined packed formats or CL_UNORM_INT8, or |
**CL_UNORM_INT16.**

read_imagef returns floating-point values in the range \([-1.0 \ldots 1.0]\) for image objects created with `image_channel_data_type` set to CL_SNORM_INT8, or CL_SNORM_INT16.

read_imagef returns floating-point values for image objects created with `image_channel_data_type` set to CL_HALF_FLOAT or CL_FLOAT.

The read_imagef calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined.

Values returned by read_imagef for image objects with `image_channel_data_type` values not specified in the description are undefined.

<table>
<thead>
<tr>
<th>Function</th>
<th>Parameters</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>int4 read_imagei</td>
<td>(image3d_t image, sampler_t sampler, int4 coord)</td>
<td>Use the coordinate ((coord.x, coord.y, coord.z)) to do an element lookup in the 3D image object specified by <code>image. coord.w</code> is ignored.</td>
</tr>
<tr>
<td>int4 read_imagei</td>
<td>(image3d_t image, sampler_t sampler, float4 coord)</td>
<td>read_imagei and read_imagei return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer.</td>
</tr>
<tr>
<td>uint4 read_imageui</td>
<td>(image3d_t image, sampler_t sampler, int4 coord)</td>
<td><strong>read_imagei</strong> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: CL_SIGNED_INT8, CL_SIGNED_INT16 and CL_SIGNED_INT32. If the <code>image_channel_data_type</code> is not one of the above values, the values returned by <strong>read_imagei</strong> are undefined.</td>
</tr>
<tr>
<td>uint4 read_imageui</td>
<td>(image3d_t image, sampler_t sampler, float4 coord)</td>
<td><strong>read_imagei</strong> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values:</td>
</tr>
</tbody>
</table>
CL_UNSIGNED_INT8, CL_UNSIGNED_INT16 and CL_UNSIGNED_INT32.

If the image_channel_data_type is not one of the above values, the values returned by read_imageui are undefined.

The read_image{i|ui} calls support a nearest filter only. The filter_mode specified in sampler must be set to CLK_FILTER_NEAREST; otherwise the values returned are undefined.

Furthermore, the read_image{i|ui} calls that take integer coordinates must use a sampler with normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined.

float4 read_imagef (image2d_array_t image, sampler_t sampler, int4 coord)

Use coord.xy to do an element lookup in the 2D image identified by coord.z in the 2D image array specified by image.

read_imagef returns floating-point values in the range [0.0 … 1.0] for image objects created with image_channel_data_type set to one of the pre-defined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16.

read_imagef returns floating-point values in the range [-1.0 … 1.0] for image objects created with image_channel_data_type set to CL_SNORM_INT8, or CL_SNORM_INT16.

read_imagef returns floating-point values for image objects created with image_channel_data_type set to CL_HALF_FLOAT or CL_FLOAT.

The read_imagef calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE,
<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
</table>
| `int4 read_imagei` | Use `coord.xy` to do an element lookup in the 2D image identified by `coord.z` in the 2D image array specified by `image`.
| `uint4 read_imageui` | The `read_imagei` and `read_imageui` return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer.

**read_imagei** can only be used with image objects created with `image_channel_data_type` set to one of the following values:
- CL_SIGNED_INT8
- CL_SIGNED_INT16
- CL_SIGNED_INT32

If the `image_channel_data_type` is not one of the above values, the values returned by `read_imagei` are undefined.

**read_imageui** can only be used with image objects created with `image_channel_data_type` set to one of the following values:
- CL_UNSIGNED_INT8
- CL_UNSIGNED_INT16
- CL_UNSIGNED_INT32

If the `image_channel_data_type` is not one of the above values, the values returned by `read_imageui` are undefined.

The `read_imagei{ui}` calls support a nearest filter only. The filter mode specified in `sampler` must be set to CLK_FILTER_NEAREST; otherwise the values returned are undefined.

Furthermore, the `read_imagei{ui}` calls that take integer coordinates must use a sampler with normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE.
<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>float4 read_imagef(image1d_t image, sampler_t sampler, int coord)</code></td>
<td>Use <code>coord</code> to do an element lookup in the 1D image object specified by <code>image</code>. <code>read_imagef</code> returns floating-point values in the range ([0.0 \ldots 1.0]) for image objects created with <code>image_channel_data_type</code> set to one of the predefined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16.</td>
</tr>
<tr>
<td><code>float4 read_imagef(image1d_t image, sampler_t sampler, float coord)</code></td>
<td><code>read_imagef</code> returns floating-point values in the range ([-1.0 \ldots 1.0]) for image objects created with <code>image_channel_data_type</code> set to CL_SNORM_INT8, or CL_SNORM_INT16.</td>
</tr>
<tr>
<td></td>
<td><code>read_imagef</code> returns floating-point values for image objects created with <code>image_channel_data_type</code> set to CL_HALF_FLOAT or CL_FLOAT.</td>
</tr>
<tr>
<td></td>
<td>The <code>read_imagef</code> calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined.</td>
</tr>
<tr>
<td></td>
<td>Values returned by <code>read_imagef</code> for image objects with <code>image_channel_data_type</code> values not specified in the description above are undefined.</td>
</tr>
<tr>
<td><code>int4 read_imagei(image1d_t image, sampler_t sampler, int coord)</code></td>
<td>Use <code>coord</code> to do an element lookup in the 1D image object specified by <code>image</code>. <code>read_imagei</code> and <code>read_imageui</code> return unnormalized signed integer and unnormalized signed integer values respectively. Each channel will be stored in a 32-bit integer.</td>
</tr>
<tr>
<td><code>int4 read_imagei(image1d_t image, sampler_t sampler, float coord)</code></td>
<td><code>read_imagei</code> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: CL_SIGNED_INT8, CL_UNSIGNED_INT8, CL_HALF_FLOAT, CL_FLOAT.</td>
</tr>
<tr>
<td><code>uint4 read_imageui(image1d_t image, sampler_t sampler, float coord)</code></td>
<td><code>read_imageui</code> only return values of CL_HALF_FLOAT, CL_FLOAT, CL_DOUBLE_FLOAT, CL_HALF_FLOAT, CL_DOUBLE_FLOAT. Each channel will be stored in a 32-bit integer.</td>
</tr>
<tr>
<td>Function</td>
<td>Description</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
</tr>
<tr>
<td><code>uint4 read_imageui (image1d_t image, sampler_t sampler, float coord)</code></td>
<td>Returns integers in the range -32768 to 32767 if <code>image_channel_data_type</code> is <code>CL_SIGNED_INT16</code> or <code>CL_SIGNED_INT32</code>. Returns an undefined value if the <code>image_channel_data_type</code> is not one of these values.</td>
</tr>
<tr>
<td><code>float4 read_imagef (image1d_array_t image, sampler_t sampler, int2 coord)</code></td>
<td>Use <code>coord.x</code> to do an element lookup in the 1D image identified by <code>coord.y</code> in the 1D image array specified by <code>image</code>. Returns floating-point values in the range [0.0 .. 1.0] for image objects created with <code>image_channel_data_type</code> set to one of the predefined packed formats or <code>CL_UNORM_INT8</code>, or <code>CL_UNORM_INT16</code>.</td>
</tr>
<tr>
<td><code>float4 read_imagef (image1d_array_t image, sampler_t sampler, float2 coord)</code></td>
<td>Use <code>coord.x</code> to do an element lookup in the 1D image identified by <code>coord.y</code> in the 1D image array specified by <code>image</code>. Returns floating-point values in the range [-1.0 .. 1.0] for image objects created with <code>image_channel_data_type</code> set to <code>CL_SNORM_INT8</code>, or <code>CL_SNORM_INT16</code>.</td>
</tr>
</tbody>
</table>

*If the `image_channel_data_type` is not one of the above values, the values returned by `read_imagei` are undefined.*

*`read_imageui` can only be used with image objects created with `image_channel_data_type` set to one of the following values: `CL_UNSIGNED_INT8`, `CL_UNSIGNED_INT16` and `CL_UNSIGNED_INT32`. If the `image_channel_data_type` is not one of the above values, the values returned by `read_imageui` are undefined.*

*The `read_image{i|ui}` calls support a nearest filter only. The filter_mode specified in `sampler` must be set to CLK_FILTER_NEAREST; otherwise the values returned are undefined.*

*Furthermore, the `read_image{i|ui}` calls that take integer coordinates must use a sampler with normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NO; otherwise the values returned are undefined.*
objects created with `image_channel_data_type` set to `CL_HALF_FLOAT` or `CL_FLOAT`.

The **read_imagef** calls that take integer coordinates must use a sampler with filter mode set to `CLK_FILTER_NEAREST`, normalized coordinates set to `CLK_NORMALIZED_COORDS_FALSE` and addressing mode set to `CLK_ADDRESS_CLAMP_TO_EDGE`, `CLK_ADDRESS_CLAMP` or `CLK_ADDRESS_NONE`; otherwise the values returned are undefined.

Values returned by **read_imagef** for image objects with `image_channel_data_type` values not specified in the description above are undefined.

```c
int4 read_imagei (image1d_array_t image, sampler_t sampler, int2 coord)
```

Use `coord.x` to do an element lookup in the 1D image identified by `coord.y` in the 1D image array specified by `image`.

**read_imagei** and **read_imageui** return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer.

**read_imagei** can only be used with image objects created with `image_channel_data_type` set to one of the following values:
- `CL_SIGNED_INT8`
- `CL_SIGNED_INT16`
- `CL_SIGNED_INT32`

If the `image_channel_data_type` is not one of the above values, the values returned by **read_imagei** are undefined.

```c
uint4 read_imageui (image1d_array_t image, sampler_t sampler, int2 coord)
```

**read_imageui** can only be used with image objects created with `image_channel_data_type` set to one of the following values:
- `CL_UNSIGNED_INT8`
- `CL_UNSIGNED_INT16`
- `CL_UNSIGNED_INT32`

If the `image_channel_data_type` is not one of the above values, the values returned by **read_imageui** are undefined.

The **read_image{i|ui}** calls support a nearest filter only. The filter_mode specified in `sampler`
must be set to \texttt{CLK\_FILTER\_NEAREST}; otherwise the values returned are undefined.

Furthermore, the \texttt{read\_image[i|ui]} calls that take integer coordinates must use a sampler with normalized coordinates set to \texttt{CLK\_NORMALIZED\_COORDS\_FALSE} and addressing mode set to \texttt{CLK\_ADDRESS\_CLAMP\_TO\_EDGE}, \texttt{CLK\_ADDRESS\_CLAMP} or \texttt{CLK\_ADDRESS\_NONE}; otherwise the values returned are undefined.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{float read_imagef (int2 coord)}</td>
<td>Use the coordinate ((\text{coord.x, coord.y})) to do an element lookup in the 2D depth image object specified by \texttt{image}. \texttt{read_imagef} returns a floating-point value in the range ([0.0 \ldots 1.0]) for depth image objects created with \texttt{image_channel_data_type} set to \texttt{CL_UNORM_INT16} or \texttt{CL_UNORM_INT24}. \texttt{read_imagef} returns a floating-point value for depth image objects created with \texttt{image_channel_data_type} set to \texttt{CL_FLOAT}. The \texttt{read_imagef} calls that take integer coordinates must use a sampler with filter mode set to \texttt{CLK_FILTER_NEAREST}, normalized coordinates set to \texttt{CLK_NORMALIZED_COORDS_FALSE} and addressing mode set to \texttt{CLK_ADDRESS_CLAMP_TO_EDGE}, \texttt{CLK_ADDRESS_CLAMP} or \texttt{CLK_ADDRESS_NONE}; otherwise the values returned are undefined. Values returned by \texttt{read_imagef} for depth image objects with \texttt{image_channel_data_type} values not specified in the description above are undefined.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{float read_imagef (float2 coord)}</td>
<td>Use \texttt{coord.xy} to do an element lookup in the 2D image identified by \texttt{coord.z} in the 2D depth image array specified by \texttt{image}. \texttt{read_imagef} returns a floating-point value in the range ([0.0 \ldots 1.0]) for depth image objects created with \texttt{image_channel_data_type} set to \texttt{CL_UNORM_INT16} or \texttt{CL_UNORM_INT24}.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{float read_imagef (int4 coord)}</td>
<td>Use \texttt{coord.xy} to do an element lookup in the 2D image identified by \texttt{coord.z} in the 2D depth image array specified by \texttt{image}. \texttt{read_imagef} returns a floating-point value in the range ([0.0 \ldots 1.0]) for depth image objects created with \texttt{image_channel_data_type} set to \texttt{CL_UNORM_INT16} or \texttt{CL_UNORM_INT24}.</td>
</tr>
</tbody>
</table>
**Table 6.23  Built-in Image Read Functions**

### 6.13.14.3  Built-in Image Sampler-less Read Functions

The sampler-less read image functions behave exactly as the corresponding read image functions described in section 6.13.14.2 that take integer coordinates and a sampler with filter mode set to CL_FILTER_NEAREST, normalized coordinates set to CL_NORMALIZED_COORDS_FALSE and addressing mode to CL_ADDRESS_NONE.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>float4 <strong>read_imagef</strong> (image2d_t image, int2 coord)</td>
<td>Use the coordinate (coord.x, coord.y) to do an element lookup in the 2D image object specified by image.</td>
</tr>
</tbody>
</table>

**read_imagef** returns floating-point values in the range [0.0 … 1.0] for image objects created with `image_channel_data_type` set to one of the pre-defined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16.

**read_imagef** returns floating-point values in the range [-1.0 … 1.0] for image objects created with `image_channel_data_type` set to CL_SNORM_INT8, or CL_SNORM_INT16.
**read_imagef** returns floating-point values for image objects created with `image_channel_data_type` set to `CL_HALF_FLOAT` or `CL_FLOAT`.

Values returned by **read_imagef** for image objects with `image_channel_data_type` values not specified in the description above are undefined.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>int4 read_imagei</strong> (image2d_t image, int2 coord)</td>
<td>Use the coordinate <code>(coord.x, coord.y)</code> to do an element lookup in the 2D image object specified by <code>image</code>. <strong>read_imagei</strong> and <strong>read_imageui</strong> return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer. <strong>read_imagei</strong> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: <code>CL_SIGNED_INT8</code>, <code>CL_SIGNED_INT16</code> and <code>CL_SIGNED_INT32</code>. If the <code>image_channel_data_type</code> is not one of the above values, the values returned by <strong>read_imagei</strong> are undefined. <strong>read_imageui</strong> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: <code>CL_UNSIGNED_INT8</code>, <code>CL_UNSIGNED_INT16</code> and <code>CL_UNSIGNED_INT32</code>. If the <code>image_channel_data_type</code> is not one of the above values, the values returned by <strong>read_imageui</strong> are undefined.</td>
</tr>
<tr>
<td><strong>uint4 read_imageui</strong> ( image2d_t image, int2 coord)</td>
<td></td>
</tr>
<tr>
<td><strong>float4 read_imagef</strong> (image3d_t image, int4 coord)</td>
<td>Use the coordinate <code>(coord.x, coord.y, coord.z)</code> to do an element lookup in the 3D image object specified by <code>image</code>. <code>coord.w</code> is ignored. <strong>read_imagef</strong> returns floating-point values in the range <code>[0.0 … 1.0]</code> for image objects created with <code>image_channel_data_type</code> set to one of the pre-defined packed formats or <code>CL_UNORM_INT8</code>, or <code>CL_UNORM_INT16</code>.</td>
</tr>
</tbody>
</table>
**read_image** returns floating-point values in the range \([-1.0 \ldots 1.0]\) for image objects created with `image_channel_data_type` set to `CL_SNORM_INT8`, or `CL_SNORM_INT16`.

**read_image** returns floating-point values for image objects created with `image_channel_data_type` set to `CL_HALF_FLOAT` or `CL_FLOAT`.

Values returned by **read_image** for image objects with `image_channel_data_type` values not specified in the description are undefined.

<table>
<thead>
<tr>
<th>int4 <code>read_imagei</code> (image3d_t <code>image</code>, int4 <code>coord</code>)</th>
<th>Use the coordinate ((coord.x, coord.y, coord.z)) to do an element lookup in the 3D image object specified by <code>image</code>. <code>coord.w</code> is ignored.</th>
</tr>
</thead>
<tbody>
<tr>
<td>uint4 <code>read_imageui</code> (image3d_t <code>image</code>, int4 <code>coord</code>)</td>
<td><strong>read_imagei</strong> and <strong>read_imageui</strong> return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer. <strong>read_imagei</strong> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: <code>CL_SIGNED_INT8</code>, <code>CL_SIGNED_INT16</code> and <code>CL_SIGNED_INT32</code>. If the <code>image_channel_data_type</code> is not one of the above values, the values returned by <strong>read_imagei</strong> are undefined. <strong>read_imageui</strong> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: <code>CL_UNSIGNED_INT8</code>, <code>CL_UNSIGNED_INT16</code> and <code>CL_UNSIGNED_INT32</code>. If the <code>image_channel_data_type</code> is not one of the above values, the values returned by <strong>read_imageui</strong> are undefined.</td>
</tr>
<tr>
<td>float4 <code>read_imagef</code> (image2d_array_t <code>image</code>, int4 <code>coord</code>)</td>
<td>Use <code>coord.xy</code> to do an element lookup in the 2D image identified by <code>coord.z</code> in the 2D image array specified by <code>image</code>.</td>
</tr>
</tbody>
</table>
read_imagef returns floating-point values in the range [0.0 ... 1.0] for image objects created with `image_channel_data_type` set to one of the pre-defined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16.

read_imagef returns floating-point values in the range [-1.0 ... 1.0] for image objects created with `image_channel_data_type` set to CL_SNORM_INT8, or CL_SNORM_INT16.

read_imagef returns floating-point values for image objects created with `image_channel_data_type` set to CL_HALF_FLOAT or CL_FLOAT.

Values returned by `read_imagef` for image objects with `image_channel_data_type` values not specified in the description above are undefined.

<table>
<thead>
<tr>
<th>int4 read_imagei (image2d_array_t image, int4 coord)</th>
<th>Use <code>coord.xy</code> to do an element lookup in the 2D image identified by <code>coord.z</code> in the 2D image array specified by <code>image</code>.</th>
</tr>
</thead>
<tbody>
<tr>
<td>uint4 read_imageui (image2d_array_t image, int4 coord)</td>
<td><code>read_imagei</code> and <code>read_imageui</code> return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer.</td>
</tr>
</tbody>
</table>

`read_imagei` can only be used with image objects created with `image_channel_data_type` set to one of the following values:
- CL_SIGNED_INT8,
- CL_SIGNED_INT16 and
- CL_SIGNED_INT32.

If the `image_channel_data_type` is not one of the above values, the values returned by `read_imagei` are undefined.

`read_imageui` can only be used with image objects created with `image_channel_data_type` set to one of the following values:
- CL UNSIGNED_INT8,
- CL UNSIGNED_INT16 and
- CL UNSIGNED_INT32.

If the `image_channel_data_type` is not one of the
Above values, the values returned by `read_imageui` are undefined.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>float4 read_imagef(image1d_t image, int coord)</code></td>
<td>Use <code>coord</code> to do an element lookup in the 1D image or 1D image buffer object specified by <code>image</code>. <code>read_imagef</code> returns floating-point values in the range $[0.0 \ldots 1.0]$ for image objects created with <code>image_channel_data_type</code> set to one of the pre-defined packed formats or <code>CL_UNORM_INT8</code>, or <code>CL_UNORM_INT16</code>. <code>read_imagef</code> returns floating-point values in the range $[-1.0 \ldots 1.0]$ for image objects created with <code>image_channel_data_type</code> set to <code>CL_SNORM_INT8</code>, or <code>CL_SNORM_INT16</code>. <code>read_imagef</code> returns floating-point values for image objects created with <code>image_channel_data_type</code> set to <code>CL_HALF_FLOAT</code> or <code>CL_FLOAT</code>. Values returned by <code>read_imagef</code> for image objects with <code>image_channel_data_type</code> values not specified in the description above are undefined.</td>
</tr>
<tr>
<td><code>int4 read_imagei(image1d_t image, int coord)</code></td>
<td>Use <code>coord</code> to do an element lookup in the 1D image or 1D image buffer object specified by <code>image</code>. <code>read_imagei</code> and <code>read_imageui</code> return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer. <code>read_imagei</code> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: <code>CL_SIGNED_INT8</code>, <code>CL_SIGNED_INT16</code> and <code>CL_SIGNED_INT32</code>. If the <code>image_channel_data_type</code> is not one of the above values, the values returned by <code>read_imagei</code> are undefined. <code>read_imageui</code> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values:</td>
</tr>
<tr>
<td><code>uint4 read_imageui(image1d_t image, int coord)</code></td>
<td></td>
</tr>
<tr>
<td><code>int4 read_imagei(image1d_buffer_t image, int coord)</code></td>
<td></td>
</tr>
<tr>
<td><code>uint4 read_imageui(image1d_buffer_t image, int coord)</code></td>
<td></td>
</tr>
</tbody>
</table>
CL_UNSIGNED_INT8, CL_UNSIGNED_INT16 and CL_UNSIGNED_INT32.
If the `image_channel_data_type` is not one of the above values, the values returned by `read_imageui` are undefined.

| float4 `read_imagef` (  
| image1d_array_t image,  
| int2 coord) | Use `coord.x` to do an element lookup in the 1D image identified by `coord.y` in the 1D image array specified by `image`.

`read_imagef` returns floating-point values in the range [0.0 … 1.0] for image objects created with `image_channel_data_type` set to one of the pre-defined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16.

`read_imagef` returns floating-point values in the range [-1.0 … 1.0] for image objects created with `image_channel_data_type` set to CL_SNORM_INT8, or CL_SNORM_INT16.

`read_imagef` returns floating-point values for image objects created with `image_channel_data_type` set to CL_HALF_FLOAT or CL_FLOAT.

Values returned by `read_imagef` for image objects with `image_channel_data_type` values not specified in the description above are undefined.

| int4 `read_imagei` (image1d_array_t image,  
| int2 coord) | Use `coord.x` to do an element lookup in the 1D image identified by `coord.y` in the 1D image array specified by `image`.

`read_imagei` and `read_imageui` return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer.

`read_imagei` can only be used with image objects created with `image_channel_data_type` set to one of the following values:
CL_SIGNED_INT8, CL_SIGNED_INT16 and CL_SIGNED_INT32.
If the `image_channel_data_type` is not one of the
above values, the values returned by `read_imageu` are undefined.

`read_imageui` can only be used with image objects created with `image_channel_data_type` set to one of the following values:
CL_UNSIGNED_INT8, CL_UNSIGNED_INT16 and CL_UNSIGNED_INT32.
If the `image_channel_data_type` is not one of the above values, the values returned by `read_imageui` are undefined.

<table>
<thead>
<tr>
<th>float <code>read_imagef</code> (image2d_depth_t <code>image</code>, int2 <code>coord</code>)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Use the coordinate <code>(coord.x, coord.y)</code> to do an element lookup in the 2D depth image object specified by <code>image</code>.</td>
</tr>
<tr>
<td><code>read_imagef</code> returns a floating-point value in the range [0.0 … 1.0] for depth image objects created with <code>image_channel_data_type</code> set to CL_UNORM_INT16 or CL_UNORM_INT24.</td>
</tr>
<tr>
<td><code>read_imagef</code> returns a floating-point value for depth image objects created with <code>image_channel_data_type</code> set to CL_FLOAT.</td>
</tr>
<tr>
<td>Values returned by <code>read_imagef</code> for image objects with <code>image_channel_data_type</code> values not specified in the description above are undefined.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>float <code>read_imagef</code> (image2d_array_depth_t <code>image</code>, int4 <code>coord</code>)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Use <code>coord.xy</code> to do an element lookup in the 2D image identified by <code>coord.z</code> in the 2D depth image array specified by <code>image</code>.</td>
</tr>
<tr>
<td><code>read_imagef</code> returns a floating-point value in the range [0.0 … 1.0] for depth image objects created with <code>image_channel_data_type</code> set to CL_UNORM_INT16 or CL_UNORM_INT24.</td>
</tr>
<tr>
<td><code>read_imagef</code> returns a floating-point value for depth image objects created with <code>image_channel_data_type</code> set to CL_FLOAT.</td>
</tr>
<tr>
<td>Values returned by <code>read_imagef</code> for image objects with <code>image_channel_data_type</code> values not specified in the description above are undefined.</td>
</tr>
</tbody>
</table>
Table 6.24  Built-in Image Sampler-less Read Functions

### 6.13.14.4 Built-in Image Write Functions

The following built-in function calls to write images are supported. Note that image writes to sRGB images are only supported if the cl_khr_srgb_image_writes extension is supported; otherwise the behavior of writing to a sRGB image is undefined.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>void write_imagef</code></td>
<td>Write color value to location specified by <code>coord.xy</code> in the 2D image object specified by <code>image</code>. Appropriate data format conversion to the specified image format is done before writing the color value. <code>coord.x</code> and <code>coord.y</code> are considered to be unnormalized coordinates and must be in the range 0 ... image width – 1, and 0 ... image height – 1. <code>write_imagef</code> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the pre-defined packed formats or set to <code>CL_SNORM_INT8</code>, <code>CL_UNORM_INT8</code>, <code>CL_SNORM_INT16</code>, <code>CL_UNORM_INT16</code>, <code>CL_HALF_FLOAT</code> or <code>CL_FLOAT</code>. Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored.</td>
</tr>
<tr>
<td><code>void write_imagei</code></td>
<td><code>write_imagei</code> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: <code>CL_SIGNED_INT8</code>, <code>CL_SIGNED_INT16</code> and <code>CL_SIGNED_INT32</code>.</td>
</tr>
<tr>
<td><code>void write_imageui</code></td>
<td><code>write_imageui</code> can only be used with image objects created with <code>image_channel_data_type</code> set to one of the following values: <code>CL_UNSIGNED_INT8</code>, <code>CL_UNSIGNED_INT16</code> and <code>CL_UNSIGNED_INT32</code>.</td>
</tr>
</tbody>
</table>

The behavior of `write_imagef`, `write_imagei` and `write_imageui` for image objects created with
image_channel_data_type values not specified in the description above or with \((x,y)\) coordinate values that are not in the range \((0 \ldots \text{image width} – 1, 0 \ldots \text{image height} – 1)\), respectively, is undefined.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>void <strong>write_imagef</strong> (</td>
<td>Write color value to location specified by (\text{coord}.xy) in the 2D image identified by (\text{coord}.z) in the 2D image array specified by (\text{image}). Appropriate data format conversion to the specified image format is done before writing the color value. (\text{coord}.x), (\text{coord}.y) and (\text{coord}.z) are considered to be unnormalized coordinates and must be in the range (0 \ldots \text{image width} – 1, 0 \ldots \text{image height} – 1) and (0 \ldots \text{image number of layers} – 1).</td>
</tr>
<tr>
<td>image2d_array_t (\text{image}), int4 (\text{coord}), float4 (\text{color})</td>
<td></td>
</tr>
<tr>
<td>void <strong>write_imagei</strong> (</td>
<td><strong>write_imagef</strong> can only be used with image objects created with (\text{image}<em>\text{channel}</em>\text{data}<em>\text{type}) set to one of the pre-defined packed formats or set to (\text{CL}</em>\text{SNORM}<em>\text{INT8}, \text{CL}</em>\text{UNORM}<em>\text{INT8}, \text{CL}</em>\text{SNORM}<em>\text{INT16}, \text{CL}</em>\text{UNORM}<em>\text{INT16}, \text{CL}</em>\text{HALF}<em>\text{FLOAT}) or (\text{CL}</em>\text{FLOAT}). Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored.</td>
</tr>
<tr>
<td>image2d_array_t (\text{image}), int4 (\text{coord}), int4 (\text{color})</td>
<td><strong>write_imagei</strong> can only be used with image objects created with (\text{image}<em>\text{channel}</em>\text{data}<em>\text{type}) set to one of the following values: (\text{CL}</em>\text{SIGNED}<em>\text{INT8}, \text{CL}</em>\text{SIGNED}<em>\text{INT16} ) and (\text{CL}</em>\text{SIGNED}_\text{INT32}).</td>
</tr>
<tr>
<td>void <strong>write_imageui</strong> (</td>
<td><strong>write_imagei</strong> can only be used with image objects created with (\text{image}<em>\text{channel}</em>\text{data}<em>\text{type}) set to one of the following values: (\text{CL}</em>\text{UNSIGN}<em>\text{INT8}, \text{CL}</em>\text{UNSIGN}<em>\text{INT16} ) and (\text{CL}</em>\text{UNSIGN}_\text{INT32}).</td>
</tr>
<tr>
<td>image2d_array_t (\text{image}), int4 (\text{coord}), uint4 (\text{color})</td>
<td>The behavior of <strong>write_imagef</strong>, <strong>write_imagei</strong> and <strong>write_imageui</strong> for image objects created with (\text{image}<em>\text{channel}</em>\text{data}_\text{type}) values not specified in the description above or with ((x,y,z)) coordinate values that are not in the range ((0 \ldots \text{image width} – 1, 0 \ldots \text{image height} – 1, 0 \ldots \text{image number of layers} – 1)), respectively, is undefined.</td>
</tr>
</tbody>
</table>
### void write_imagef

```c
void write_imagef (image1d_t image, int coord, float4 color)
```

Write color value to location specified by `coord` in the 1D image or 1D image buffer object specified by `image`. Appropriate data format conversion to the specified image format is done before writing the color value. `coord` is considered to be unnormalized coordinates and must be in the range 0 ... image width – 1.

### void write_imagei

```c
void write_imagei (image1d_t image, int coord, int4 color)
```

### void write_imageui

```c
void write_imageui (image1d_t image, int coord, uint4 color)
```

### void write_imagef

```c
void write_imagef (image1d_buffer_t image, int coord, float4 color)
```

### void write_imagei

```c
void write_imagei (image1d_buffer_t image, int coord, int4 color)
```

### void write_imageui

```c
void write_imageui (image1d_buffer_t image, int coord, uint4 color)
```

write_imagef can only be used with image objects created with `image_channel_data_type` set to one of the pre-defined packed formats or set to `CL_SNORM_INT8, CL_UNORM_INT8, CL_SNORM_INT16, CL_UNORM_INT16, CL_HALF_FLOAT or CL_FLOAT`. Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored.

write_imagei can only be used with image objects created with `image_channel_data_type` set to one of the following values:
- `CL_SIGNED_INT8`
- `CL_SIGNED_INT16`
- `CL_SIGNED_INT32`

write_imageui can only be used with image objects created with `image_channel_data_type` set to one of the following values:
- `CL_UNSIGNED_INT8`
- `CL_UNSIGNED_INT16`
- `CL_UNSIGNED_INT32`

The behavior of `write_imagef`, `write_imagei` and `write_imageui` for image objects created with `image_channel_data_type` values not specified in the description above or with coordinate values that is not in the range (0 … image width – 1), is undefined.

### void write_imagef

```c
void write_imagef (image1d_array_t image, int2 coord, float4 color)
```

Write color value to location specified by `coord.x` in the 1D image identified by `coord.y` in the 1D image array specified by `image`. Appropriate data format conversion to the specified image format is done before writing the color value. `coord.x` and `coord.y` are considered to be unnormalized coordinates and
image1d_array_t image, int2 coord, int4 color)

void write_imageu ( 
    image1d_array_t image, int2 coord, uint4 color)

must be in the range 0 ... image width – 1 and 0 ... image number of layers – 1.

write_imagef can only be used with image objects created with image_channel_data_type set to one of the pre-defined packed formats or set to CL_SNORM_INT8, CL_UNORM_INT8, CL_SNORM_INT16, CL_UNORM_INT16, CL_HALF_FLOAT or CL_FLOAT. Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored.

write_imagei can only be used with image objects created with image_channel_data_type set to one of the following values:
CL_SIGNED_INT8,
CL_SIGNED_INT16 and
CL_SIGNED_INT32.

write_imageui can only be used with image objects created with image_channel_data_type set to one of the following values:
CL_UNSIGNED_INT8,
CL_UNSIGNED_INT16 and
CL_UNSIGNED_INT32.

The behavior of write_imagef, write_imagei and write_imageui for image objects created with image_channel_data_type values not specified in the description above or with (x, y) coordinate values that are not in the range (0 ... image width – 1, 0 ... image number of layers – 1), respectively, is undefined.

void write_imagef ( 
    image2d_depth_t image, int2 coord, float depth)

Write depth value to location specified by coord.xy in the 2D depth image object specified by image. Appropriate data format conversion to the specified image format is done before writing the depth value. coord.x and coord.y are considered to be unnormalized coordinates and must be in the range 0 ... image width – 1, and 0 ... image height – 1.

write_imagef can only be used with image objects created with image_channel_data_type set to CL_UNORM_INT16, CL_UNORM_INT24 or
Appropriate data format conversion will be done to convert depth value from a floating-point value to actual data format associated with the image.

The behavior of **write_imagef**, **write_imagei** and **write_imageui** for image objects created with **image_channel_data_type** values not specified in the description above or with \((x, y)\) coordinate values that are not in the range \((0 \ldots \text{image width} – 1, 0 \ldots \text{image height} – 1)\), respectively, is undefined.

```c
void write_imagef(
    image2d_array_depth_t image,
    int4 coord,
    float depth)
```

Write depth value to location specified by \(coord.xy\) in the 2D image identified by \(coord.z\) in the 2D depth image array specified by \(image\). Appropriate data format conversion to the specified image format is done before writing the depth value. \(coord.x\), \(coord.y\) and \(coord.z\) are considered to be unnormalized coordinates and must be in the range \(0 \ldots \text{image width} – 1, 0 \ldots \text{image height} – 1\) and \(0 \ldots \text{image number of layers} – 1\).

**write_imagef** can only be used with image objects created with **image_channel_data_type** set to **CL_UNORM_INT16**, **CL_UNORM_INT24** or **CL_FLOAT**. Appropriate data format conversion will be done to convert depth value from a floating-point value to actual data format associated with the image.

The behavior of **write_imagef**, **write_imagei** and **write_imageui** for image objects created with **image_channel_data_type** values not specified in the description above or with \((x, y, z)\) coordinate values that are not in the range \((0 \ldots \text{image width} – 1, 0 \ldots \text{image height} – 1, 0 \ldots \text{image number of layers} – 1)\), respectively, is undefined.

```c
void write_imagef (image3d_t image, int4 coord, float4 color)
void write_imagei (image3d_t image, int4 coord, int4 color)
```

Write color value to location specified by \(coord.xyz\) in the 3D image object specified by \(image\). Appropriate data format conversion to the specified image format is done before writing the color value. \(coord.x\), \(coord.y\) and \(coord.z\) are considered to be unnormalized coordinates and must be in the range \(0 \ldots \text{image width} – 1, 0 \ldots \text{image height} – 1\) and \(0 \ldots \text{image depth} – 1\).
void write_imageui (image3d_t image,
    int4 coord,
    uint4 color)

write_imagef can only be used with image objects created with image_channel_data_type set to one of the pre-defined packed formats or set to
CL_SNORM_INT8, CL_UNORM_INT8,
CL_SNORM_INT16, CL_UNORM_INT16,
CL_HALF_FLOAT or CL_FLOAT. Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored.

write_imagei can only be used with image objects created with image_channel_data_type set to one of the following values:
CL_SIGNED_INT8,
CL_SIGNED_INT16 and
CL_SIGNED_INT32.

write_imageui can only be used with image objects created with image_channel_data_type set to one of the following values:
CL_UNSIGNED_INT8,
CL_UNSIGNED_INT16 and
CL_UNSIGNED_INT32.

The behavior of write_imagef, write_imagei and write_imageui for image objects with image_channel_data_type values not specified in the description above or with (x, y, z) coordinate values that are not in the range (0 … image width – 1, 0 … image height – 1, 0 … image depth – 1) respectively is undefined.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>int get_image_width (image1d_t image)</td>
<td>Return the image width in pixels.</td>
</tr>
<tr>
<td>int get_image_width (image1d_buffer_t image)</td>
<td></td>
</tr>
<tr>
<td>int get_image_width (image2d_t image)</td>
<td></td>
</tr>
</tbody>
</table>
### get_image_width

- `int get_image_width (image3d_t image)`
- `int get_image_width (image1d_array_t image)`
- `int get_image_width (image2d_array_t image)`
- `int get_image_width (image2d_depth_t image)`
- `int get_image_width (image2d_array_depth_t image)`

Return the image width in pixels.

### get_image_height

- `int get_image_height (image2d_t image)`
- `int get_image_height (image3d_t image)`
- `int get_image_height (image1d_array_t image)`
- `int get_image_height (image2d_array_t image)`
- `int get_image_height (image2d_depth_t image)`
- `int get_image_height (image2d_array_depth_t image)`

Return the image height in pixels.

### get_image_depth

- `int get_image_depth (image3d_t image)`

Return the image depth in pixels.

### get_image_channel_data_type

- `int get_image_channel_data_type (image1d_t image)`
- `int get_image_channel_data_type (image1d_buffer_t image)`
- `int get_image_channel_data_type (image2d_t image)`
- `int get_image_channel_data_type (image2d_depth_t image)`
- `int get_image_channel_data_type (image2d_array_depth_t image)`
- `int get_image_channel_data_type (image3d_t image)`
- `int get_image_channel_data_type (image1d_array_t image)`
- `int get_image_channel_data_type (image2d_array_t image)`

Return the channel data type. Valid values are:
- CLK_SNORM_INT8
- CLK_SNORM_INT16
- CLK_UNORM_INT8
- CLK_UNORM_INT16
- CLK_UNORM_SHORT_565
- CLK_UNORM_SHORT_555
- CLK_UNORM_SHORT_101010
- CLK_SIGNED_INT8
- CLK_SIGNED_INT16
- CLK_SIGNED_INT32
- CLK_UNSIGNED_INT8
- CLK_UNSIGNED_INT16
- CLK_UNSIGNED_INT32
- CLK_HALF_FLOAT
- CLK_FLOAT

### get_image_channel_order

- `int get_image_channel_order (image1d_t image)`
- `int get_image_channel_order (image1d_buffer_t image)`
- `int get_image_channel_order (image2d_array_t image)`

Return the image channel order. Valid values are:
- CLK_A
- CLK_R
### Table 6.26  Built-in Image Query Functions

The values returned by `get_image_channel_data_type` and `get_image_channel_order` as specified in table 6.26 with the CLK prefixes correspond to the CL prefixes used to describe the image channel order and data type in tables 5.4 and 5.5. For example, both `CL_UNORM_INT8` and `CLK_UNORM_INT8` refer to an image channel data type that is an unnormalized unsigned 8-bit integer.
6.13.14.6Reading and writing to the same image in a kernel

The atomic_work_item_fence(CLK_IMAGE_MEM_FENCE) built-in function can be used to make sure that sampler-less writes are visible to later reads by the same work-item. Only a scope of memory_order_acq_rel is valid for atomic_work_item_fence when passed the CLK_IMAGE_MEM_FENCE flag. If multiple work-items are writing to and reading from multiple locations in an image, the work_group_barrier(CLK_IMAGE_MEM_FENCE) should be used.

Consider the following example:

```c
kernel void foo(read_write image2d_t img, ... )
{
    int2 coord;
    coord.x = (int)get_global_id(0);
    coord.y = (int)get_global_id(1);

    float4 clr = read_imagef(img, coord);
    ...
    write_imagef(img, coord, clr);

    // required to ensure that following read from image at
    // location coord returns the latest color value.
    atomic_work_item_fence(CLK_IMAGE_MEM_FENCE,
                            memory_order_acq_rel,
                            memory_scope_work_item);

    float4 clr_new = read_imagef(img, coord);
    ...
}
```

6.13.14.7Mapping image channels to color values returned by read_image and color values passed to write_image to image channels

The following table describes the mapping of the number of channels of an image element to the appropriate components in the float4, int4 or uint4 vector data type for the color values returned by read_image{f|i|ui} or supplied to write_image{f|i|ui}. The unmapped components will be set to 0.0 for red, green and blue channels and will be set to 1.0 for the alpha channel.

<table>
<thead>
<tr>
<th>Channel Order</th>
<th>float4, int4 or uint4 components of channel data</th>
</tr>
</thead>
<tbody>
<tr>
<td>CL_R, CL_Rx</td>
<td>(r, 0.0, 0.0, 1.0)</td>
</tr>
<tr>
<td>CL_A</td>
<td>(0.0, 0.0, 0.0, a)</td>
</tr>
<tr>
<td>CL_RG, CL_RGx</td>
<td>(r, g, 0.0, 1.0)</td>
</tr>
<tr>
<td>----------------</td>
<td>-----------------</td>
</tr>
<tr>
<td>CL_RA</td>
<td>(r, 0.0, 0.0, a)</td>
</tr>
<tr>
<td>CL_RGB, CL_RGBx, CL_sRGB, CL_sRGBx</td>
<td>(r, g, b, 1.0)</td>
</tr>
<tr>
<td>CL_RGBA, CL_BGRA, CL_ARGB, CL_ABGR, CL_sRGBA, CL_sBGRA</td>
<td>(r, g, b, a)</td>
</tr>
<tr>
<td>CL_INTENSITY</td>
<td>(i, i, i, i)</td>
</tr>
<tr>
<td>CL_LUMINANCE</td>
<td>(L, L, L, 1.0)</td>
</tr>
</tbody>
</table>

For CL_DEPTH images, a scalar value is returned by `read_imagef` or supplied to `write_imagef`.

**NOTE:** A kernel that uses a sampler with the CL_ADDRESS_CLAMP addressing mode with multiple images may result in additional samplers being used internally by an implementation. If the same sampler is used with multiple images called via `read_image{f | i | ui}`, then it is possible that an implementation may need to allocate an additional sampler to handle the different border color values that may be needed depending on the image formats being used. These implementation allocated samplers will count against the maximum sampler values supported by the device and given by CL_DEVICE_MAX_SAMPLERS. Enqueuing a kernel that requires more samplers than the implementation can support will result in a CL_OUT_OF_RESOURCES error being returned.
6.13.15 Work-group Functions

The OpenCL C programming language implements the following built-in functions that operate on a work-group level. These built-in functions must be encountered by all work-items in a work-group executing the kernel. We use the generic type name gentype to indicate the built-in data types half\textsuperscript{59}, int, uint, long, ulong, float or double\textsuperscript{60} as the type for the arguments.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>int work_group_all (int predicate)</td>
<td>Evaluates predicate for all work-items in the work-group and returns a non-zero value if predicate evaluates to non-zero for all work-items in the work-group.</td>
</tr>
<tr>
<td>int work_group_any (int predicate)</td>
<td>Evaluates predicate for all work-items in the work-group and returns a non-zero value if predicate evaluates to non-zero for any work-items in the work-group.</td>
</tr>
<tr>
<td>gentype work_group_broadcast (</td>
<td>Broadcast the value of x for work-item identified by local_id to all work-items in the work-group.</td>
</tr>
<tr>
<td>gentype a,</td>
<td>local_id)</td>
</tr>
<tr>
<td>size_t local_id)</td>
<td></td>
</tr>
<tr>
<td>gentype work_group_broadcast (</td>
<td>local_id must be the same value for all work-items in the work-group.</td>
</tr>
<tr>
<td>gentype a,</td>
<td></td>
</tr>
<tr>
<td>size_t local_id_x,</td>
<td></td>
</tr>
<tr>
<td>size_t local_id_y)</td>
<td></td>
</tr>
<tr>
<td>gentype work_group_broadcast (</td>
<td></td>
</tr>
<tr>
<td>gentype a,</td>
<td></td>
</tr>
<tr>
<td>size_t local_id_x,</td>
<td></td>
</tr>
<tr>
<td>size_t local_id_y,</td>
<td></td>
</tr>
<tr>
<td>size_t local_id_z)</td>
<td></td>
</tr>
<tr>
<td>gentype work_group_reduce_&lt;op&gt; (</td>
<td>Return result of reduction operation specified by &lt;op&gt; for all values of x specified by work-items in a work-group.</td>
</tr>
<tr>
<td>gentype x)</td>
<td></td>
</tr>
<tr>
<td>gentype work_group_scan_exclusive_&lt;op&gt; (</td>
<td>Do an exclusive scan operation specified by &lt;op&gt; of all values specified by work-items in the work-group. The scan results are returned for each work-item.</td>
</tr>
<tr>
<td>gentype x)</td>
<td></td>
</tr>
</tbody>
</table>

\textsuperscript{59} Only if the cl_khr_fp16 extension is supported.

\textsuperscript{60} Only if double precision is supported.
linear global ID within the work-group.

<table>
<thead>
<tr>
<th>gentype</th>
<th>work_group_scan_inclusive_&lt;op&gt;(gentype x)</th>
<th>Do an inclusive scan operation specified by &lt;op&gt; of all values specified by work-items in the work-group. The scan results are returned for each work-item.</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td>The scan order is defined by increasing 1D linear global ID within the work-group.</td>
</tr>
</tbody>
</table>

Table 6.27  Built-in Work-group Functions

The <op> in work_group_reduce_<op>, work_group_scan_exclusive_<op> and work_group_scan_inclusive_<op> defines the operator and can be add, min or max.

The inclusive scan operation takes a binary operator op with an identity I and n (where n is the size of the work-group) elements \([a_0, a_1, \ldots, a_{n-1}]\) and returns \([a_0, (a_0 \text{ op } a_1), \ldots, (a_0 \text{ op } a_1 \text{ op } \ldots \text{ op } a_{n-1})]\). If \(<op> = \text{ add}\), the identity I is 0. If \(<op> = \text{ min}\), the identity I is INT_MAX, UINT_MAX, LONG_MAX, ULONG_MAX, for int, uint, long, ulong types and is +INF for floating-point types. Similarly if \(<op> = \text{ max}\), the identity I is INT_MIN, 0, LONG_MIN, 0 and -INF.

Consider the following example:

```c
void foo(int *p)
{
    ...
    int prefix_sum_val = work_group_scan_inclusive_add(
        p[get_local_id(0)]);
}
```

For the example above, let’s assume that the work-group size is 8 and p points to the following elements \([3 \ 1 \ 7 \ 0 \ 4 \ 1 \ 6 \ 3]\). Work-item 0 calls work_group_scan_inclusive_add with 3 and returns 3. Work-item 1 calls work_group_scan_inclusive_add with 1 and returns 4. The full set of values returned by work_group_scan_inclusive_add for work-items 0 … 7 are \([3 \ 4 \ 11 \ 11 \ 15 \ 16 \ 22 \ 25]\).

The exclusive scan operation takes a binary associative operator op with an identity I and n (where n is the size of the work-group) elements \([a_0, a_1, \ldots, a_{n-1}]\) and returns \([I, a_0, (a_0 \text{ op } a_1), \ldots, (a_0 \text{ op } a_1 \text{ op } \ldots \text{ op } a_{n-2})]\). For the example above, the exclusive scan add operation on the ordered set \([3 \ 1 \ 7 \ 0 \ 4 \ 1 \ 6 \ 3]\) would return \([0 \ 3 \ 4 \ 11 \ 11 \ 15 \ 16 \ 22]\).

NOTE: The order of floating-point operations is not guaranteed for the work_group_reduce_<op>, work_group_scan_inclusive_<op> and work_group_scan_exclusive_<op> built-in functions that operate on half, float and double data.
types. The order of these floating-point operations is also non-deterministic for a given work-group.
### 6.13.16 Pipe Functions

A pipe is identified by specifying the `pipe` keyword with a type. The data type specifies the size of each packet in the pipe. The `pipe` keyword is a type modifier. When it is applied to another type `T`, the result is a pipe type whose elements (or packets) are of type `T`. The packet type `T` may be any supported OpenCL C scalar and vector integer or floating-point data types, or a user-defined type built from these scalar and vector data types.

Examples:

```c
pipe int4 pipeA; // a pipe with int4 packets
pipe user_type_t pipeB; // a pipe with user_type_t packets
```

The `read_only` (or `__read_only`) and `write_only` (or `__write_only`) qualifiers can also be used with the `pipe` qualifier to identify if a pipe can be read from or written to by a kernel and its callees and enqueued child kernels. The default qualifier is `read_only`.

A kernel cannot read from and write to the same pipe object. Using the `read_write` (or `__read_write`) qualifier with the `pipe` qualifier is a compilation error.

In the following example

```c
kernel void foo (read_only pipe fooA_t pipeA,
    write_only pipe fooB_t pipeB)
{
    ...
}
```

`pipeA` is a read-only pipe object, and `pipeB` is a write-only pipe object.

The macro `CLK_NULL_RESERVE_ID` refers to an invalid reservation ID.

#### 6.13.16.1 Restrictions

- Pipes can only be passed as arguments to a function (including kernel functions). The C operators (refer to section 6.3 of the OpenCL 2.0 specification) cannot be used with variables declared with the `pipe` qualifier.

- The `pipe` qualifier cannot be used with variables declared inside a kernel, a structure or union field, a pointer type, an array, global variables declared in program scope or the return type of a function.
### 6.13.16.2 Built-in Pipe Read and Write Functions

The OpenCL C programming language implements the following built-in functions that read from or write to a pipe. We use the generic type name `gentype` to indicate the built-in OpenCL C scalar or vector integer or floating-point data types or any user defined type built from these scalar and vector data types can be used as the type for the arguments to the pipe functions listed in Table 6.28.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>int read_pipe</strong> (pipe gentype <em>p</em>, gentype *ptr)</td>
<td>Read packet from pipe <em>p</em> into <em>ptr</em>. Returns 0 if <strong>read_pipe</strong> is successful and a negative value if the pipe is empty.</td>
</tr>
<tr>
<td><strong>int write_pipe</strong> (pipe gentype <em>p</em>, const gentype *ptr)</td>
<td>Write packet specified by <em>ptr</em> to pipe <em>p</em>. Returns 0 if <strong>write_pipe</strong> is successful and a negative value if the pipe is full.</td>
</tr>
<tr>
<td><strong>int read_pipe</strong> (pipe gentype <em>p</em>, reserve_id_t <em>reserve_id</em>, uint index, gentype *ptr)</td>
<td>Read packet from the reserved area of the pipe referred to by <em>reserve_id</em> and <em>index</em> into <em>ptr</em>. The reserved pipe entries are referred to by indices that go from 0 … num_packets – 1. Returns 0 if <strong>read_pipe</strong> is successful and a negative value otherwise.</td>
</tr>
<tr>
<td><strong>int write_pipe</strong> (pipe gentype <em>p</em>, reserve_id_t <em>reserve_id</em>, uint index, const gentype *ptr)</td>
<td>Write packet specified by <em>ptr</em> to the reserved area of the pipe referred to by <em>reserve_id</em> and <em>index</em>. The reserved pipe entries are referred to by indices that go from 0 … num_packets – 1. Returns 0 if <strong>write_pipe</strong> is successful and a negative value otherwise.</td>
</tr>
<tr>
<td><strong>reserve_id_t reserve_read_pipe</strong> (pipe gentype <em>p</em>, uint num_packets)</td>
<td>Reserve num_packets entries for reading from or writing to pipe <em>p</em>. Returns a valid reservation ID if the reservation is successful.</td>
</tr>
<tr>
<td><strong>reserve_id_t reserve_write_pipe</strong> (pipe gentype <em>p</em>, uint num_packets)</td>
<td></td>
</tr>
<tr>
<td><strong>void commit_read_pipe</strong> (pipe gentype <em>p</em>,)</td>
<td>Indicates that all reads and writes to num_packets associated with reservation <em>reserve_id</em> are</td>
</tr>
</tbody>
</table>

---

61 The half scalar and vector types can only be used if the **cl_khr_fp16** extension is supported. The double scalar and vector types can only be used if double precision is supported.
### 6.13.16.3 Built-in Work-group Pipe Read and Write Functions

The OpenCL C programming language implements the following built-in pipe functions that operate at a work-group level. These built-in functions must be encountered by all work-items in a work-group executing the kernel with the same argument values; otherwise the behavior is undefined. We use the generic type name `gentype` to indicate the built-in OpenCL C scalar or vector integer or floating-point data types or any user-defined type built from these scalar and vector data types can be used as the type for the arguments to the pipe functions listed in table 6.29.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>reserve_id_t reserve_id_t reserve_id_t</code></td>
<td><code>work_group_reserve_read_pipe (pipe gentype p, uint num_packets)</code> Reserve <code>num_packets</code> entries for reading from or writing to pipe <code>p</code>. Returns a valid reservation ID if the reservation is successful. The reserved pipe entries are referred to by indices that go from 0 … <code>num_packets</code> – 1.</td>
</tr>
<tr>
<td><code>reserve_id_t reserve_id_t</code></td>
<td><code>work_group_reserve_write_pipe (pipe gentype p, uint num_packets)</code></td>
</tr>
<tr>
<td><code>void work_group_commit_read_pipe (pipe gentype p, reserve_id_t reserve_id)</code></td>
<td>Indicates that all reads and writes to <code>num_packets</code> associated with reservation <code>reserve_id</code> are completed.</td>
</tr>
<tr>
<td><code>void work_group_commit_write_pipe (pipe gentype p, reserve_id_t reserve_id)</code></td>
<td></td>
</tr>
</tbody>
</table>

Table 6.29  Built-in Pipe Work-group Functions

---

62 The half scalar and vector types can only be used if the `cl_khr_fp16` extension is supported. The double scalar and vector types can only be used if double precision is supported.
NOTE: The read_pipe and write_pipe functions that take a reservation ID as an argument can be used to read from or write to a packet index. These built-ins can be used to read from or write to a packet index one or multiple times. If a packet index that is reserved for writing is not written to using the write_pipe function, the contents of that packet in the pipe are undefined. commit_read_pipe and work_group_commit_read_pipe remove the entries reserved for reading from the pipe. commit_write_pipe and work_group_commit_write_pipe ensures that the entries reserved for writing are all added in-order as one contiguous set of packets to the pipe.

There can only be CL_DEVICE_PIPE_MAX_ACTIVE_RESERVATIONS (refer to table 4.3) reservations active (i.e. reservation IDs that have been reserved but not committed) per work-item or work-group for a pipe in a kernel executing on a device.

Work-item based reservations made by a work-item are ordered in the pipe as they are ordered in the program. Reservations made by different work-items that belong to the same work-group can be ordered using the work-group barrier function. The order of work-item based reservations that belong to different work-groups is implementation defined.

Work-group based reservations made by a work-group are ordered in the pipe as they are ordered in the program. The order of work-group based reservations by different work-groups is implementation defined.

### 6.13.16.4 Built-in Pipe Query Functions

The OpenCL C programming language implements the following built-in query functions for a pipe. We use the generic type name gentype to indicate the built-in OpenCL C scalar or vector integer or floating-point data types or any user defined type built from these scalar and vector data types or any user defined type built from these scalar and vector data types can be used as the type for the arguments to the pipe functions listed in table 6.30.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>uint get_pipe_num_packets (pipe gentype p)</td>
<td>Returns the number of available entries in the pipe. The number of available entries in a pipe is a dynamic value. The value returned should be considered immediately stale.</td>
</tr>
<tr>
<td>uint get_pipe_max_packets (pipe gentype p)</td>
<td>Returns the maximum number of packets specified when pipe was created.</td>
</tr>
</tbody>
</table>

| Table 6.30 | Built-in Pipe Query Functions |

---

63 The half scalar and vector types can only be used if the cl_khr_fp16 extension is supported. The double scalar and vector types can only be used if double precision is supported.
6.13.16.5 Restrictions

The following behavior is undefined:

- A kernel fails to call `reserve_pipe` before calling `read_pipe` or `write_pipe` that take a reservation ID.
- A kernel calls `read_pipe`, `write_pipe`, `commit_read_pipe` or `commit_write_pipe` with an invalid reservation ID.
- A kernel calls `read_pipe` or `write_pipe` with an valid reservation ID but with an index that is not a value from 0 … `num_packets` – 1 specified to the corresponding call to `reserve_pipe`.
- A kernel calls `read_pipe` or `write_pipe` with a reservation ID that has already been committed (i.e. a `commit_read_pipe` or `commit_write_pipe` with this reservation ID has already been called).
- A kernel fails to call `commit_read_pipe` for any reservation ID obtained by a prior call to `reserve_read_pipe`.
- A kernel fails to call `commit_write_pipe` for any reservation ID obtained by a prior call to `reserve_write_pipe`.
- The contents of the reserved data packets in the pipe are undefined if the kernel does not call `write_pipe` for all entries that were reserved by the corresponding call to `reserve_pipe`.
- Calls to `read_pipe` that takes a reservation ID and `commit_read_pipe` or `write_pipe` that takes a reservation ID and `commit_write_pipe` for a given reservation ID must be called by the same kernel that made the reservation using `reserve_read_pipe` or `reserve_write_pipe`. The reservation ID cannot be passed to another kernel including child kernels.
6.13.17 Enqueuing Kernels

OpenCL 2.0 allows a kernel to independently enqueue to the same device, without host interaction. A kernel may enqueue code represented by Block syntax, and control execution order with event dependencies including user events and markers. There are several advantages to using the Block syntax: it is more compact; it does not require a cl_kernel object; and enqueuing can be done as a single semantic step.

The following table describes the list of built-in functions that can be used to enqueue a kernel(s).

The macro CLK_NULL_EVENT refers to an invalid device event. The macro CLK_NULL_QUEUE refers to an invalid device queue.

### 6.13.17.1 Built-in Functions – Enqueuing a kernel

<table>
<thead>
<tr>
<th>Built-in Function</th>
<th>Description</th>
</tr>
</thead>
</table>
| int enqueue_kernel (  
    queue_t queue,  
    kernel_enqueue_flags_t flags,  
    const ndrange_t ndrange,  
    void (^block)(void)) | Enqueue the block for execution to queue.  
If an event is returned, enqueue_kernel performs an implicit retain on the returned event. |

| int enqueue_kernel (  
    queue_t queue,  
    kernel_enqueue_flags_t flags,  
    const ndrange_t ndrange,  
    u_int num_events_in_wait_list,  
    const clk_event_t *event_wait_list,  
    clk_event_t *event_ret,  
    void (^block)(void)) | |

| int enqueue_kernel (  
    queue_t queue,  
    kernel_enqueue_flags_t flags,  
    const ndrange_t ndrange,  
    void (^block)(local void *, ...),  
    u_int size0, ...) | |

| int enqueue_kernel (  
    queue_t queue,  
    kernel_enqueue_flags_t flags,  
    const ndrange_t ndrange,  
    u_int num_events_in_wait_list, | |
Table 6.31  Built-in Kernel Enqueue Functions

The **enqueue_kernel** built-in function allows a work-item to enqueue a block. Work-items can enqueue multiple blocks to a device queue(s).

The **enqueue_kernel** built-in function returns CL_SUCCESS if the block is enqueued successfully and returns CL_ENQUEUE_FAILURE otherwise. If the –g compile option is specified in compiler options passed to **clCompileProgram** or **clBuildProgram** when compiling or building the parent program, the following errors may be returned instead of CL_ENQUEUE_FAILURE to indicate why **enqueue_kernel** failed to enqueue the block:

- **CLK_INVALID_QUEUE** if queue is not a valid device queue.
- **CLK_INVALID_NDRANGE** if ndrange is not a valid ND-range descriptor or if the program was compiled with –cl-uniform-work-group-size and the local_work_size is specified in ndrange but the global_work_size specified in ndrange is not a multiple of the local_work_size.
- **CLK_INVALID_EVENT_WAIT_LIST** if event_wait_list is NULL and num_events_in_wait_list > 0, or if event_wait_list is not NULL and num_events_in_wait_list is 0, or if event objects in event_wait_list are not valid events.
- **CLK_DEVICE_QUEUE_FULL** if queue is full.
- **CLK_INVALID_ARG_SIZE** if size of local memory arguments is 0.
- **CLK_EVENT_ALLOCATION_FAILURE** if event_ret is not NULL and an event could not be allocated.
- **CLK_OUT_OF_RESOURCES** if there is a failure to queue the block in queue because of insufficient resources needed to execute the kernel.

Below are some examples of how to enqueue a block.

```c
kernel void
my_func_A(global int *a, global int *b, global int *c)
{
    ...
}
```
kernel void
my_func_B(global int *a, global int *b, global int *c)
{
    ndrange_t ndrange;
    // build ndrange information
    ...

    // example – enqueue a kernel as a block
    enqueue_kernel(get_default_queue(), ndrange,
        ^{my_func_A(a, b, c);});
    ...
}

kernel void
my_func_C(global int *a, global int *b, global int *c)
{
    ndrange_t ndrange;
    // build ndrange information
    ...

    // note that a, b and c are variables in scope of
    // the block
    void (^my_block_A)(void) = ^{my_func_A(a, b, c);};

    // enqueue the block variable
    enqueue_kernel(get_default_queue(),
        CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
        ndrange,
        my_block_A);
    ...
}

The example below shows how to declare a block literal and enqueue it.

kernel void
my_func(global int *a, global int *b)
{
    ndrange_t ndrange;
    // build ndrange information
    ...

    // note that a, b and c are variables in scope of
    // the block
    void (^my_block_A)(void) = ^{ size_t id = get_global_id(0);
        b[id] += a[id];
    }
};

// enqueue the block variable
enqueue_kernel(get_default_queue(),
    CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
    ndrange,
    my_block_A);

// or we could have done the following
enqueue_kernel(get_default_queue(),
    CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
    ndrange,
    ^{
        size_t id = get_global_id(0);
        b[id] += a[id];
    });

NOTE: Blocks passed to enqueue_kernel cannot use global variables or stack variables local to the enclosing lexical scope that are a pointer type in the local or private address space.

Example:

kernel void foo(global int *a, local int *lptr, ...)
{
    enqueue_kernel(get_default_queue(),
        CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
        ndrange,
        ^{
            size_t id = get_global_id(0);
            local int *p = lptr; // undefined behavior
        });
}

6.13.17.2 Arguments that are a pointer type to local address space

A block passed to enqueue_kernel can have arguments declared to be a pointer to local memory. The enqueue_kernel built-in function variants allow blocks to be enqueued with a variable number of arguments. Each argument must be declared to be a pointer of a data type to local memory. These enqueue_kernel built-in function variants also have a corresponding number of arguments each of type uint that follow the block argument. These arguments specify the size of each local memory pointer argument of the enqueued block.

Some examples follow:
kernel void
my_func_A_local_arg1(global int *a, local int *lptr, ...)
{
    ...
}

kernel void
my_func_A_local_arg2(global int *a,
                      local int *lptr1, local float4 *lptr2, ...)
{
    ...
}

kernel void
my_func_B(global int *a, ...)
{
    ...
    ndrange_t ndrange = ndrange_1d(...);

    uint local_mem_size = compute_local_mem_size();

    enqueue_kernel(get_default_queue(),
                   CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
                   ndrange,
                   ^(local int *p){my_func_A_local_arg1(a, p, ...);},
                   local_mem_size);
}

kernel void
my_func_C(global int *a, ...)
{
    ...
    ndrange_t ndrange = ndrange_1d(...);

    void (^my_blk_A)(local int *, local float4 *) =
                   ^(local int *lptr1, local float4 *lptr2)
                   {my_func_A_local_arg2(a, lptr1, lptr2, ...);};

    // calculate local memory size for lptr
    // argument in local address space for my_blk_A
    uint local_mem_size = compute_local_mem_size();

    enqueue_kernel(get_default_queue(),
                   CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
                   ndrange,
                   my_blk_A,
local_mem_size, local_mem_size*4);
}

6.13.17.3 A Complete Example

The example below shows how to implement an iterative algorithm where the host enqueues the first instance of the nd-range kernel (dp_func_A). The kernel dp_func_A will launch a kernel (evaluate_dp_work_A) that will determine if new nd-range work needs to be performed. If new nd-range work does need to be performed, then evaluate_dp_work_A will enqueue a new instance of dp_func_A. This process is repeated until all the work is completed.

```
kernelpvoid
dp_func_A(queue_t q, ...)
{
  ...

  // queue a single instance of evaluate_dp_work_A to
  // device queue q. queued kernel begins execution after
  // kernel dp_func_A finishes

  if (get_global_id(0) == 0)
  {
    enqueue_kernel(q,
                    CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
                    ndrange_1d(1),
                    ^{evaluate_dp_work_A(q, ...);});
  }
}

kernelpvoid
evaluate_dp_work_A(queue_t q, ...)
{
  // check if more work needs to be performed
  bool more_work = check_new_work(...);
  if (more_work)
  {
    size_t global_work_size = compute_global_size(...);

    void (^dp_func_A_blk)(void) =
      ^{dp_func_A(q, ...)};

    // get local WG-size for kernel dp_func_A
    size_t local_work_size =
```

Last Revision Date: 11/1/14
get_kernel_work_group_size(dp_func_A_blk);

// build nd-range descriptor
ndrange_t ndrange = ndrange_1D(global_work_size,
                                local_work_size);

// enqueue dp_func_A
enqueue_kernel(q,
               CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
               ndrange,
               dp_func_A_blk);
}

...}

6.13.17.4 Determining when a child kernel begins execution

The kernel_enqueue_flags_t\textsuperscript{64} argument to enqueue_kernel built-in functions can be used to specify when the child kernel begins execution. Supported values are described in the table below:

<table>
<thead>
<tr>
<th>kernel_enqueue_flags_t enum</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>CLK_ENQUEUE_FLAGS_NO_WAIT</td>
<td>Indicates that the enqueued kernels do not need to wait for the parent kernel to finish execution before they begin execution.</td>
</tr>
<tr>
<td>CLK_ENQUEUE_FLAGS_WAIT_KERNEL</td>
<td>Indicates that all work-items of the parent kernel must finish executing and all immediate\textsuperscript{65} side effects committed before the enqueued child kernel may begin execution.</td>
</tr>
<tr>
<td>CLK_ENQUEUE_FLAGS_WAIT_WORK_GROUP\textsuperscript{66}</td>
<td>Indicates that the enqueued kernels wait only for the workgroup that enqueued the kernels to finish before they begin execution.</td>
</tr>
</tbody>
</table>

Table 6.32 Kernel Enqueue Flags

\textsuperscript{64} Implementations are not required to honor this flag. Implementations may not schedule kernel launch earlier than the point specified by this flag, however.

\textsuperscript{65} Immediate meaning not side effects resulting from child kernels. The side effects would include stores to global memory and pipe reads and writes.

\textsuperscript{66} This acts as a memory synchronization point between work-items in a work-group and child kernels enqueued by work-items in the work-group.
NOTE: The `kernel_enqueue_flags_t` flags are useful when a kernel enqueued from the host and executing on a device enqueues kernels on the device. The kernel enqueued from the host may not have an event associated with it. The `kernel_enqueue_flags_t` flags allow the developer to indicate when the child kernels can begin execution.

### 6.13.17.5 Determining when a parent kernel has finished execution

A parent kernel’s execution status is considered to be complete when it and all its child kernels have finished execution. The execution status of a parent kernel will be `CL_COMPLETE` if this kernel and all its child kernels finish execution successfully. The execution status of the kernel will be an error code (given by a negative integer value) if it or any of its child kernels encounter an error, or are abnormally terminated.

For example, assume that the host enqueues a kernel `k` for execution on a device. Kernel `k` when executing on the device enqueues kernels `A` and `B` to a device queue(s). The `enqueue_kernel` call to enqueue kernel `B` specifies the event associated with kernel `A` in the `event_wait_list` argument i.e. wait for kernel `A` to finish execution before kernel `B` can begin execution. Let’s assume kernel `A` enqueues kernels `X, Y and Z`. Kernel `A` is considered to have finished execution i.e. its execution status is `CL_COMPLETE` only after `A` and the kernels `A` enqueued (and any kernels these enqueued kernels enqueue and so on) have finished execution.

### 6.13.17.6 Built-in Functions – Kernel Query Functions

<table>
<thead>
<tr>
<th>Built-in Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>uint get_kernel_work_group_size ( void (^block)(void));</code></td>
<td>This provides a mechanism to query the maximum work-group size that can be used to execute a block on a specific device given by <code>device</code>. <code>block</code> specifies the block to be enqueued.</td>
</tr>
<tr>
<td><code>uint get_kernel_work_group_size ( void (^block)(local void *, ...));</code></td>
<td></td>
</tr>
<tr>
<td><code>uint get_kernel_preferred_work_group_size_multiple ( void (^block)(void));</code></td>
<td>Returns the preferred multiple of work-group size for launch. This is a performance hint. Specifying a work-group size that is not a multiple of the value returned by this query as the value of the local work size argument to <code>enqueue_kernel</code> will not fail to enqueue the block for execution unless the work-group size specified is larger than the device maximum.</td>
</tr>
<tr>
<td><code>uint get_kernel_preferred_work_group_size_multiple ( void (^block)(local void *, ...));</code></td>
<td></td>
</tr>
</tbody>
</table>

**Table 6.33**  Built-in Kernel Query Functions
6.13.17.7 Built-in Functions – Queuing other commands

The following table describes the list of built-in functions that can be used to enqueue commands such as a marker.

<table>
<thead>
<tr>
<th>Built-in Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>int enqueue_marker (</td>
<td>Enqueue a marker command to queue.</td>
</tr>
<tr>
<td>queue_t queue,</td>
<td>The marker command waits for a list of events specified by event_wait_list</td>
</tr>
<tr>
<td>uint num_events_in_wait_list,</td>
<td>to complete before the marker completes.</td>
</tr>
<tr>
<td>const clk_event_t *event_wait_list,</td>
<td>event_wait_list must not be NULL as otherwise this is a no-op.</td>
</tr>
<tr>
<td>clk_event_t *event_ret)</td>
<td>If an event is returned, enqueue_marker performs an implicit retain on the returned event.</td>
</tr>
</tbody>
</table>

Table 6.34  Built-in Other Enqueue Functions

The enqueue_marker built-in function returns CLK_SUCCESS if the marked command is enqueued successfully and returns CLK_ENQUEUE_FAILURE otherwise. If the –g compile option is specified in compiler options passed to `clCompileProgram` or `clBuildProgram`, the following errors may be returned instead of CLK_ENQUEUE_FAILURE to indicate why enqueue_marker failed to enqueue the marker command:

- CLK_INVALID_QUEUE if `queue` is not a valid device queue.
- CLK_INVALID_EVENT_WAIT_LIST if `event_wait_list` is NULL, or if `event_wait_list` is not NULL and `num_events_in_wait_list` is 0, or if event objects in `event_wait_list` are not valid events.
- CLK_DEVICE_QUEUE_FULL if `queue` is full.
- CLK_EVENT_ALLOCATION_FAILURE if `event_ret` is not NULL and an event could not be allocated.
- CLK_OUT_OF_RESOURCES if there is a failure to queue the block in `queue` because of insufficient resources needed to execute the kernel.

6.13.17.8 Built-in Functions – Event Functions

The following table describes the list of built-in functions that work on events.
<table>
<thead>
<tr>
<th>Built-in Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>void <strong>retain_event</strong> (clk_event_t event)</td>
<td>Increments the event reference count. <strong>event</strong> must be an event returned by <code>enqueue_kernel</code> or <code>enqueue_marker</code> or a user event.</td>
</tr>
<tr>
<td>void <strong>release_event</strong> (clk_event_t event)</td>
<td>Decrements the event reference count. The event object is deleted once the event reference count is zero, the specific command identified by this event has completed (or terminated) and there are no commands in any device command queue that require a wait for this event to complete. <strong>event</strong> must be an event returned by <code>enqueue_kernel</code>, <code>enqueue_marker</code> or a user event.</td>
</tr>
<tr>
<td><strong>clk_event_t create_user_event</strong> ()</td>
<td>Create a user event. Returns the user event. The execution status of the user event created is set to CL_SUBMITTED.</td>
</tr>
<tr>
<td><strong>bool is_valid_event</strong> (clk_event_t event)</td>
<td>Returns true if <strong>event</strong> is a valid event. Otherwise returns false.</td>
</tr>
<tr>
<td>void <strong>set_user_event_status</strong> (clk_event_t event, int status)</td>
<td>Sets the execution status of a user event. <strong>event</strong> must be a user-event. <strong>status</strong> can be either CL_COMPLETE or a negative integer value indicating an error.</td>
</tr>
<tr>
<td>void <strong>capture_event_profiling_info</strong> (clk_event_t event, clk_profiling_info name, global void *value)</td>
<td>Captures the profiling information for functions that are enqueued as commands. The specific function being referred to is: <code>enqueue_kernel</code>. These enqueued commands are identified by unique event objects. The profiling information will be available in <strong>value</strong> once the command identified by <strong>event</strong> has completed. <strong>event</strong> must be an event returned by <code>enqueue_kernel</code>. <strong>name</strong> identifies which profiling information is to be queried and can be: CLK_PROFILING_COMMAND_EXEC_TIME <strong>value</strong> is a pointer to two 64-bit values.</td>
</tr>
</tbody>
</table>
The first 64-bit value describes the elapsed time 
CL_PROFILING_COMMAND_END – 
CL_PROFILING_COMMAND_START for the 
command identified by event in nanoseconds.

The second 64-bit value describes the elapsed 
time CL_PROFILING_COMMAND_COMPLETE – CL_PROFILING_COMMAND_START for the 
command identified by event in nanoseconds.

NOTE: The behavior of 
capture_event_profiling_info when 
called multiple times for the same event is 
undefined.

Table 6.35  Built-in Event Functions

Events can be used to identify commands enqueued to a command-queue from the host. These 
events created by the OpenCL runtime can only be used on the host i.e. as events passed in 

**event_wait_list** argument to various clEnqueue APIs or runtime APIs that take events as 
arguments such as clRetainEvent, clReleaseEvent, clGetEventProfilingInfo.

Similarly, events can be used to identify commands enqueued to a device queue (from a kernel). 
These event objects cannot be passed to the host or used by OpenCL runtime APIs such as the 
cEnqueueAPIs or runtime APIs that take event arguments.

**clRetainEvent** and **clReleaseEvent** will return CL_INVALID_OPERATION if event 
specified is an event that refers to any kernel enqueued to a device queue using 
enqueue_kernel or enqueue_marker or is a user event created by 
create_user_event.

Similarly, **clSetUserEventStatus** can only be used to set the execution status of events 
created using clCreateUserEvent. User events created on the device can be set using 
set_user_event_status built-in function.

The example below shows how events can be used with kernels enqueued to multiple device 
queues.

```c
extern void barA_kernel(...);
extern void barB_kernel(...);

kernel void
foo(queue_t q0, queue q1, ...) {
  ...
```
clk_event_t evt0;

// enqueue kernel to queue q0
enqueue_kernel(q0,
CLK_ENQUEUE_FLAGS_NO_WAIT,
ndrange_A,
0, NULL, &evt0,
^{barA_kernel(...);} );

// enqueue kernel to queue q1
enqueue_kernel(q1,
CLK_ENQUEUE_FLAGS_NO_WAIT,
ndrange_B,
1, &evt0, NULL,
^{barB_kernel(...);} );

// release event evt0. This will get released
// after barA_kernel enqueued in queue q0 has finished
// execution and barB_kernel enqueued in queue q1 and
// waits for evt0 is submitted for execution i.e. wait
// for evt0 is satisfied.
release_event(evt0);
}

The example below shows how the marker command can be used with kernels enqueued to a device queue.

kernel void
foo(queue_t q, ...)
{
...
clk_event_t marker_event;
clk_event_t events[2];

enqueue_kernel(q,
CLK_ENQUEUE_FLAGS_NO_WAIT,
ndrange,
0, NULL, &events[0],
^{barA_kernel(...);} );

enqueue_kernel(q,
CLK_ENQUEUE_FLAGS_NO_WAIT,
ndrange,
0, NULL, &events[1],
^{barB_kernel(...);} );

// barA_kernel and barB_kernel can be executed
// out of order. we need to wait for both these
// kernels to finish execution before barC_kernel
// starts execution so we enqueue a marker command and
// then enqueue barC_kernel that waits on the event
// associated with the marker.

enqueue_marker(q, 2, events, &marker_event);

enqueue_kernel(q,
    CLK_ENQUEUE_FLAGS_NO_WAIT,
    1, &marker_event, NULL,
    ^{barC_kernel(...);} );

release_event(events[0];
release_event(events[1]);
release_event(marker_event);
}

6.13.17.9 Built-in Functions – Helper Functions

<table>
<thead>
<tr>
<th>Built-in Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>queue_t get_default_queue (void)</td>
<td>Returns the default device queue. If a default device queue has not been created, CLK_NULL_QUEUE is returned.</td>
</tr>
<tr>
<td>ndrange_t ndrange_1D (</td>
<td>Builds a 1D, 2D or 3D ND-range descriptor.</td>
</tr>
<tr>
<td>size_t global_work_size)</td>
<td></td>
</tr>
<tr>
<td>ndrange_t ndrange_1D (</td>
<td></td>
</tr>
<tr>
<td>size_t global_work_size, size_t local_work_size)</td>
<td></td>
</tr>
<tr>
<td>ndrange_t ndrange_1D (</td>
<td></td>
</tr>
<tr>
<td>size_t global_work_offset, size_t global_work_size, size_t local_work_size)</td>
<td></td>
</tr>
<tr>
<td>ndrange_t ndrange_2D (</td>
<td></td>
</tr>
<tr>
<td>const size_t global_work_size[2])</td>
<td></td>
</tr>
<tr>
<td>ndrange_t ndrange_2D (</td>
<td></td>
</tr>
<tr>
<td>const size_t global_work_size[2], const size_t local_work_size[2])</td>
<td></td>
</tr>
<tr>
<td>ndrange_t ndrange_2D (</td>
<td></td>
</tr>
<tr>
<td>const size_t global_work_offset[2],</td>
<td></td>
</tr>
<tr>
<td>const size_t global_work_size[2],</td>
<td></td>
</tr>
<tr>
<td>const size_t local_work_size[2])</td>
<td></td>
</tr>
<tr>
<td>ndrange_t ndrange_3D (</td>
<td></td>
</tr>
<tr>
<td>const size_t global_work_size[3])</td>
<td></td>
</tr>
<tr>
<td>ndrange_t ndrange_3D (</td>
<td></td>
</tr>
<tr>
<td>const size_t global_work_size[3],</td>
<td></td>
</tr>
<tr>
<td>const size_t local_work_size[3])</td>
<td></td>
</tr>
<tr>
<td>ndrange_t ndrange_3D (</td>
<td></td>
</tr>
<tr>
<td>const size_t global_work_offset[3],</td>
<td></td>
</tr>
<tr>
<td>const size_t global_work_size[3],</td>
<td></td>
</tr>
<tr>
<td>const size_t local_work_size[3])</td>
<td></td>
</tr>
</tbody>
</table>

Table 6.36  Built-in Helper Functions
7. **OpenCL Numerical Compliance**

This section describes features of the C99 and IEEE 754 standards that must be supported by all OpenCL compliant devices.

This section describes the functionality that must be supported by all OpenCL devices for single precision floating-point numbers. Currently, only single precision floating-point is a requirement. Double precision floating-point is an optional feature.

### 7.1 Rounding Modes

Floating-point calculations may be carried out internally with extra precision and then rounded to fit into the destination type. IEEE 754 defines four possible rounding modes:

- **Round to nearest even**
- **Round toward +∞**
- **Round toward -∞**
- **Round toward zero**

*Round to nearest even* is currently the only rounding mode required\(^{67}\) by the OpenCL specification for single precision and double precision operations and is therefore the default rounding mode. In addition, only static selection of rounding mode is supported. Dynamically reconfiguring the rounding modes as specified by the IEEE 754 spec is unsupported.

### 7.2 INF, NaN and Denormalized Numbers

INF and NaNs must be supported. Support for signaling NaNs is not required.

Support for denormalized numbers with single precision floating-point is optional. Denormalized single precision floating-point numbers passed as input or produced as the output of single precision floating-point operations such as add, sub, mul, divide, and the functions defined in sections 6.11.2 (math functions), 6.11.4 (common functions) and 6.11.5 (geometric functions) may be flushed to zero.

---

\(^{67}\) Except for the embedded profile whether either round to zero or round to nearest rounding mode may be supported for single precision floating-point.
### 7.3 Floating-Point Exceptions

Floating-point exceptions are disabled in OpenCL. The result of a floating-point exception must match the IEEE 754 spec for the exceptions not enabled case. Whether and when the implementation sets floating-point flags or raises floating-point exceptions is implementation-defined. This standard provides no method for querying, clearing or setting floating-point flags or trapping raised exceptions. Due to non-performance, non-portability of trap mechanisms and the impracticality of servicing precise exceptions in a vector context (especially on heterogeneous hardware), such features are discouraged.

Implementations that nevertheless support such operations through an extension to the standard shall initialize with all exception flags cleared and the exception masks set so that exceptions raised by arithmetic operations do not trigger a trap to be taken. If the underlying work is reused by the implementation, the implementation is however not responsible for reclearing the flags or resetting exception masks to default values before entering the kernel. That is to say that kernels that do not inspect flags or enable traps are licensed to expect that their arithmetic will not trigger a trap. Those kernels that do examine flags or enable traps are responsible for clearing flag state and disabling all traps before returning control to the implementation. Whether or when the underlying work-item (and accompanying global floating-point state if any) is reused is implementation-defined.

The expressions `math_errorhandling` and `MATH_ERREXCEPT` are reserved for use by this standard, but not defined. Implementations that extend this specification with support for floating-point exceptions shall define `math_errorhandling` and `MATH_ERREXCEPT` per ISO / IEC 9899 : TC2.

### 7.4 Relative Error as ULPs

In this section we discuss the maximum relative error defined as `ulp` (units in the last place). Addition, subtraction, multiplication, fused multiply-add and conversion between integer and a single precision floating-point format are IEEE 754 compliant and are therefore correctly rounded. Conversion between floating-point formats and explicit conversions specified in section 6.2.3 must be correctly rounded.

The ULP is defined as follows:

If \( x \) is a real number that lies between two finite consecutive floating-point numbers \( a \) and \( b \), without being equal to one of them, then \( \text{ulp}(x) = |b - a| \), otherwise \( \text{ulp}(x) \) is the distance between the two non-equal finite floating-point numbers nearest \( x \). Moreover, \( \text{ulp}(\text{NaN}) \) is \( \text{NaN} \).
Table 7.1\textsuperscript{68} describes the minimum accuracy of single precision floating-point arithmetic operations given as ULP values. The reference value used to compute the ULP value of an arithmetic operation is the infinitely precise result.

<table>
<thead>
<tr>
<th>Function</th>
<th>Min Accuracy - ULP values\textsuperscript{69}</th>
</tr>
</thead>
<tbody>
<tr>
<td>$x + y$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$x - y$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$x \times y$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$1.0 / x$</td>
<td>$\leq 2.5$ ulp</td>
</tr>
<tr>
<td>$x / y$</td>
<td>$\leq 2.5$ ulp</td>
</tr>
<tr>
<td>$\arccos$</td>
<td>$\leq 4$ ulp</td>
</tr>
<tr>
<td>$\arccos\pi$</td>
<td>$\leq 5$ ulp</td>
</tr>
<tr>
<td>$\arcsin$</td>
<td>$\leq 4$ ulp</td>
</tr>
<tr>
<td>$\arcsin\pi$</td>
<td>$\leq 5$ ulp</td>
</tr>
<tr>
<td>$\arctan$</td>
<td>$\leq 5$ ulp</td>
</tr>
<tr>
<td>$\arctan\pi$</td>
<td>$\leq 5$ ulp</td>
</tr>
<tr>
<td>$\arctan2\pi$</td>
<td>$\leq 6$ ulp</td>
</tr>
<tr>
<td>$\cosh$</td>
<td>$\leq 4$ ulp</td>
</tr>
<tr>
<td>$\sinh$</td>
<td>$\leq 4$ ulp</td>
</tr>
<tr>
<td>$\tanh$</td>
<td>$\leq 5$ ulp</td>
</tr>
<tr>
<td>$\ceiling$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$\text{copysign}$</td>
<td>0 ulp</td>
</tr>
<tr>
<td>$\cos$</td>
<td>$\leq 4$ ulp</td>
</tr>
<tr>
<td>$\cosh$</td>
<td>$\leq 4$ ulp</td>
</tr>
<tr>
<td>$\cos\pi$</td>
<td>$\leq 4$ ulp</td>
</tr>
<tr>
<td>$\text{erfc}$</td>
<td>$\leq 16$ ulp</td>
</tr>
<tr>
<td>$\text{erf}$</td>
<td>$\leq 16$ ulp</td>
</tr>
<tr>
<td>$\exp$</td>
<td>$\leq 3$ ulp</td>
</tr>
<tr>
<td>$\exp2$</td>
<td>$\leq 3$ ulp</td>
</tr>
<tr>
<td>$\exp10$</td>
<td>$\leq 3$ ulp</td>
</tr>
<tr>
<td>$\expm1$</td>
<td>$\leq 3$ ulp</td>
</tr>
<tr>
<td>$\text{fabs}$</td>
<td>0 ulp</td>
</tr>
<tr>
<td>$\text{fdim}$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$\text{floor}$</td>
<td>Correctly rounded</td>
</tr>
</tbody>
</table>

\textsuperscript{68} The ULP values for built-in math functions $\text{lgamma}$ and $\text{lgamma}_r$ is currently undefined.

\textsuperscript{69} 0 ulp is used for math functions that do not require rounding.
<table>
<thead>
<tr>
<th>Function</th>
<th>Property</th>
</tr>
</thead>
<tbody>
<tr>
<td>fma</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>fmax</td>
<td>0 ulp</td>
</tr>
<tr>
<td>fmin</td>
<td>0 ulp</td>
</tr>
<tr>
<td>fmod</td>
<td>0 ulp</td>
</tr>
<tr>
<td>fract</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>frexp</td>
<td>0 ulp</td>
</tr>
<tr>
<td>hypot</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>ilogb</td>
<td>0 ulp</td>
</tr>
<tr>
<td>ldexp</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>log</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>log2</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>log10</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>log1p</td>
<td>&lt;= 2 ulp</td>
</tr>
<tr>
<td>logb</td>
<td>0 ulp</td>
</tr>
<tr>
<td>mad</td>
<td>Implemented either as a correctly rounded fma or as a multiply followed by an add both of which are correctly rounded</td>
</tr>
<tr>
<td>maxmag</td>
<td>0 ulp</td>
</tr>
<tr>
<td>minmag</td>
<td>0 ulp</td>
</tr>
<tr>
<td>modf</td>
<td>0 ulp</td>
</tr>
<tr>
<td>nan</td>
<td>0 ulp</td>
</tr>
<tr>
<td>nextafter</td>
<td>0 ulp</td>
</tr>
<tr>
<td>pow(x, y)</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>pown(x, y)</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>powr(x, y)</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>remainder</td>
<td>0 ulp</td>
</tr>
<tr>
<td>remquo</td>
<td>0 ulp</td>
</tr>
<tr>
<td>rint</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>rootn</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>round</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>rsqrt</td>
<td>&lt;= 2 ulp</td>
</tr>
<tr>
<td>sin</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>sin cos</td>
<td>&lt;= 4 ulp for sine and cosine values</td>
</tr>
<tr>
<td>sinh</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>sin pi</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>sqrt</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>tan</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>tanh</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>tan pi</td>
<td>&lt;= 6 ulp</td>
</tr>
<tr>
<td>tgamma</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>trunc</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>half cos</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half divide</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>Function</td>
<td>Min Accuracy - ULP values</td>
</tr>
<tr>
<td>-------------------</td>
<td>----------------------------</td>
</tr>
<tr>
<td>half_exp</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_exp2</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_exp10</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_log</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_log2</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_log10</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_powr</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_recip</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_rsqrt</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_sin</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_sqrt</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>half_tan</td>
<td>&lt;= 8192 ulp</td>
</tr>
<tr>
<td>native_cos</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_divide</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_exp</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_exp2</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_exp10</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_log</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_log2</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_log10</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_powr</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_recip</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_rsqrt</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_sin</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_sqrt</td>
<td>Implementation-defined</td>
</tr>
<tr>
<td>native_tan</td>
<td>Implementation-defined</td>
</tr>
</tbody>
</table>

**Table 7.1** ULP values for single precision built-in math functions

Table 7.2 describes the minimum accuracy of commonly used single precision floating-point arithmetic operations given as ULP values if the –cl-fast-relaxed-math compiler option is specified when compiling or building an OpenCL program. The minimum accuracy of math functions not defined in *table 7.2* when the –cl-fast-relaxed-math compiler option is specified is as defined in *table 7.1*. The reference value used to compute the ULP value of an arithmetic operation is the infinitely precise result.

<table>
<thead>
<tr>
<th>Function</th>
<th>Min Accuracy - ULP values</th>
</tr>
</thead>
<tbody>
<tr>
<td>1.0 / x</td>
<td>&lt;= 2.5 ulp for x in the domain of $2^{-126}$ to $2^{126}$</td>
</tr>
</tbody>
</table>
| x / y             | <= 2.5 ulp for x in the domain of $2^{-62}$ to $2^{62}$ and y in the domain of $2^{-62}$ to $2^{62}$.

---

70 0 ulp is used for math functions that do not require rounding.
<table>
<thead>
<tr>
<th>Function</th>
<th>Error Bound</th>
</tr>
</thead>
<tbody>
<tr>
<td>( \text{acos}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{acos}(x) ) Implemented as ( \text{acos}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
<tr>
<td>( \text{asin}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{asin}(x) ) Implemented as ( \text{asin}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
<tr>
<td>( \text{atan}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{atan}(x) ) Implemented as ( \text{atan}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
<tr>
<td>( \text{atan}(y, x) )</td>
<td>Implemented as ( \text{atan}(y/x) ) for ( x &gt; 0 ), ( \text{atan}(y/x) + M_{\text{PI}} \text{F} ) for ( x &lt; 0 ) and ( y &gt; 0 ) and ( \text{atan}(y/x) - M_{\text{PI}} \text{F} ) for ( x &lt; 0 ) and ( y &lt; 0 ).</td>
</tr>
<tr>
<td>( \text{atan}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{atan}(x) ) Implemented as ( \text{atan}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
<tr>
<td>( \text{atan}(y, x) )</td>
<td>Implemented as ( \text{atan}(y/x) ) for ( x &gt; 0 ), ( \text{atan}(y/x) + M_{\text{PI}} \text{F} ) for ( x &lt; 0 ) and ( y &gt; 0 ) and ( \text{atan}(y/x) - M_{\text{PI}} \text{F} ) for ( x &lt; 0 ) and ( y &lt; 0 ).</td>
</tr>
<tr>
<td>( \text{atan}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{atan}(x) ) Implemented as ( \text{atan}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
<tr>
<td>( \text{atan}(y, x) )</td>
<td>Implemented as ( \text{atan}(y/x) ) for ( x &gt; 0 ), ( \text{atan}(y/x) + M_{\text{PI}} \text{F} ) for ( x &lt; 0 ) and ( y &gt; 0 ) and ( \text{atan}(y/x) - M_{\text{PI}} \text{F} ) for ( x &lt; 0 ) and ( y &lt; 0 ).</td>
</tr>
<tr>
<td>( \text{acos}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{acos}(x) ) Implemented as ( \text{acos}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
<tr>
<td>( \text{asin}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{asin}(x) ) Implemented as ( \text{asin}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
<tr>
<td>( \text{atan}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{atan}(x) ) Implemented as ( \text{atan}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
<tr>
<td>( \text{atan}(x) )</td>
<td>( \leq 4096 \text{ ulp} )</td>
</tr>
<tr>
<td>( \text{atan}(x) ) Implemented as ( \text{atan}(x) \cdot M_{\text{PI}} \text{F} ). For non-derived implementations, the error is ( \leq 8192 \text{ ulp} ).</td>
<td></td>
</tr>
</tbody>
</table>

For \( x \in [-\pi, \pi] \), the maximum absolute error is \( \leq 2^{-11} \) and larger otherwise.

\[ \text{cosh}(x) \] Implemented as \( 0.5 \times \exp(x) + \exp(-x) \). For non-derived implementations, the error is \( \leq 8192 \text{ ulp} \).

\[ \cos(x) \] For \( x \) in the domain \([0.5, 2]\), the maximum absolute error is \( \leq 3 \text{ ulp} \) for the full profile and \( \leq 4 \text{ ulp} \) for the embedded profile.

\[ \ln(x) \] For \( x \) in the domain \([0.5, 2]\), the maximum absolute error is \( \leq 3 \text{ ulp} \) for the full profile and \( \leq 4 \text{ ulp} \) for the embedded profile.

\[ \log(x) \] For \( x \) in the domain \([0.5, 2]\), the maximum absolute error is \( \leq 2^{-21} \); otherwise the maximum error is \( \leq 3 \text{ ulp} \) for the full profile and \( \leq 4 \text{ ulp} \) for the embedded profile.

\[ \log(x) \] For \( x \) in the domain \([0.5, 2]\), the maximum absolute error is \( \leq 2^{-21} \); otherwise the maximum error is \( \leq 3 \text{ ulp} \) for the full profile and \( \leq 4 \text{ ulp} \) for the embedded profile.

\[ \log(x) \] For \( x \) in the domain \([0.5, 2]\), the maximum absolute error is \( \leq 2^{-21} \); otherwise the maximum error is \( \leq 3 \text{ ulp} \) for the full profile and \( \leq 4 \text{ ulp} \) for the embedded profile.

\[ \log1p(x) \] Derived implementations implement this as \( \log(x + 1) \). For non-derived implementations, the error is
<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pow(x, y)</td>
<td>Undefined for x = 0 and y = 0 or for x &lt; 0 and non-integer y. For x &gt;= 0 or x &lt; 0 and even y, derived implementations implement this as ( \exp2(y \cdot \log2(x)) ). For x &lt; 0 and odd y, derived implementations implement this as (-\exp2(y \cdot \log2(\text{fabs}(x)))). For non-derived implementations, the error is &lt;= 8192 ulp.</td>
</tr>
<tr>
<td>pown(x, y)</td>
<td>Defined only for integer values of y. Undefined for x = 0 and y = 0. For x &gt;= 0 or x &lt; 0 and even y, derived implementations implement this as ( \exp2(y \cdot \log2(x)) ). For x &lt; 0 and odd y, derived implementations implement this as (-\exp2(y \cdot \log2(\text{fabs}(x)))). For non-derived implementations, the error is &lt;= 8192 ulp.</td>
</tr>
<tr>
<td>powr(x, y)</td>
<td>Defined only for x &gt;= 0. Undefined for x = 0 and y = 0. Derived implementations implement this as ( \exp2(y \cdot \log2(x)) ). For non-derived implementations, the error is &lt;= 8192 ulp.</td>
</tr>
<tr>
<td>rootn(x, y)</td>
<td>Defined for x &gt; 0 and for x &lt; 0 when y is odd. Undefined for x = 0 and y = 0. Derived implementations implement this as ( \exp2(\log2(x) / y) ) for x &gt; 0. Derived implementations implement this as (-\exp2(\log2(-x) / y) ) for x &lt; 0. For non-derived implementations, the error is &lt;= 8192 ulp.</td>
</tr>
<tr>
<td>sin(x)</td>
<td>For x in the domain ([-\pi, \pi]), the maximum absolute error is &lt;= (2^{-11}) and larger otherwise.</td>
</tr>
<tr>
<td>sincos(x)</td>
<td>ulp values as defined for (\sin(x)) and (\cos(x))</td>
</tr>
<tr>
<td>sinh(x)</td>
<td>Implemented as (0.5 \cdot \exp(x) - \exp(-x)). For non-derived implementations, the error is &lt;= 8192 ulp.</td>
</tr>
<tr>
<td>sinpi(x)</td>
<td>For x in the domain ([-1, 1]), the maximum absolute error is &lt;= (2^{-11}) and larger otherwise.</td>
</tr>
<tr>
<td>tan(x)</td>
<td>Derived implementations implement this as (\sin(x) \cdot (1.0f / \cos(x))). For non-derived implementations, the error is &lt;= 8192 ulp.</td>
</tr>
<tr>
<td>tanpi(x)</td>
<td>Derived implementations implement this as (\tan(x \cdot \text{M_PI_F})). For non-derived implementations, the error is &lt;= 8192 ulp for x in the domain ([-1, 1]).</td>
</tr>
<tr>
<td>x * y + z</td>
<td>Implemented either as a correctly rounded \text{fma} or as a multiply and an add both of which are correctly rounded.</td>
</tr>
</tbody>
</table>

**Table 7.2** ULP values for single precision built-in math functions with fast relaxed math

*Table 7.3* describes the minimum accuracy of double precision floating-point arithmetic.
operations given as ULP values. The reference value used to compute the ULP value of an arithmetic operation is the infinitely precise result.

<table>
<thead>
<tr>
<th>Function</th>
<th>Min Accuracy - ULP values</th>
</tr>
</thead>
<tbody>
<tr>
<td>$x + y$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$x - y$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$x \times y$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$1.0 / x$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$x / y$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$\cos$</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>$\cos\pi$</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>$\sin$</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>$\sin\pi$</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>$\tan$</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>$\tan2$</td>
<td>&lt;= 6 ulp</td>
</tr>
<tr>
<td>$\tan\pi$</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>$\tan2\pi$</td>
<td>&lt;= 6 ulp</td>
</tr>
<tr>
<td>$\cosh$</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>$\sinh$</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>$\tanh$</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>$\text{cbrt}$</td>
<td>&lt;= 2 ulp</td>
</tr>
<tr>
<td>$\text{ceil}$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$\text{copysign}$</td>
<td>0 ulp</td>
</tr>
<tr>
<td>$\text{cos}$</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>$\text{cosh}$</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>$\text{cos}\pi$</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>$\text{erfc}$</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>$\text{erf}$</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>$\exp$</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>$\exp2$</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>$\exp10$</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>$\expm1$</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>$\text{fabs}$</td>
<td>0 ulp</td>
</tr>
<tr>
<td>$\text{fdim}$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$\text{floor}$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$\text{fma}$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$\text{fmax}$</td>
<td>0 ulp</td>
</tr>
<tr>
<td>$\text{fmin}$</td>
<td>0 ulp</td>
</tr>
<tr>
<td>$\text{fmod}$</td>
<td>0 ulp</td>
</tr>
<tr>
<td>$\text{fract}$</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>$\text{fexp}$</td>
<td>0 ulp</td>
</tr>
<tr>
<td>$\text{hypot}$</td>
<td>&lt;= 4 ulp</td>
</tr>
</tbody>
</table>

71 0 ulp is used for math functions that do not require rounding.
### Table 7.3
ULP values for double precision built-in math functions

<table>
<thead>
<tr>
<th>Function</th>
<th>ULP Bound</th>
</tr>
</thead>
<tbody>
<tr>
<td>ilogb</td>
<td>0 ulp</td>
</tr>
<tr>
<td>ldexp</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>log</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>log2</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>log10</td>
<td>&lt;= 3 ulp</td>
</tr>
<tr>
<td>log1p</td>
<td>&lt;= 2 ulp</td>
</tr>
<tr>
<td>logb</td>
<td>0 ulp</td>
</tr>
<tr>
<td>mad</td>
<td>Any value allowed (infinite ulp)</td>
</tr>
<tr>
<td>maxmag</td>
<td>0 ulp</td>
</tr>
<tr>
<td>minmag</td>
<td>0 ulp</td>
</tr>
<tr>
<td>modf</td>
<td>0 ulp</td>
</tr>
<tr>
<td>nan</td>
<td>0 ulp</td>
</tr>
<tr>
<td>nextafter</td>
<td>0 ulp</td>
</tr>
<tr>
<td>pow(x, y)</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>pown(x, y)</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>powr(x, y)</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>remainder</td>
<td>0 ulp</td>
</tr>
<tr>
<td>remqu</td>
<td>0 ulp</td>
</tr>
<tr>
<td>rint</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>rootn</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>round</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>rsqrt</td>
<td>&lt;= 2 ulp</td>
</tr>
<tr>
<td>sin</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>sincos</td>
<td>&lt;= 4 ulp for sine and cosine values</td>
</tr>
<tr>
<td>sinh</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>sinpi</td>
<td>&lt;= 4 ulp</td>
</tr>
<tr>
<td>sqrt</td>
<td>Correctly rounded</td>
</tr>
<tr>
<td>tan</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>tanh</td>
<td>&lt;= 5 ulp</td>
</tr>
<tr>
<td>tanpi</td>
<td>&lt;= 6 ulp</td>
</tr>
<tr>
<td>tgamma</td>
<td>&lt;= 16 ulp</td>
</tr>
<tr>
<td>trunc</td>
<td>Correctly rounded</td>
</tr>
</tbody>
</table>

### 7.5 Edge Case Behavior

The edge case behavior of the math functions (section 6.13.2) shall conform to sections F.9 and G.6 of ISO/IEC 9899:TC 2 (commonly known as C99, TC2), except where noted below in section 7.5.1.
7.5.1 Additional Requirements Beyond C99 TC2

Functions that return a NaN with more than one NaN operand shall return one of the NaN operands. Functions that return a NaN operand may silence the NaN if it is a signaling NaN. A non-signaling NaN shall be converted to a non-signaling NaN. A signaling NaN shall be converted to a NaN, and should be converted to a non-signaling NaN. How the rest of the NaN payload bits or the sign of NaN is converted is undefined.

half_<funcname> functions behave identically to the function of the same name without the half_prefix. They must conform to the same edge case requirements (see sections F.9 and G.6 of C99, TC2). For other cases, except where otherwise noted, these single precision functions are permitted to have up to 8192 ulps of error (as measured in the single precision result), although better accuracy is encouraged.

The usual allowances for rounding error (section 7.4) or flushing behavior (section 7.5.3) shall not apply for those values for which section F.9 of C99, TC2, or sections 7.5.1 and 7.5.3 below (and similar sections for other floating-point precisions) prescribe a result (e.g. ceil ( -1 < x < 0 ) returns -0). Those values shall produce exactly the prescribed answers, and no other. Where the ± symbol is used, the sign shall be preserved. For example, sin(±0) = ±0 shall be interpreted to mean sin(+0) is +0 and sin(-0) is -0.

\[
\begin{align*}
\text{acospi} (1) &= +0. \\
\text{acospi} (x) &\text{ returns a NaN for } |x| > 1. \\
\text{asinpi} (±0) &= ±0. \\
\text{asinpi} (x) &\text{ returns a NaN for } |x| > 1. \\
\text{atanpi} (±0) &= ±0. \\
\text{atanpi} (±∞) &= ±0.5.
\end{align*}
\]

\[
\begin{align*}
\text{atan2pi} (±0, -0) &= ±1. \\
\text{atan2pi} (±0, +0) &= ±0. \\
\text{atan2pi} (±0, x) &\text{ returns } ±1 \text{ for } x < 0. \\
\text{atan2pi} (±0, x) &\text{ returns } ±0 \text{ for } x > 0. \\
\text{atan2pi} (y, ±0) &\text{ returns } -0.5 \text{ for } y < 0. \\
\text{atan2pi} (y, ±0) &\text{ returns } 0.5 \text{ for } y > 0. \\
\text{atan2pi} (±y, -∞) &\text{ returns } ±1 \text{ for finite } y > 0. \\
\text{atan2pi} (±y, +∞) &\text{ returns } ±0 \text{ for finite } y > 0. \\
\text{atan2pi} (±∞, x) &\text{ returns } ±0.5 \text{ for finite } x. \\
\text{atan2pi} (±∞, -∞) &\text{ returns } ±0.75. \\
\text{atan2pi} (±∞, +∞) &\text{ returns } ±0.25.
\end{align*}
\]

\[
\begin{align*}
\text{ceil} (-1 < x < 0) &\text{ returns } -0. \\
\cospi (±0) &\text{ returns } 1 \\
\cospi (n + 0.5) &\text{ is } +0 \text{ for any integer } n \text{ where } n + 0.5 \text{ is representable.}
\end{align*}
\]
cospi ( ±∞ ) returns a NaN.

exp10 ( ±0 ) returns 1.
exp10 ( -∞ ) returns +0.
exp10 ( +∞ ) returns +∞.

distance (x, y) calculates the distance from x to y without overflow or extraordinary precision loss due to underflow.

fdim ( any, NaN ) returns NaN.
fdim ( NaN, any ) returns NaN.

fmod ( ±0, NaN )  returns NaN.

frexp ( ±∞, exp ) returns ±∞ and stores 0 in exp.
frexp ( NaN, exp ) returns the NaN and stores 0 in exp.

fract ( x, iptr ) shall not return a value greater than or equal to 1.0, and shall not return a value less than 0.
fract ( +0, iptr ) returns +0 and +0 in iptr.
fract ( -0, iptr ) returns -0 and -0 in iptr.
fract ( +inf, iptr ) returns +0 and +inf in iptr.
fract ( -inf, iptr ) returns -0 and -inf in iptr.
fract ( NaN, iptr ) returns the NaN and NaN in iptr.

length calculates the length of a vector without overflow or extraordinary precision loss due to underflow.

lgamma_r (x, signp) returns 0 in signp if x is zero or a negative integer.

nextafter ( -0, y > 0 ) returns smallest positive denormal value.
nextafter ( +0, y < 0 ) returns smallest negative denormal value.

normalize shall reduce the vector to unit length, pointing in the same direction without overflow or extraordinary precision loss due to underflow.
normalize ( v ) returns v if all elements of v are zero.
normalize ( v ) returns a vector full of NaNs if any element is a NaN.

normalize ( v ) for which any element in v is infinite shall proceed as if the elements in v were replaced as follows:

for( i = 0; i < sizeof(v) / sizeof(v[0]); i++)
    v[i] = isinf(v[i]) ? copysign(1.0, v[i]) : 0.0 * v[i];

pow ( ±0, -∞ ) returns +∞
\textbf{pown} (x, 0) is 1 for any x, even zero, \text{NaN} or infinity.
\textbf{pown} (±0, n) is ±∞ for odd n < 0.
\textbf{pown} (±0, n) is +∞ for even n < 0.
\textbf{pown} (±0, n) is +0 for even n > 0.
\textbf{pown} (±0, n) is ±0 for odd n > 0.

\textbf{powr} (x, ±0) is 1 for finite x > 0.
\textbf{powr} (±0, y) is +∞ for finite y < 0.
\textbf{powr} (±0, -∞) is +∞.
\textbf{powr} (±0, y) is +0 for y > 0.
\textbf{powr} (+1, y) is 1 for finite y.
\textbf{powr} (x, y) returns \text{NaN} for x < 0.
\textbf{powr} (±0, ±0) returns \text{NaN}.
\textbf{powr} (+∞, ±0) returns \text{NaN}.
\textbf{powr} (+1, ±∞) returns \text{NaN}.
\textbf{powr} (x, \text{NaN}) returns the \text{NaN} for x >= 0.
\textbf{powr} (\text{NaN}, y) returns the \text{NaN}.

\textbf{rint} (-0.5 <= x < 0) returns -0.

\textbf{remquo} (x, y, &quo) returns a \text{NaN} and 0 in quo if x is ±∞, or if y is 0 and the other argument is non-NaN or if either argument is a NaN.

\textbf{rootn} (±0, n) is ±∞ for odd n < 0.
\textbf{rootn} (±0, n) is +∞ for even n < 0.
\textbf{rootn} (±0, n) is +0 for even n > 0.
\textbf{rootn} (±0, n) is ±0 for odd n > 0.
\textbf{rootn} (x, n) returns a \text{NaN} for x < 0 and n is even.
\textbf{rootn} (x, 0) returns a \text{NaN}.

\textbf{round} (-0.5 < x < 0) returns -0.

\textbf{sinpi} (±0) returns ±0.
\textbf{sinpi} (+n) returns +0 for positive integers n.
\textbf{sinpi} (-n) returns -0 for negative integers n.
\textbf{sinpi} (±∞) returns a \text{NaN}.

\textbf{tanpi} (±0) returns ±0.
\textbf{tanpi} (±∞) returns a \text{NaN}.
\textbf{tanpi} (n) is \text{copysign}(0.0, n) for even integers n.
\textbf{tanpi} (n) is \text{copysign}(0.0, -n) for odd integers n.
\textbf{tanpi} (n + 0.5) for even integer n is +∞ where n + 0.5 is representable.
\textbf{tanpi} (n + 0.5) for odd integer n is -∞ where n + 0.5 is representable.

\textbf{trunc} (-1 < x < 0) returns -0.
7.5.2 Changes to C99 TC2 Behavior

`modf` behaves as though implemented by:

```c
gentype modf ( gentype value, gentype *iptr )
{
    *
    i
    ptr
    =
    trunc(
    value );
    return copysign( isinf( value ) ? 0.0 : value – *iptr, value );
}
```

`rint` always rounds according to round to nearest even rounding mode even if the caller is in some other rounding mode.

7.5.3 Edge Case Behavior in Flush To Zero Mode

If denormals are flushed to zero, then a function may return one of four results:

1. Any conforming result for non-flush-to-zero mode

2. If the result given by 1. is a sub-normal before rounding, it may be flushed to zero

3. Any non-flushed conforming result for the function if one or more of its sub-normal operands are flushed to zero.

4. If the result of 3. is a sub-normal before rounding, the result may be flushed to zero.

In each of the above cases, if an operand or result is flushed to zero, the sign of the zero is undefined.

If subnormals are flushed to zero, a device may choose to conform to the following edge cases for `nextafter` instead of those listed in section 7.5.1:

- `nextafter` ( +smallest normal, $y < +$smallest normal ) = +0.
- `nextafter` ( -smallest normal, $y > -$smallest normal ) = -0.
- `nextafter` ( -0, $y > 0$ ) returns smallest positive normal value.
- `nextafter` ( +0, $y < 0$ ) returns smallest negative normal value.

For clarity, subnormals or denormals are defined to be the set of representable numbers in the range $0 < \times < \text{TYPE}\_\text{MIN}$ and $-\text{TYPE}\_\text{MIN} < \times < -0$. They do not include ±0. A non-zero number is said to be sub-normal before rounding if after normalization, its radix-2 exponent is less than $(\text{TYPE}\_\text{MIN}\_\text{EXP} - 1)$.\textsuperscript{72}

\textsuperscript{72} Here \text{TYPE}\_\text{MIN} and \text{TYPE}\_\text{MIN}\_\text{EXP} should be substituted by constants appropriate to the floating-point type under consideration, such as \text{FLT}\_\text{MIN} and \text{FLT}\_\text{MIN}\_\text{EXP} for float.
8. Image Addressing and Filtering

Let $w_t$, $h_t$ and $d_t$ be the width, height (or image array size for a 1D image array) and depth (or image array size for a 2D image array) of the image in pixels. Let $\text{coord.x}$ also referred to as $(s,t)$ or $\text{coord.x.y.z}$ also referred to as $(s,t,r)$ be the coordinates specified to \texttt{read_image{f|i|ui}}. The sampler specified in \texttt{read_image{f|i|ui}} is used to determine how to sample the image and return an appropriate color.

8.1 Image Coordinates

This affects the interpretation of image coordinates. If image coordinates specified to \texttt{read_image{f|i|ui}} are normalized (as specified in the sampler), the $s$, $t$, and $r$ coordinate values are multiplied by $w_t$, $h_t$, and $d_t$ respectively to generate the unnormalized coordinate values. For image arrays, the image array coordinate (i.e. $t$ if it is a 1D image array or $r$ if it is a 2D image array) specified to \texttt{read_image{f|i|ui}} must always be the un-normalized image coordinate value.

Let $(u, v, w)$ represent the unnormalized image coordinate values.

8.2 Addressing and Filter Modes

We first describe how the addressing and filter modes are applied to generate the appropriate sample locations to read from the image if the addressing mode is not \text{CLK_ADDRESS_REPEAT} or \text{CLK_ADDRESS_MIRRORED_REPEAT}.

After generating the image coordinate $(u, v, w)$ we apply the appropriate addressing and filter mode to generate the appropriate sample locations to read from the image.

If values in $(u, v, w)$ are INF or NaN, the behavior of \texttt{read_image{f|i|ui}} is undefined.

Filter Mode = \text{CLK_FILTER_NEAREST}

When filter mode is \text{CLK_FILTER_NEAREST}, the image element in the image that is nearest (in Manhattan distance) to that specified by $(u, v, w)$ is obtained. This means the image element at location $(i, j, k)$ becomes the image element value, where

\[
\begin{align*}
i &= \text{address_mode}((\text{int})\text{floor}(u)) \\
j &= \text{address_mode}((\text{int})\text{floor}(v)) \\
k &= \text{address_mode}((\text{int})\text{floor}(w))
\end{align*}
\]
For a 3D image, the image element at location \((i, j, k)\) becomes the color value. For a 2D image, the image element at location \((i, j)\) becomes the color value.

Table 8.1 describes the \texttt{address\_mode} function.

<table>
<thead>
<tr>
<th>Addressing Mode</th>
<th>Result of \texttt{address_mode(coord)}</th>
</tr>
</thead>
<tbody>
<tr>
<td>exttt{CLK_ADDRESS_CLAMP_TO_EDGE}</td>
<td>\texttt{clamp(coord, 0, size - 1)}</td>
</tr>
<tr>
<td>exttt{CLK_ADDRESS_CLAMP}</td>
<td>\texttt{clamp(coord, -1, size)}</td>
</tr>
<tr>
<td>exttt{CLK_ADDRESS_NONE}</td>
<td>coord</td>
</tr>
</tbody>
</table>

\textbf{Table 8.1} \textemdash \textit{Addressing modes to generate texel location.}

The \texttt{size} term in table 8.1 is \(w_t\) for \(u\), \(h_t\) for \(v\) and \(d_t\) for \(w\).

The \texttt{clamp} function used in \textit{table 8.1} is defined as:

\[
\texttt{clamp}(a, b, c) = \text{return } (a < b) ? b : ((a > c) ? c : a)
\]

If the selected texel location \((i, j, k)\) refers to a location outside the image, the border color is used as the color value for this texel.

\textbf{Filter Mode} = \texttt{CLK\_FILTER\_LINEAR}

When filter mode is \texttt{CLK\_FILTER\_LINEAR}, a \(2 \times 2\) square of image elements for a 2D image or a \(2 \times 2 \times 2\) cube of image elements for a 3D image is selected. This \(2 \times 2\) square or \(2 \times 2 \times 2\) cube is obtained as follows.

Let

\[
\begin{align*}
  i_0 &= \text{address\_mode}((\text{int})\text{floor}(u - 0.5)) \\
  j_0 &= \text{address\_mode}((\text{int})\text{floor}(v - 0.5)) \\
  k_0 &= \text{address\_mode}((\text{int})\text{floor}(w - 0.5)) \\
  i_1 &= \text{address\_mode}((\text{int})\text{floor}(u - 0.5) + 1) \\
  j_1 &= \text{address\_mode}((\text{int})\text{floor}(v - 0.5) + 1) \\
  k_1 &= \text{address\_mode}((\text{int})\text{floor}(w - 0.5) + 1) \\
  a &= \text{frac}(u - 0.5) \\
  b &= \text{frac}(v - 0.5) \\
  c &= \text{frac}(w - 0.5)
\end{align*}
\]

where \texttt{frac(x)} denotes the fractional part of \(x\) and is computed as \(x - \text{floor}(x)\).

For a 3D image, the image element value is found as

\[
T = (1 - a) \times (1 - b) \times (1 - c) \times T_{i_0j_0k_0} \\
+ a \times (1 - b) \times (1 - c) \times T_{i_1j_0k_0}
\]
\[ + (1 - a) \times b \times (1 - c) \times T_{i0j1k0} \]
\[ + a \times b \times (1 - c) \times T_{i1j1k0} \]
\[ + (1 - a) \times (1 - b) \times c \times T_{i0j0k1} \]
\[ + a \times (1 - b) \times c \times T_{i1j0k1} \]
\[ + (1 - a) \times b \times c \times T_{i0j1k1} \]
\[ + a \times b \times c \times T_{i1j1k1} \]

where \( T_{ijk} \) is the image element at location \((i, j, k)\) in the 3D image.

For a 2D image, the image element value is found as

\[ T = (1 - a) \times (1 - b) \times T_{i0j0} \]
\[ + a \times (1 - b) \times T_{i1j0} \]
\[ + (1 - a) \times b \times T_{i0j1} \]
\[ + a \times b \times T_{i1j1} \]

where \( T_{ij} \) is the image element at location \((i, j)\) in the 2D image.

If any of the selected \( T_{ijk} \) or \( T_{ij} \) in the above equations refers to a location outside the image, the border color is used as the color value for \( T_{ijk} \) or \( T_{ij} \).

We now discuss how the addressing and filter modes are applied to generate the appropriate sample locations to read from the image if the addressing mode is CLK_ADDRESS_REPEAT.

If values in \((s, t, r)\) are INF or NaN, the behavior of the built-in image read functions is undefined.

**Filter Mode = CLK_FILTER_NEAREST**

When filter mode is CLK_FILTER_NEAREST, the image element at location \((i, j, k)\) becomes the image element value, with \( i, j \) and \( k \) computed as

\[
  u = (s - \text{floor}(s)) \times w_t \\
  i = (\text{int})\text{floor}(u) \\
  \text{if} \ (i > w_t - 1) \\
      i = i - w_t \\
\]
\[
  v = (t - \text{floor}(t)) \times h_t \\
  j = (\text{int})\text{floor}(v) \\
  \text{if} \ (j > h_t - 1) \\
      j = j - h_t \\
\]
\[
  w = (r - \text{floor}(r)) \times d_t \\
\]
\[ k = (\text{int}) \lfloor \frac{w}{d_t} \rfloor \]
\[ \text{if } (k > d_t - 1) \]
\[ k = k - d_t \]

For a 3D image, the image element at location \((i, j, k)\) becomes the color value. For a 2D image, the image element at location \((i, j)\) becomes the color value.

**Filter Mode = CLK_FILTER_LINEAR**

When filter mode is CLK_FILTER_LINEAR, a \(2 \times 2\) square of image elements for a 2D image or a \(2 \times 2 \times 2\) cube of image elements for a 3D image is selected. This \(2 \times 2\) square or \(2 \times 2 \times 2\) cube is obtained as follows.

Let

\[ u = (s - \lfloor s \rfloor) \times w_t \]
\[ i_0 = (\text{int}) \lfloor u - 0.5 \rfloor \]
\[ i_1 = i_0 + 1 \]
\[ \text{if } (i_0 < 0) \]
\[ i_0 = w_t + i_0 \]
\[ \text{if } (i_1 > w_t - 1) \]
\[ i_1 = i_1 - w_t \]

\[ v = (t - \lfloor t \rfloor) \times h_t \]
\[ j_0 = (\text{int}) \lfloor v - 0.5 \rfloor \]
\[ j_1 = j_0 + 1 \]
\[ \text{if } (j_0 < 0) \]
\[ j_0 = h_t + j_0 \]
\[ \text{if } (j_1 > h_t - 1) \]
\[ j_1 = j_1 - h_t \]

\[ w = (r - \lfloor r \rfloor) \times d_t \]
\[ k_0 = (\text{int}) \lfloor w - 0.5 \rfloor \]
\[ k_1 = k_0 + 1 \]
\[ \text{if } (k_0 < 0) \]
\[ k_0 = d_t + k_0 \]
\[ \text{if } (k_1 > d_t - 1) \]
\[ k_1 = k_1 - d_t \]

\[ a = \text{frac}(u - 0.5) \]
\[ b = \text{frac}(v - 0.5) \]
\[ c = \text{frac}(w - 0.5) \]

where \(\text{frac}(x)\) denotes the fractional part of \(x\) and is computed as \(x - \lfloor x \rfloor\).

For a 3D image, the image element value is found as
\[ T = (1 - a) \times (1 - b) \times (1 - c) \times T_{i0j0k0} \]
\[ + a \times (1 - b) \times (1 - c) \times T_{i1j0k0} \]
\[ + (1 - a) \times b \times (1 - c) \times T_{i0j1k0} \]
\[ + a \times b \times (1 - c) \times T_{i1j1k0} \]
\[ + (1 - a) \times (1 - b) \times c \times T_{i1j0k1} \]
\[ + a \times (1 - b) \times c \times T_{i1j1k1} \]
\[ + (1 - a) \times b \times c \times T_{i0j1k1} \]
\[ + a \times b \times c \times T_{i1j1k1} \]

where \( T_{ijk} \) is the image element at location \((i, j, k)\) in the 3D image.

For a 2D image, the image element value is found as

\[ T = (1 - a) \times (1 - b) \times T_{i0j0} \]
\[ + a \times (1 - b) \times T_{i1j0} \]
\[ + (1 - a) \times b \times T_{i0j1} \]
\[ + a \times b \times T_{i1j1} \]

where \( T_{ij} \) is the image element at location \((i, j)\) in the 2D image.

We now discuss how the addressing and filter modes are applied to generate the appropriate sample locations to read from the image if the addressing mode is CLK_ADDRESS_MIRRORED_REPEAT. The CLK_ADDRESS_MIRRORED_REPEAT addressing mode causes the image to be read as if it is tiled at every integer seam with the interpretation of the image data flipped at each integer crossing. For example, the \((s, t, r)\) coordinates between 2 and 3 are addressed into the image as coordinates from 1 down to 0. If values in \((s, t, r)\) are INF or NaN, the behavior of the built-in image read functions is undefined.

**Filter Mode = CLK_FILTER_NEAREST**

When filter mode is CLK_FILTER_NEAREST, the image element at location \((i, j, k)\) becomes the image element value, with \(i, j\) and \(k\) computed as

\[ s' = 2.0f \times \text{rint}(0.5f \times s) \]
\[ s' = \text{fabs}(s - s') \]
\[ u = s' \times w_t \]
\[ i = \text{(int)}\text{floor}(u) \]
\[ i = \min(i, w_t - 1) \]
\[ t' = 2.0f \times \text{rint}(0.5f \times t) \]
\[ t' = \text{fabs}(t - t') \]
\[ v = t' \times h_t \]
j = (int)floor(v)
j = min(j, h_t - 1)

r' = 2.0f * rint(0.5f * r)
r' = fabs(r - r')
w = r' * d_t
k = (int)floor(w)
k = min(k, d_t - 1)

For a 3D image, the image element at location (i, j, k) becomes the color value. For a 2D image, the image element at location (i, j) becomes the color value.

**Filter Mode = CLK_FILTER_LINEAR**

When filter mode is CLK_FILTER_LINEAR, a 2 x 2 square of image elements for a 2D image or a 2 x 2 x 2 cube of image elements for a 3D image is selected. This 2 x 2 square or 2 x 2 x 2 cube is obtained as follows.

Let

s' = 2.0f * rint(0.5f * s)
s' = fabs(s - s')
u = s' * w_t
i0 = (int)floor(u - 0.5f)
i1 = i0 + 1
i0 = max(i0, 0)
i1 = min(i1, w_t - 1)

t' = 2.0f * rint(0.5f * t)
t' = fabs(t - t')
v = t' * h_t
j0 = (int)floor(v - 0.5f)
j1 = j0 + 1
j0 = max(j0, 0)
j1 = min(j1, h_t - 1)

r' = 2.0f * rint(0.5f * r)
r' = fabs(r - r')
w = r' * d_t
k0 = (int)floor(w - 0.5f)
k1 = k0 + 1
k0 = max(k0, 0)
k1 = min(k1, d_t - 1)

a = frac(u - 0.5)
b = frac(v - 0.5)
\[ c = \frac{w - 0.5}{w - 0.5} \]

where \( \text{frac}(x) \) denotes the fractional part of \( x \) and is computed as \( x - \text{floor}(x) \).

For a 3D image, the image element value is found as

\[
 T = (1 - a) \cdot (1 - b) \cdot (1 - c) \cdot T_{i0j0k0} + a \cdot (1 - b) \cdot (1 - c) \cdot T_{i1j0k0} + (1 - a) \cdot b \cdot (1 - c) \cdot T_{i0j1k0} + a \cdot b \cdot (1 - c) \cdot T_{i1j1k0} + (1 - a) \cdot c \cdot T_{i0j0k1} + a \cdot c \cdot c \cdot T_{i1j0k1} + (1 - a) \cdot (1 - b) \cdot c \cdot T_{i0j1k1} + a \cdot b \cdot c \cdot T_{i1j1k1}
\]

where \( T_{ijk} \) is the image element at location \((i, j, k)\) in the 3D image.

For a 2D image, the image element value is found as

\[
 T = (1 - a) \cdot (1 - b) \cdot T_{i0j0} + a \cdot (1 - b) \cdot T_{i1j0} + (1 - a) \cdot b \cdot T_{i0j1} + a \cdot b \cdot T_{i1j1}
\]

where \( T_{ij} \) is the image element at location \((i, j)\) in the 2D image.

For a 1D image, the image element value is found as

\[
 T = (1 - a) \cdot T_{i0} + a \cdot T_{i1}
\]

where \( T_{i} \) is the image element at location \((i)\) in the 1D image.

**NOTE**

*If the sampler is specified as using unnormalized coordinates (floating-point or integer coordinates), filter mode set to CLK_FILTER_NEAREST and addressing mode set to one of the following modes - CLK_ADDRESS_NONE, CLK_ADDRESS_CLAMP_TO_EDGE or CLK_ADDRESS_CLAMP, the location of the image element in the image given by \((i, j, k)\) in section 8.2 will be computed without any loss of precision.*

*For all other sampler combinations of normalized or unnormalized coordinates, filter and addressing modes, the relative error or precision of the addressing mode calculations and the*
image filter operation are not defined by this revision of the OpenCL specification. To ensure a minimum precision of image addressing and filter calculations across any OpenCL device, for these sampler combinations, developers should unnormalize the image coordinate in the kernel and implement the linear filter in the kernel with appropriate calls to \texttt{read\_imagef[i]ui} with a sampler that uses unnormalized coordinates, filter mode set to \texttt{CLK\_FILTER\_NEAREST}, addressing mode set to \texttt{CLK\_ADDRESS\_NONE}, \texttt{CLK\_ADDRESS\_CLAMP\_TO\_EDGE} or \texttt{CLK\_ADDRESS\_CLAMP} and finally performing the interpolation of color values read from the image to generate the filtered color value.

### 8.3 Conversion Rules

In this section we discuss conversion rules that are applied when reading and writing images in a kernel.

#### 8.3.1 Conversion rules for normalized integer channel data types

In this section we discuss converting normalized integer channel data types to floating-point values and vice-versa.

##### 8.3.1.1 Converting normalized integer channel data types to floating-point values

For images created with image channel data type of \texttt{CL\_UNORM\_INT8} and \texttt{CL\_UNORM\_INT16}, \texttt{read\_imagef} will convert the channel values from an 8-bit or 16-bit unsigned integer to normalized floating-point values in the range \([0.0 \ f \ ... \ 1.0]\).

For images created with image channel data type of \texttt{CL\_SNORM\_INT8} and \texttt{CL\_SNORM\_INT16}, \texttt{read\_imagef} will convert the channel values from an 8-bit or 16-bit signed integer to normalized floating-point values in the range \([-1.0 \ ... \ 1.0]\).

These conversions are performed as follows:

\[
\text{CL\_UNORM\_INT8} \ (8\text{-bit unsigned integer}) \rightarrow \text{float} \\
\text{normalized float value} = (\text{float})c / 255.0f
\]

\[
\text{CL\_UNORM\_INT101010} \ (10\text{-bit unsigned integer}) \rightarrow \text{float} \\
\text{normalized float value} = (\text{float})c / 1023.0f
\]

\[
\text{CL\_UNORM\_INT16} \ (16\text{-bit unsigned integer}) \rightarrow \text{float}
\]
normalized float value = (float)c / 65535.0f

CL_SNORM_INT8 (8-bit signed integer) → float

normalized float value = \max(-1.0f, \frac{c}{127.0f})

CL_SNORM_INT16 (16-bit signed integer) → float

normalized float value = \max(-1.0f, \frac{c}{32767.0f})

The precision of the above conversions is <= 1.5 ulp except for the following cases.

For CL_UNORM_INT8

0 must convert to 0.0f and
255 must convert to 1.0f

For CL_UNORM_INT_101010

0 must convert to 0.0f and
1023 must convert to 1.0f

For CL_UNORM_INT16

0 must convert to 0.0f and
65535 must convert to 1.0f

For CL_SNORM_INT8

-128 and -127 must convert to -1.0f,
0 must convert to 0.0f and
127 must convert to 1.0f

For CL_SNORM_INT16

-32768 and -32767 must convert to -1.0f,
0 must convert to 0.0f and
32767 must convert to 1.0f

8.3.1.2 Converting floating-point values to normalized integer channel data types

For images created with image channel data type of CL_UNORM_INT8 and CL_UNORM_INT16, write_imagef will convert the floating-point color value to an 8-bit or 16-bit unsigned integer.
For images created with image channel data type of CL_SNORM_INT8 and CL_SNORM_INT16, write_imagef will convert the floating-point color value to an 8-bit or 16-bit signed integer.

The preferred method for how conversions from floating-point values to normalized integer values are performed is as follows:

- float $\rightarrow$ CL_UNORM_INT8 (8-bit unsigned integer)
  
  \[ \text{convert}_\text{uchar}_\text{sat}_\text{rte}(f \times 255.0f) \]

- float $\rightarrow$ CL_UNORM_INT101010 (10-bit unsigned integer)
  
  \[ \min(\text{convert}_\text{ushort}_\text{sat}_\text{rte}(f \times 1023.0f), 0x3ff) \]

- float $\rightarrow$ CL_UNORM_INT16 (16-bit unsigned integer)
  
  \[ \text{convert}_\text{ushort}_\text{sat}_\text{rte}(f \times 65535.0f) \]

- float $\rightarrow$ CL_SNORM_INT8 (8-bit signed integer)
  
  \[ \text{convert}_\text{char}_\text{sat}_\text{rte}(f \times 127.0f) \]

- float $\rightarrow$ CL_SNORM_INT16 (16-bit signed integer)
  
  \[ \text{convert}_\text{short}_\text{sat}_\text{rte}(f \times 32767.0f) \]

Please refer to section 6.2.3.3 for out-of-range behavior and saturated conversions rules.

OpenCL implementations may choose to approximate the rounding mode used in the conversions described above. If a rounding mode other than round to nearest even (_rte) is used, the absolute error of the implementation dependant rounding mode vs. the result produced by the round to nearest even rounding mode must be $\leq 0.6$.

- float $\rightarrow$ CL_UNORM_INT8 (8-bit unsigned integer)
  
  Let \( f_{\text{preferred}} = \text{convert}_\text{uchar}_\text{sat}_\text{rte}(f \times 255.0f) \)

  Let \( f_{\text{approx}} = \text{convert}_\text{uchar}_\text{sat}_\langle \text{impl\,-\,rounding\,-\,mode} \rangle(f \times 255.0f) \)

  \[ \text{fabs}(f_{\text{preferred}} - f_{\text{approx}})\text{ must be } \leq 0.6 \]

- float $\rightarrow$ CL_UNORM_INT101010 (10-bit unsigned integer)
  
  Let \( f_{\text{preferred}} = \text{convert}_\text{ushort}_\text{sat}_\text{rte}(f \times 1023.0f) \)

  Let \( f_{\text{approx}} = \text{convert}_\text{ushort}_\text{sat}_\langle \text{impl\,-\,rounding\,-\,mode} \rangle(f \times 1023.0f) \)
fabs(f_{preferred} - f_{approx}) must be <= 0.6

float \rightarrow CL\_UNORM\_INT16 (16-bit unsigned integer)

Let \( f_{preferred} = \text{convert\_ushort\_sat\_rte}(f \times 65535.0f) \)
Let \( f_{approx} = \text{convert\_ushort\_sat\_impl\_rounding\_mode}(f \times 65535.0f) \)

fabs(f_{preferred} - f_{approx}) must be <= 0.6

float \rightarrow CL\_SNORM\_INT8 (8-bit signed integer)

Let \( f_{preferred} = \text{convert\_char\_sat\_rte}(f \times 127.0f) \)
Let \( f_{approx} = \text{convert\_char\_sat\_impl\_rounding\_mode}(f \times 127.0f) \)

fabs(f_{preferred} - f_{approx}) must be <= 0.6

float \rightarrow CL\_SNORM\_INT16 (16-bit signed integer)

Let \( f_{preferred} = \text{convert\_short\_sat\_rte}(f \times 32767.0f) \)
Let \( f_{approx} = \text{convert\_short\_sat\_impl\_rounding\_mode}(f \times 32767.0f) \)

fabs(f_{preferred} - f_{approx}) must be <= 0.6

**8.3.2 Conversion rules for half precision floating-point channel data type**

For images created with a channel data type of CL\_HALF\_FLOAT, the conversions from half to float are lossless (as described in section 6.1.1.1). Conversions from float to half round the mantissa using the round to nearest even or round to zero rounding mode. Denormalized numbers for the half data type which may be generated when converting a float to a half may be flushed to zero. A float NaN must be converted to an appropriate NaN in the half type. A float INF must be converted to an appropriate INF in the half type.

**8.3.3 Conversion rules for floating-point channel data type**

The following rules apply for reading and writing images created with channel data type of CL\_FLOAT.
NaNs may be converted to a NaN value(s) supported by the device.

Denorms can be flushed to zero.

All other values must be preserved.

8.3.4 Conversion rules for signed and unsigned 8-bit, 16-bit and 32-bit integer channel data types

Calls to read_imagei with channel data type values of CL_SIGNED_INT8, CL_SIGNED_INT16 and CL_SIGNED_INT32 return the unmodified integer values stored in the image at specified location.

Calls to read_imageui with channel data type values of CL_UNSIGNED_INT8, CL_UNSIGNED_INT16 and CL_UNSIGNED_INT32 return the unmodified integer values stored in the image at specified location.

Calls to write_imagei will perform one of the following conversions:

32 bit signed integer → 8-bit signed integer
    convert_char_sat(i)

32 bit signed integer → 16-bit signed integer
    convert_short_sat(i)

32 bit signed integer → 32-bit signed integer
    no conversion is performed

Calls to write_imageui will perform one of the following conversions:

32 bit unsigned integer → 8-bit unsigned integer
    convert_uchar_sat(i)

32 bit unsigned integer → 16-bit unsigned integer
    convert_ushort_sat(i)

32 bit unsigned integer → 32-bit unsigned integer
    no conversion is performed
The conversions described in this section must be correctly saturated.

### 8.3.5 Conversion rules for sRGBA and sBGRA images

Standard RGB data, which roughly displays colors in a linear ramp of luminosity levels such that an average observer, under average viewing conditions, can view them as perceptually equal steps on an average display. All 0’s maps to 0.0f, and all 1’s maps to 1.0f. The sequence of unsigned integer encodings between all 0’s and all 1’s represent a nonlinear progression in the floating-point interpretation of the numbers between 0.0f to 1.0f. For more detail, see the SRGB color standard, IEC 61996-2-1, at IEC (International Electrotechnical Commission).

Conversion from sRGB space is automatically done by `read_imagef` built-in functions if the image channel order is one of the sRGB values described above. When reading from an sRGB image, the conversion from sRGB to linear RGB is performed before the filter specified in the sampler specified to read_imagef is applied. If the format has an alpha channel, the alpha data is stored in linear color space. Conversion to sRGB space is automatically done by `write_imagef` built-in functions if the image channel order is one of the sRGB values described above and the device supports writing to sRGB images.

If the format has an alpha channel, the alpha data is stored in linear color space.

The following is the conversion rule for converting a normalized 8-bit unsigned integer sRGB color value to a floating-point linear RGB color value using `read_imagef`.

```c
if (c <= 0.04045),
    result = c / 12.92;
else
    result = powr((c + 0.055) / 1.055, 2.4);
```

The resulting floating point value, if converted back to an sRGB value without rounding to a 8-bit unsigned integer value, must be within 0.5 ulp of the original sRGB value.

The following are the conversion rules for converting a linear RGB floating-point color value (call it c) to a normalized 8-bit unsigned integer sRGB value using `write_imagef`.

```c
if (isnan(c)) c = 0.0;
if (c > 1.0) c = 1.0;
else if (c < 0.0) c = 0.0;
```
else if (c < 0.0031308)
    c = 12.92 * c;
else
    c = 1.055 * powr(c, 1.0/2.4) - 0.055;

convert to integer scale i.e. c = c * 255.0
convert to integer:
    c = c + 0.5
    drop the decimal fraction, and the remaining
    floating-point(integral) value is converted directly
to an integer.

The precision of the above conversion should be such that \( \text{fabs(} \text{reference result - integer result}) \leq 0.6. \)

### 8.4 Selecting an Image from an Image Array

Let \((u, v, w)\) represent the unnormalized image coordinate values for reading from and/or writing to a 2D image in a 2D image array.

The 2D image layer selected is computed as:

\[
\text{layer} = \text{clamp(rint(w), 0, d_t - 1)}
\]

Let \((u, v)\) represent the unnormalized image coordinate values for reading from and/or writing to a 1D image in a 1D image array.

The 1D image layer selected is computed as:

\[
\text{layer} = \text{clamp(rint(v), 0, h_t - 1)}
\]